A tour problem on a toroidal board

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Abstract

In this paper we study a tour problem that we came across while studying biembeddings and Heffter arrays; see [D.S. Archdeacon, *Electron. J. Combin.* 22 (2015) #P1.74]. Let A be an $n \times m$ toroidal array consisting of filled cells and empty cells. Assume that an orientation $R = (r_1, \ldots, r_n)$ of the cells in each row and an orientation $C = (c_1, \ldots, c_m)$ of the cells in each column of A are fixed. Given an initial filled cell (i_1, j_1) , consider the list $L_{R,C} = ((i_1, j_1), (i_2, j_2), \ldots, (i_k, j_k), (i_{k+1}, j_{k+1}), \ldots)$, where j_{k+1} is the column index of the filled cell (i_k, j_{k+1}) of the row R_{i_k} next to (i_k, j_k) in the orientation r_{i_k} , and where i_{k+1} is the row index of the filled cell of the column $C_{j_{k+1}}$ next to (i_k, j_{k+1}) in the orientation $c_{j_{k+1}}$. We propose the following.

Crazy Knight's Tour Problem. Do there exist R and C such that the list $L_{R,C}$ covers all the filled cells of A?

Here we provide a complete solution for the case with no empty cells, and we obtain partial results for square arrays where the filled cells follow some specific regular patterns.

1 Introduction

Chessboards have been a classical setting for many challenging problems starting from very ancient times up to the most recent ones, for a survey see [19, 24]. Most of these problems are not only interesting *per se*, but have many applications and often arise from a mathematical context. For example classical problems of covering chessboards with the minimum number of chess pieces are related to the study of dominating sets in graphs, see for instance [9, 10]. As well as domination problems, even tour problems are much studied. The most famous problem of this type is the *Knight's tour problem* (see for instance [11]), whose goal is to determine a series of moves made by a knight so that it visits every square on a chessboard exactly once. Then several variations of the original statement have been proposed and investigated, see [4, 8, 17, 20, 21, 22] and the references therein.

Our aim is to find out how to obtain a closed tour on an $n \times m$ board wrapped onto a torus so that the *n* rows and the *m* columns go around the torus. Hence, in the following we will consider all the row indices modulo *n* and all the column indices modulo *m*. Also, the general board we consider might have some squares removed, and we thus represent it as a matrix *A* with some empty and some filled cells, naming it a *partially filled array*. The tour has to visit every filled cell exactly once. Since we are only interested in which cells are empty and which filled, but not in the elements of *A*, we will denote the filled positions simply with a \bullet .

In order to state our problem we introduce a move function which defines the rules for the considered tour on the board. Given two positive integers a and b with $a \leq b$, by [a, b] we will mean the set $\{a, a + 1, \ldots, b\}$. Let us consider a partially filled $n \times m$ array A; we denote by F(A) the subset of $[1, n] \times [1, m]$ given by the filled positions of A. Given $(i, j) \in F(A)$, we define the row successor $s_r((i, j))$ to be (i, j + k) where $k \geq 1$ is the minimum value such that $(i, j + k) \in F(A)$. Similarly we define the column successor $s_c((i, j))$ to be (i + k, j) where $k \geq 1$ is the minimum value such that $(i + k, j) \in F(A)$. Given two vectors $R := (r_1, \ldots, r_n) \in \{-1, 1\}^n$ and $C := (c_1, \ldots, c_m) \in \{-1, 1\}^m$, we denote $S_{R,C} : F(A) \to F(A)$ to be the move function, defined as follows:

$$S_{R,C}((i,j)) := \begin{cases} s_c(s_r((i,j))) \text{ if } r_i = 1 \text{ and } c_{j'} = 1 \text{ where } s_r((i,j)) = (i,j'); \\ s_c(s_r^{-1}((i,j))) \text{ if } r_i = -1 \text{ and } c_{j'} = 1 \text{ where } s_r^{-1}((i,j)) = (i,j'); \\ s_c^{-1}(s_r((i,j))) \text{ if } r_i = 1 \text{ and } c_{j'} = -1 \text{ where } s_r((i,j)) = (i,j'); \\ s_c^{-1}(s_r^{-1}((i,j))) \text{ if } r_i = -1 \text{ and } c_{j'} = -1 \text{ where } s_r^{-1}((i,j)) = (i,j'). \end{cases}$$

We can see the vector R as a choice of the direction of each row: from left to right if $r_i = 1$ and from right to left if $r_i = -1$. Similarly the vector C can be seen as a choice of the direction of each column: from top to bottom if $c_i = 1$ and in the reverse way if $c_i = -1$. Moreover, we can give the following interpretation to the move $S_{R,C}((i, j))$: from the position (i, j) we move first in the *i*-th row following the direction of r_i and then, from the arrival position (i, j') we move in the *j'*-th column in direction of $c_{j'}$ arriving in the position $S_{R,C}((i, j))$. For any $(i, j) \in F(A)$, we set

$$L(i,j) := ((i,j), S_{R,C}((i,j)), S_{R,C}^2((i,j)), \dots, S_{R,C}^p((i,j))),$$

where p is the minimum positive integer such that $S_{R,C}^{p+1}((i,j)) = (i,j)$, namely L(i,j) is the tour that we obtain starting from the cell (i,j). It is natural to ask whether, in this way, we can cover all the filled positions or not.

Crazy Knight's Tour Problem. Given a partially filled $n \times m$ array A, determine whether there exist vectors $R \in \{-1, 1\}^n$ and $C \in \{-1, 1\}^m$ such that, given a filled position (i, j), the list L(i, j) covers all the filled positions of A.

By P(A) we will denote the Crazy Knight's Tour Problem for a given array A. Also, given a filled cell (i, j), if L(i, j) covers all the filled positions of A we will say that the vectors R and C are a solution of P(A).

In this paper, firstly, we will present some necessary conditions for the existence of a solution of P(A) where A is a given array, see Section 2. Then, in Section 3, we present a complete solution when A is a totally filled rectangular array. We focus also on square arrays with some empty cells. In particular in Section 4 and in Section 5 we obtain partial results when A has exactly k filled cells in each row and in each column and when A has exactly k filled cells in each row and in each for one row and column which have k + 1 filled positions, respectively. Finally, in Section 6, we present a recursive construction which allows us to obtain other infinite classes of arrays A such that P(A) has a solution.

1.1 Motivation

In [1], Archdeacon introduced the concept of a Heffter array as a useful tool for determining biembeddings of complete graphs, that is 2-colorable embeddings. In particular, in that paper the author investigated the case in which the face boundaries of the two colour classes form two cycle systems, for a survey see [18]. Heffter arrays which give rise to biembeddings have been constructed in [7, 12, 14, 15]. In particular, in [12] the authors introduced the class of *globally simple* Heffter arrays and, implicitly, studied the Crazy Knight's Tour Problem in some special instances (see [12, Propositions 3.4 and 3.6]) obtaining in such a way new biembeddings (see [12, Theorem 1.11]). The relationship between globally simple Heffter arrays, Crazy Knight's Tour Problem and biembeddings is explained in the following result, that is a reformulation of [1, Theorem 1.1], in the case of square globally simple Heffter arrays, in terms of P(A). Clearly, a similar theorem holds in the rectangular case.

Theorem 1.1. Let A be a globally simple Heffter array $n \times n$ such that each row and each column has exactly k filled cells. If there exists a solution of P(A), then there exists a biembedding of the complete graph of order 2nk + 1 on orientable surface whose face boundaries are k-cycles.

We point out that Heffter arrays are considered interesting as combinatorial objects on their own and not only in view of their relationship with biembeddings. For instance, in [2, 3, 5, 6, 13, 16, 23] the authors construct new infinite classes of Heffter arrays without investigating biembeddings. Analogously, we believe that also the Crazy Knight's Tour Problem is interesting *per se*, as many other tour problems which have been largely studied.

2 Preliminary considerations

Clearly, we can suppose that each row and each column of A has at least one filled position. We note that if the *i*-th row R_i or the *j*-th column C_j has exactly just one filled position (i, j) then $s_r((i, j)) = (i, j)$ or $s_c((i, j)) = (i, j)$, respectively. Also if there exist a row R_i and a column C_j such that $R_i \cup C_j$ has exactly one filled position, then this position is (i, j), since we have assumed that we have no empty row or empty column. In this case, since $S_{R,C}((i, j)) = (i, j)$, it is immediate that the Crazy Knight's Tour Problem has no solution, except for the trivial case in which Ais a square array of size 1.

Remark 2.1. Given an array A and a filled position (i, j), the list L(i, j) covers all the filled positions of A if and only if the list L(i', j') covers all the filled positions of A, for any $(i', j') \in F(A)$. Since we are dealing with toroidal boards, without loss of generality we may assume that $(1, 1) \in F(A)$, namely that the cell (1, 1) of the given array A is filled.

Example 2.2. Let A be

•			•
	٠	٠	
	•	•	•
٠		٠	
~	1.		

Choosing R := (-1, 1, 1, -1) and C := (1, -1, 1, 1) we can cover all the filled positions of A, as shown in the table below where in each filled position we write j if we reach that position after having applied $S_{R,C}$ to (1, 1) exactly j times. Here we represent the elements of R and C by an arrow.

	\downarrow	\uparrow	\downarrow	\downarrow
\rightarrow	0			4
\rightarrow		2	6	
\rightarrow		7	3	1
\leftarrow	5		8	

Remark 2.3. Given a partially filled $n \times m$ array A, let $R := (r_1, r_2, \ldots, r_n)$ and $C := (c_1, c_2, \ldots, c_m)$ be a solution of P(A). Note that this does not imply that $R^i := (r_i, \ldots, r_n, r_1, \ldots, r_{i-1})$ and $C^j := (c_j, \ldots, c_m, c_1, \ldots, c_{j-1})$ are a solution of P(A) too, where $i = 2, \ldots, n$ and $j = 2, \ldots, m$.

Example 2.4. Let A be the array of Example 2.2. If instead of R = (-1, 1, 1, -1) and C = (1, -1, 1, 1) we take $R^2 = (1, 1, -1, -1)$ and $C^2 = (-1, 1, 1, 1)$ it is easy to see that we do not cover all the filled cells of A.

In order to determine the necessary conditions for the existence of a solution of the Crazy Knight's Tour Problem, we introduce the notion of *closed subarray*.

Definition 2.5. Let A be an $n \times m$ toroidal array with no empty row or column. Let \mathcal{R} be a list of rows of A and let \mathcal{C} be a list of columns of A. We say that the subarray $\mathcal{R} \cap \mathcal{C}$ of A is *closed* if $F(\mathcal{R} \cap \mathcal{C}) = F(\mathcal{R} \cup \mathcal{C})$. We say that a closed subarray is minimal if it is minimal with respect to the inclusion. **Example 2.6.** Let A be the following 7×9 array.

			•					
•		۲		٠		٠	۲	٠
			•		•			
٠	•	•		•		•	•	•
•	٠	٠		•		•		
			•		•			
			٠		٠			

Consider the lists $\mathcal{R} = (R_2, R_4, R_5)$ and $\mathcal{C} = (C_1, C_2, C_3, C_5, C_7, C_8, C_9)$. Then it is easy to see that $\mathcal{R} \cap \mathcal{C}$ is a closed subarray of A.

Theorem 2.7. Let A be an $n \times m$ array with |F(A)| filled cells. Necessary conditions for the existence of a solution of P(A) are:

- 1) the array A is a minimal closed subarray;
- 2) $|F(A)| \equiv m + n 1 \pmod{2}$.

Proof. 1) Let $\mathcal{R} \cap \mathcal{C}$ be a closed subarray of A such that $\mathcal{R} \cap \mathcal{C} \neq A$. Since A has no empty row and no empty column, $F(\mathcal{R} \cap \mathcal{C}) \subset F(A)$. It is immediate to see that for any $(i, j) \in F(\mathcal{R} \cap \mathcal{C}), L(i, j) \subseteq F(\mathcal{R} \cap \mathcal{C})$, hence it does not cover all the filled positions of A.

2) The function $S_{R,C}$ defines a permutation ω on F(A). Let ω_r and ω_c be the permutations on the rows and on the columns, respectively, obtained from the move function. By definition of $S_{R,C}$, we have $\omega = \omega_c \circ \omega_r$. Hence the parity of ω , that is |F(A)| - 1, has to be equal to the parity of $\omega_c \circ \omega_r$, that is $\sum_{i=1}^n (|F(R_i)| - 1) + \sum_{j=1}^m (|F(C_j)| - 1) = \sum_{i=1}^n |F(R_i)| + \sum_{j=1}^m |F(C_j)| - n - m = 2|F(A)| - n - m$. It follows that $|F(A)| \equiv n + m - 1 \pmod{2}$.

Corollary 2.8. Let A be a square array with |F(A)| filled cells. If there exists a solution of P(A) then |F(A)| is odd.

3 Totally filled arrays

In this section, we solve Problem P(A) when A is a totally filled $n \times m$ array (i.e. |F(A)| = mn). In this case the move function is simply given by:

$$S_{R,C}((i,j)) := \begin{cases} (i+1,j+1) \text{ if } r_i = 1 \text{ and } c_{j+1} = 1; \\ (i+1,j-1) \text{ if } r_i = -1 \text{ and } c_{j-1} = 1; \\ (i-1,j+1) \text{ if } r_i = 1 \text{ and } c_{j+1} = -1; \\ (i-1,j-1) \text{ if } r_i = -1 \text{ and } c_{j-1} = -1. \end{cases}$$

We first introduce a useful lemma.

Lemma 3.1. Let A be a totally filled $n \times m$ array. Let R := (1, 1, ..., 1) and $C := (1, ..., c_{m-l} = 1, c_{m-l+1} = -1, ..., -1)$, where $0 \le l \le m$. Then R and C are a solution of P(A) if and only if gcd(m - 2l, n) = 1.

Proof. Note that, for i = 1, ..., n, the first m cells of L(i, 1) form the list

$$\Gamma_i := ((i,1), (i+1,2), \dots, (i+m-l-1, m-l), (i+m-l-2, m-l+1), (i+m-l-3, m-l+2), \dots, (i+m-2l-1, m)).$$
(3.1)

Since the Γ_i 's contain one element in each column and are vertical shift of one another, it is clear that

$$\bigcup_{i=1}^{n} \Gamma_i = F(A).$$

Now, since $c_1 = 1$, $L(1, 1) = (\Gamma_1, \Gamma_{1+(m-2l)}, \Gamma_{1+2(m-2l)}, \ldots)$, where the indices are taken modulo n. The indices cover all the values $1, 2, \ldots, n$ if and only if gcd(m - 2l, n) = 1, so this is equivalent to the condition L(1, 1) = F(A).

Example 3.2. Consider the totally filled array A of size $n \times m$ with n = 5 and m = 14. Also let l = 3. In the table below, cells in Γ_i are flagged with γ_i .

γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_2	γ_3	γ_4
γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_3	γ_4	γ_5
γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_4	γ_5	γ_1
γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_5	γ_1	γ_2
γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_4	γ_3	γ_2	γ_1	γ_5	γ_1	γ_2	γ_3

Since gcd(m-2l, n) = gcd(8, 5) = 1 we can apply Lemma 3.1. So we start from the position (1, 1) and apply $S_{R,C}$ until we arrive again in the position (1, 1) of A. In the following table in each position we write j if we reach that position after having applied $S_{R,C}$ to (1, 1) exactly j times.

	\downarrow	\uparrow	\uparrow	\uparrow										
\rightarrow	0	43	16	59	32	5	48	21	64	37	10	39	68	27
\rightarrow	28	1	44	17	60	33	6	49	22	65	38	67	26	55
\rightarrow	56	29	2	45	18	61	34	7	50	23	66	25	54	13
\rightarrow	14	57	30	3	46	19	62	35	8	51	24	53	12	41
\rightarrow	42	15	58	31	4	47	20	63	36	9	52	11	40	69

It is easy to see that we obtain $L(1,1) = (\Gamma_1, \Gamma_4, \Gamma_2, \Gamma_5, \Gamma_3)$.

Theorem 3.3. Let us consider a totally filled $n \times m$ array A. Then there exists a solution of P(A) if and only if n and m are not both even.

Proof. If m and n are both even, the second necessary condition of Theorem 2.7 is not satisfied, hence P(A) has no solution. So suppose that m and n are not both even.

Firstly suppose m odd. Let l be such that m - 2l = 1. Because of Lemma 3.1 we have that R := (1, 1, ..., 1) and $C := (1, ..., c_{m-l} = 1, c_{m-l+1} = -1, ..., -1)$ are a solution of P(A).

Suppose now *m* even, which implies that *n* is odd. Let *l* be such that m - 2l = 2. Because of Lemma 3.1 we have that R := (1, 1, ..., 1) and $C := (1, ..., c_{m-l} = 1, c_{m-l+1} = -1, ..., -1)$ are a solution of P(A).

We point out that the same result of the above theorem, stated in terms of biembeddings, was already obtained in [15, Theorem 2.3].

4 Square arrays with k filled diagonals

In this section we focus on partially filled square arrays having the same number of filled cells in each row and column. Let A be a square array of size n, for i = 1, ..., n we define the *i*-th diagonal D_i to be the set of cells

$$D_i := \{(i, 1), (i + 1, 2), \dots, (i - 1, n)\}$$

where all the arithmetic is performed in \mathbb{Z}_n using the reduced residues $\{1, 2, \ldots, n\}$. We will say that $D_i, D_{i+1}, \ldots, D_{i+k-1}$ are k consecutive diagonals.

Definition 4.1. Let $k \ge 1$ be an integer. We will say that a square array A of size $n \ge k$ is k-diagonal if the non-empty cells of A are exactly those of k diagonals $D_{i_1}, D_{i_2}, \ldots, D_{i_k}$, where $i_1 < i_2 < \ldots < i_k$.

Obviously, if n = k, A is a totally filled array which has already considered in previous section. So, let A be a k-diagonal array of size n > k. In order to find solutions of P(A) we assume that R = (1, ..., 1), even if this is not a necessary condition as shown in the following example.

Example 4.2. Let A be a 3-diagonal array of size 7 whose filled diagonals are D_1, D_2, D_3 . We note that R := (-1, 1, 1, 1, 1, 1, 1) and C := (-1, -1, 1, 1, 1, 1, -1, -1) are a solution of P(A) even though that R is not $(1, \ldots, 1)$. In the following figure we show this graphically by labeling the filled cells as done in previous examples.

		\uparrow	\uparrow	\downarrow	\downarrow	\downarrow	\uparrow	\uparrow
	\leftarrow	0					15	18
	\rightarrow	4	17					20
A ·	\rightarrow	16	11	3				
л.—	\rightarrow		5	12	10			
	\rightarrow			6	13	2		
	\rightarrow				7	14	9	
	\rightarrow					8	19	1

The motivation of this choice for R is explained in the following remark that we state in the case of square arrays, but it holds also for the rectangular ones.

Remark 4.3. Let A be an array of size n such that the cells (i, j), (i, j + 1), (i + 1, j), (i + 1, j + 1) are filled

		j	j+1	
	•	••••		·
i	• • •	٠	•	• • •
i + 1	• • •	•	•	• • •
	·	:	:	·

Then we have that

- $R := (r_1, \ldots, r_{i-1}, r_i = 1, r_{i+1} = -1, r_{i+2}, \ldots, r_n)$ and $C := (c_1, \ldots, c_{j-1}, c_j = -1, c_{j+1} = 1, c_{j+2}, \ldots, c_n)$ are not a solution of P(A) since $S^2_{R,C}((i, j)) = (i, j)$;
- $R := (r_1, \ldots, r_{i-1}, r_i = -1, r_{i+1} = 1, r_{i+2}, \ldots, r_n)$ and $C := (c_1, \ldots, c_{j-1}, c_j = 1, c_{j+1} = -1, c_{j+2}, \ldots, c_n)$ are not a solution of P(A) since $S^2_{R,C}((i+1,j)) = (i+1,j)$.

So, sign changes in both rows and columns can create very short paths on the board whenever there are a relevant number of adjacent filled positions. For this reason, from now on we look for solutions with R = (1, ..., 1). Obviously, since A is a k-diagonal array of size n > k, this implies that $C \neq (1, ..., 1)$. We define $E = (e_1, ..., e_t)$, where $e_1 < e_2 < ... < e_t$, to be the list of the positions of -1's in C. For instance if C = (1, -1, 1, 1, -1, -1) we have E = (2, 5, 6).

4.1 Cyclically *k*-diagonal square arrays

In this subsection we investigate the case in which the k filled diagonals $D_{i_1}, D_{i_2}, \ldots, D_{i_k}$ of A are consecutive.

Definition 4.4. Let k be an integer. A square array A of size n > k is said to be *cyclically k-diagonal* if it is *k*-diagonal and the non-empty diagonals are consecutive (modulo n).

Example 4.5. The following is a cyclically 5-diagonal array of size 9 whose filled diagonals are D_8, D_9, D_1, D_2, D_3 .

•	•	•					•	•
•	•	•	•					•
•	•	•	•	٠				
	•	•	•	٠	•			
		•	•	٠	•	•		
			•	٠	•	٠	•	
				٠	٠	٠	•	•
•					٠	•	٠	•
•	•					•	•	•

Given a cyclically k-diagonal array A, we will say that A is written in the standard form if its nonempty diagonals are D_1, \ldots, D_k . Note that the array of Example 4.2 is in standard form, while the array in Example 4.5 is not. In the following, by Remark 2.1, we can assume without loss of generality that a cyclically k-diagonal array A is written in the standard form.

We start, determining a necessary condition for the existence of a solution.

Proposition 4.6. Let A be a cyclically k-diagonal array of size n > k. If P(A) admits a solution then n and k are both odd and $k \neq 1$.

Proof. Since |F(A)| = kn, from Corollary 2.8 we have n and k odd. Also, note that if k = 1, then every cell of the filled diagonal of A is a closed subarray. Therefore the thesis follows from Theorem 2.7.

Hence, in the following let $k \geq 3$ be an odd integer.

Lemma 4.7. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of size n > k. Then the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, where the positions of each -1 in C are described by E, are a solution of P(A) if and only if:

- 1) the list E covers all the congruence classes modulo d, where $d = \gcd(n, k-1)$;
- 2) the list L(1,1) covers all the positions of $\{(e,e) \mid e \in E\}$.

Proof. If R and C are a solution, the list L(1, 1) covers all the positions of F(A)and hence all the positions of $\{(e, e) \mid e \in E\} \subseteq D_1$. Let us suppose, for the sake of contradiction, that R and C are a solution and that there exists a congruence class [f]modulo d that is not covered by E. We note that, for $(i, i) \in D_1$ and $i - (k - 1) \notin E$, we have $S_{R,C}((i, i)) = (i - (k - 1), i - (k - 1)) \in D_1$ where $i - (k - 1) \equiv i \pmod{d}$. Therefore the list L(f, f) is contained in D_1 , which is a contradiction.

Conversely, let us suppose, for the sake of contradiction, that the hypotheses 1) and 2) of the statement are satisfied, but L(1,1) does not cover F(A). Let $(f,g) \in F(A) \setminus L(1,1)$ and consider the list L(f,g). Obviously $L(1,1) \cap L(f,g) = \emptyset$, hence, by hypothesis 2)

$$\{(e,e)|e \in E\} \cap L(f,g) = \emptyset. \tag{*}$$

We note that, given a position $(i, j), S_{R,C}((i, j)) \in D_1$ only in the following two cases:

- A) $(i, j) \in D_1$ and $i \notin \{e + (k 1) \mid e \in E\};$
- B) $(i, j) \in D_3$ and $S_{R,C}((i, j)) = (e, e)$ for some $e \in E$.

Since from A) for $(i,i) \in D_1 \setminus \{(e+(k-1), e+(k-1)) \mid e \in E\}$ we have $S_{R,C}((i,i)) = (i-(k-1), i-(k-1))$ and, by hypothesis 1), the list E covers all the congruence classes modulo d, it follows that $L(f,g) \not\subseteq D_1$. Suppose now that there exists $(i',j') \in L(f,g) \setminus D_1$ such that $S_{R,C}((i',j')) \in D_1$. From B) it follows that $S_{R,C}((i',j')) \in \{(e,e) \mid e \in E\}$, but obviously $S_{R,C}((i',j')) \in L(f,g)$, which is in contradiction with (*). Hence $L(f,g) \cap D_1 = \emptyset$.

Now we note that, given $(i, j) \in D_h$ with $h \neq 1$, we have $S_{R,C}((i, j)) = (i+1, j+1) \in D_h$ if $j+1 \notin E$ and $S_{R,C}((i, j)) \in D_{h-2}$ if $j+1 \in E$ where the subscripts of the

diagonals are considered modulo k. Since k is odd, this means that the list L(f,g) will reach every non-empty diagonal and hence also the diagonal D_1 , but this is a contradiction because $L(f,g) \cap D_1 = \emptyset$.

As a consequence we have the following result.

Proposition 4.8. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of size n > k. If gcd(n, k - 1) = 1, then the vectors R := (1, ..., 1) and C := (-1, 1, ..., 1) are a solution of P(A).

We point out that the result of Proposition 4.8 was previously obtained in [12] (see Proposition 3.4).

Given a cyclically k-diagonal array of size n > k and vectors R = (1, ..., 1) and $C \in \{-1, 1\}^n$, whose -1 are in positions $E = (e_1, ..., e_t)$ where $e_1 < e_2 < \cdots < e_t$, we would like to study some properties of the move function $S_{R,C}$. In particular, consider an element $(e, e) \in D_1$ with $e \in E$, there exists a minimum $m \ge 1$ such that $S_{R,C}^m((e, e)) = (e', e')$ for some $e' \in E$. We define the permutation ω_C on E as $\omega_C(e) = e'$. Similarly, given $e \in E$, there exists a minimum $m \ge 1$ such that $e - m(k - 1) \equiv e'' \pmod{n}$ for some $e'' \in E$. We define the permutation $\omega_{1,C}$ on E as $\omega_{1,C}(e) = e''$. Finally we define the permutation $\omega_{2,C}$ on $E = (e_1, \ldots, e_t)$ as $\omega_{2,C}(e_i) = e_{i+(k-1)}$ where the indices are considered modulo t.

Lemma 4.9. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of size n > k. Then the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, whose -1 are in positions $E = (e_1, ..., e_t)$ where $e_1 < e_2 < \cdots < e_t$, are a solution of P(A) if and only if:

- 1) the list E covers all the congruence classes modulo d, where $d = \gcd(n, k-1)$;
- 2) the permutation $\omega_{2,C} \circ \omega_{1,C}$ on E is a cycle of length t = |E|.

Proof. Given $e \in E$, the second cell of the form (e', e') with $e' \in E$ we meet in the list L(e, e) is reached after the following moves:

- from (e, e) we move backward into the diagonal D_1 with steps of length k 1 until we reach a cell of the form $(e_i + (k 1), e_i + (k 1))$ with $e_i \in E$;
- from $(e_i + (k-2), e_i) = S_{R,C}(e_i + (k-1), e_i + (k-1))$ we move forward into the diagonal D_{k-1} with steps of length 1 until we reach the cell $(e_{i+1} - 1 + (k-2), e_{i+1} - 1)$, where the indices are considered modulo t (as for the rest of this proof);
- from $(e_{i+1}+(k-4), e_{i+1}) = S_{R,C}(e_{i+1}-1+(k-2), e_{i+1}-1)$ we move forward into the diagonal D_{k-3} with steps of length 1 until we reach the cell $(e_{i+2}-1+(k-4), e_{i+2}-1)$; we reiterate this procedure into the diagonals $D_{k-5}, D_{k-7}, \ldots, D_4$;
- since k is odd, we arrive at the cell $(e_{i+(k-3)/2} + 1, e_{i+(k-3)/2}) \in D_2$ from which we move forward with steps of length 1 until we reach the cell $(e_{i+(k-1)/2}, e_{i+(k-1)/2} - 1);$

- from $(e_{i+(k-1)/2} + (k-1), e_{i+(k-1)/2}) = S_{R,C}(e_{i+(k-1)/2}, e_{i+(k-1)/2} 1)$ we move forward into the diagonal D_k with steps of length 1 until we reach the cell $(e_{i+(k+1)/2} - 1 + (k-1), e_{i+(k+1)/2} - 1)$; we reiterate this procedure into the diagonals $D_{k-2}, D_{k-4}, \ldots, D_3$;
- since k is odd, we arrive at the cell $(e_{i+(k-1)}, e_{i+(k-1)}) \in D_1$ that is the second one of the form $(e', e') \in D_1$ with $e' \in E$ we meet in the list L(e, e).

Because of the definitions, $\omega_{1,C}(e) = e_i$ and $\omega_{2,C}(e_i) = e_{i+(k-1)}$. It follows that $\omega_C = \omega_{2,C} \circ \omega_{1,C}$. Since ω_C is a cycle of length t if and only if the list L(1,1) covers all the positions of $\{(e,e) \mid e \in E\}$, the claim follows from Lemma 4.7.

Example 4.10. The following is a cyclically 7-diagonal array of size 13 in standard form. We consider R := (1, ..., 1) and $C \in \{-1, 1\}^n$, whose -1 are in positions E = (1, 2, 3, 4). We can easily check that $\omega_{1,C} = (1, 2, 3, 4)$, $\omega_{2,C} = (1, 3)(2, 4)$ and $\omega_C = \omega_{2,C} \circ \omega_{1,C} = (1, 4, 3, 2)$. Since the hypotheses of Lemma 4.9 are satisfied, R and C are a solution of P(A). We show in the table below the tour we obtain with this solution.

	\uparrow	\uparrow	\uparrow	\uparrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow
\rightarrow	0							31	46	60	75	89	13
\rightarrow	78	65							32	47	61	76	90
\rightarrow	64	52	39							33	48	62	77
\rightarrow	51	38	26	17							34	49	63
\rightarrow	37	25	16	4	19							35	50
\rightarrow	24	15	3	81	5	21							36
\rightarrow	14	2	80	68	82	6	23						
\rightarrow		79	67	54	69	83	7	1					
\rightarrow			53	41	55	70	84	8	66				
\rightarrow				27	42	56	71	85	9	40			
\rightarrow					28	43	57	72	86	10	18		
\rightarrow						$\overline{29}$	44	$\overline{58}$	$\overline{73}$	87	11	$\overline{20}$	
\rightarrow							30	45	$\overline{59}$	$\overline{74}$	88	$1\overline{2}$	$\overline{22}$

Proposition 4.11. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of size n > k. Let us assume that the vectors R := (1, ..., 1) and C are a solution of P(A). Then there exists a solution of P(A') for any cyclically k-diagonal array A' of size $n' = n + \lambda(k - 1)$ for any integer $\lambda \ge 0$.

Proof. Let us consider the vector R' := (1, ..., 1) and the vector $C' \in \{-1, 1\}^{n'}$ that has the -1 in the same positions, denoted by $E = (e_1, ..., e_t)$, of C. Then the permutations $\omega_{1,C}$ and $\omega_{1,C'}$ on E are identical and the same holds for the permutations $\omega_{2,C}$ and $\omega_{2,C'}$. Thus $\omega_C = \omega_{C'}$. Since $gcd(n, k - 1) = gcd(n + \lambda(k - 1), k - 1)$, the claim follows from Lemma 4.9.

As a consequence of Proposition 4.11 we obtain the following result:

Theorem 4.12. Let $3 \le k < 200$ be an integer and let A be a cyclically k-diagonal array of size n > k. Then there exists a solution of P(A) if and only if n and k are both odd.

Proof. The necessary condition n and k odd follows from Proposition 4.6. For any odd integer k, with $3 \le k < 200$, we checked, using a computer, that the odd values $n \in [k+2, 2k-1]$ have a solution with vector $R = (1, \ldots, 1)$. Hence the thesis follows from Proposition 4.11.

At a first reading it seems that, in Lemma 4.9 (and in the analogous Lemma 4.7), we should also require that n is not an even number, according to Proposition 4.6. In the next result we show that this condition is hidden in the hypothesis of those lemmas.

Proposition 4.13. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of size n > k. Let us consider the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, where the positions of each -1 in C are described by E. Let us suppose that the list E covers all the congruence classes modulo d, where d = gcd(n, k - 1). Then the permutation $\omega_{2,C} \circ \omega_{1,C}$ on E can be a cycle of length t = |E| only if n is odd.

Proof. Suppose, for the sake of contradiction, that n is even (which implies d even) and that $\omega_C = \omega_{2,C} \circ \omega_{1,C}$ is a cycle of length t. Denoted by τ the permutation of the list E such that $\tau(e_i) = e_{i+1}$, we have $\omega_{2,C} = \tau^{k-1}$. Therefore $\omega_{2,C}$ is always an even permutation. Let us partition E in the congruence classes modulo d by setting $E^h := \{e_i \in E | e_i \equiv h \pmod{d}\}$. Since $d = \gcd(n, k-1)$, the permutation $\omega_{1,C}$ is the product of cyclic permutations on the sets E^h , each of which has length $|E^h|$. Since d is even and no E^h is empty, $\omega_{1,C}$ has parity $\sum_{h=0}^{d-1}(|E^h|-1) \equiv (\sum_{h=0}^{d-1}|E^h|) - d \equiv t$ (mod 2), which is also the parity of $\omega_{2,C} \circ \omega_{1,C}$. By hypothesis we have assumed that ω_C is a cycle of length t, then it has parity t - 1, but this is a contradiction because $\omega_C = \omega_{2,C} \circ \omega_{1,C}$ should have parity t.

Remark 4.14. Since the hypotheses of Lemma 4.7 imply those of Lemma 4.9, the condition n odd is also hidden in Lemma 4.7.

Now we introduce a more general application of Lemma 4.9 that helps us in finding a solution of the Crazy Knight's Tour Problem whenever n is sufficiently large.

Theorem 4.15. Let $k \ge 3$ be an odd integer and let A be a cyclically k-diagonal array of odd size $n \ge (k-2)(k-1)$. Then there exists a solution of P(A).

Proof. Set again R := (1, ..., 1) and let d = gcd(n, k - 1), because of Proposition 4.8 we can assume that $d \ge 3$. Now we want to define a vector C that satisfies the hypothesis of Lemma 4.9. Since $n \ge (k-2)(k-1) \ge (k-d+1)(k-1)$, there exist k columns e_1, \ldots, e_k such that:

- $e_i = i$ for any $i \in [1, d 1];$
- $e_i = d + (k-1)(i-d)$ for any $i \in [d, k]$.

Consider now the vector $C \in \{-1,1\}^n$, whose -1 are in positions (e_1,\ldots,e_k) . Clearly the set $\{e_1,\ldots,e_k\}$ covers all the congruence classes modulo d. Since the e_i 's in the congruence class of d are at distance k-1, $\omega_{1,C}$ is the cyclic permutation $(e_d, e_k, e_{k-1}, \ldots, e_{d+1})$. In this context t = |E| = k and $k-1 \equiv -1 \pmod{t}$, hence we have that $\omega_{2,C}$ is the cyclic permutation $(e_1, e_k, e_{k-1}, \ldots, e_2)$. Since k - d + 1 is odd, one can easily check that $\omega_{2,C} \circ \omega_{1,C}$ is given by the cyclic permutation:

$$(e_1, e_k, e_{k-2}, \dots, e_{d+2}, e_d, e_{k-1}, e_{k-3}, \dots, e_{d+1}, e_{d-1}, \dots, e_2)$$

Therefore, because of Lemma 4.9, the vectors R and C are a solution for the Crazy Knight's Tour Problem for any cyclically k-diagonal array of size n where $n \ge (k - 2)(k - 1)$.

4.2 *k*-diagonal square arrays

In this subsection we consider the case in which the filled diagonals of A are not necessarily consecutive. In order to study the solutions of P(A) it is important to know the number of empty diagonals between two filled diagonals, we thus introduce the notion of *empty strip*.

Example 4.16. Here we have a 5-diagonal array of size 11, the non-empty diagonals are D_1 , D_4 , D_6 , D_7 and D_{11} .

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	٠			•	٠				•	•
•		•			•	•				•
٠	•		•			•	•			
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Let A be a k-diagonal array of size n > k. A set $S = \{D_{r+1}, D_{r+2}, \ldots, D_{r+t}\}$ is said to be an *empty strip of width* t if $D_{r+1}, D_{r+2}, \ldots, D_{r+t}$ are empty diagonals, while D_r and D_{r+t+1} are filled diagonals. The array of Example 4.16 has three empty strips $S_1 = \{D_2, D_3\}, S_2 = \{D_5\}$ and $S_3 = \{D_8, D_9, D_{10}\}.$

Definition 4.17. Let A be a k-diagonal array of size n > k. We will say that A is a k-diagonal array with width s if all the empty strips of A have width s.

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	•	•	•				٠	•		
		•	•	•				•	•	
			•	•	٠				•	٠
٠				•	٠	•				•
٠	•				٠	•	٠			
	•	•				•	•	•		
		•	•				•	•	•	
			•	•				•	•	•

Example 4.18. The following is a 5-diagonal array of size 11 with width 3.

We will say that a k-diagonal array A is written in standard form if D_1 is a filled diagonal and D_n is an empty diagonal. For instance the array of Example 4.16 is not in a standard form, while that of Example 4.18 is written in a standard form. In the following we will suppose that A is written in a standard form, since this is not restrictive in order to study the solution of P(A).

Now we present a result which generalizes Lemma 4.7.

Lemma 4.19. Let $k \ge 3$ be an odd integer and let A be a k-diagonal array of size n > k and width s. Then the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, where the positions of each -1 in C are described by E, are a solution of P(A) if and only if:

1) the list E covers all the congruence classes modulo d, where d = gcd(n, k-1);

2) the list L(1,1) covers all the positions of $\{(e,e)|e \in E\}$.

Proof. Let $D_{i_1} = D_1, D_{i_2}, \ldots, D_{i_k}$ with $i_j < i_{j+1}$, for any $j = 1, \ldots, k-1$, be the k non-empty diagonals of A.

If R and C are a solution, the list L(1,1) covers all the filled positions of Aand hence all the positions of $\{(e, e) | e \in E\} \subseteq D_1$. Let us suppose, for the sake of contradiction, that R and C are a solution and that there exists a congruence class [f] modulo d that is not covered by the list E. We note that, for $(i, i) \in D_1$ and $i + s + 1 \notin E$, we have $S_{R,C}((i, i)) = (i + s + 1, i + s + 1) \in D_1$ where $i + s + 1 \equiv i$ (mod d). Therefore the list L(f, f) is contained in D_1 , which is a contradiction.

Conversely, let us suppose, for the sake of contradiction, that the hypotheses 1) and 2) of the statement are satisfied, but L(1,1) does not cover F(A). Let $(f,g) \in F(A) \setminus L(1,1)$ and consider the list L(f,g); obviously $L(1,1) \cap L(f,g) = \emptyset$, hence by 2) we have that

$$\{(e,e)|e \in E\} \cap L(f,g) = \emptyset. \tag{(*)}$$

We note that, given a position $(i, j), S_{R,C}((i, j)) \in D_1$ only in the following two cases:

- A) $(i, j) \in D_1$ and $i + s + 1 \notin E$;
- B) $(i, j) \in D_{i_3}$ and $S_{R,C}((i, j)) = (e, e)$ for some $e \in E$.

Since for $i + s + 1 \notin E$ we have $S_{R,C}((i,i)) = (i + s + 1, i + s + 1)$ and, by hypothesis 1), the list E covers all the congruence classes modulo d, it follows that $L(f,g) \not\subseteq D_1$. Suppose now that there exists $(i', j') \in L(f,g) \setminus D_1$ such that $S_{R,C}((i', j')) \in D_1$. From B) it follows that $S_{R,C}((i', j')) \in \{(e, e) | e \in E\}$, but obviously $S_{R,C}((i', j')) \in L(f,g)$, which is in contradiction with (*). Hence $L(f,g) \cap D_1 = \emptyset$. Given $(i,j) \in D_{i_h}$, set $S_{R,C}((i,j)) = (i'',j'')$. We note that $S_{R,C}((i,j)) \in D_{i_h}$ if $j'' \notin E$, otherwise $S_{R,C}((i,j)) \in D_{i_{h-2}}$ where the subscripts of the diagonals are considered modulo k. Since k is odd it means that the list L(f,g) will reach every non-empty diagonal and hence also the diagonal D_1 , but this is a contradiction because $L(f,g) \cap D_1 = \emptyset$. \Box

The following result, obtained independently and in a more general form in [7] (see Lemma 2.2), is an easy consequence of the previous lemma.

Proposition 4.20. Let $k \ge 3$ be an odd integer and let A be a k-diagonal array of size n > k and width s. If gcd(n, s + 1) = 1, then the vectors R := (1, ..., 1) and C := (-1, 1, ..., 1) are a solution of P(A).

Proof. It is immediate to see that conditions 1) and 2) of Lemma 4.19 are satisfied.

If the array A has exactly two strips of the same width the following holds.

Corollary 4.21. Let $n > k \ge 3$ be odd integers and let A be a k-diagonal array of odd size n and width $s = \frac{n-k}{2}$. If gcd(n, k-2) = 1, then the vectors R := (1, ..., 1) and C := (-1, 1, ..., 1) are a solution of P(A).

Proof. Since n and k are odd, gcd(n, k-2) = 1 if and only if $gcd(n, \frac{n-k+2}{2}) = 1$. Note that $\frac{n-k+2}{2} = s+1$, hence the thesis follows from Proposition 4.20.

Example 4.22. The array of Example 4.18 satisfies the hypotheses of Corollary 4.21 with n = 11, k = 5 and s = 3. We show the tour obtained with the solution presented in the corollary writing in a position j if we reach that position after having applied $S_{R,C}$ to (1, 1) exactly j times.

	\uparrow	\downarrow									
\rightarrow	0				37	15				53	32
\rightarrow	22	3				38	18				54
\rightarrow	44	23	6				39	21			
\rightarrow		45	24	9				40	13		
\rightarrow			46	25	1				41	16	
\rightarrow				47	26	4				42	19
\rightarrow	11				48	27	7				43
\rightarrow	33	14				49	28	10			
\rightarrow		34	17				50	29	2		
\rightarrow			35	20				51	30	5	
\rightarrow				36	12				52	31	8

5 Cyclically almost k-diagonal square arrays

We have seen that the necessary conditions for the existence of a solution of the problem when we have a k-diagonal array of size n are n and k odd. Hence if we want to obtain some results about arrays of even size and/or with an even number of filled diagonals, we have to change the structure of filled cells. In order to break the symmetry as little as possible we consider an array with k filled diagonals and another "extra" filled position.

Definition 5.1. Let $k \ge 1$ be an integer. We will say that a square array A of size n > k is almost k-diagonal if the non-empty cells of A are exactly those of k distinct diagonals $D_{i_1}, D_{i_2}, \ldots, D_{i_k}$ together with another extra filled position belonging to a diagonal D_j with $j \ne i_1, \ldots, i_k$.

In other words, an almost k-diagonal array can be obtained by adding an extra filled cell to a k-diagonal array, hence it has exactly kn + 1 filled positions.

Example 5.2. The following is an almost 2-diagonal array of size 5, whose totally filled diagonals are D_3 and D_4 , and the extra filled position is (4, 4).

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First of all, we determine a necessary condition for the existence of a solution of P(A) when A is an almost k-diagonal array.

Proposition 5.3. Let A be an almost k-diagonal array of size n > k. If P(A) admits a solution then kn is even with $k \neq 1$ or (n, k) = (2, 1).

Proof. Since |F(A)| = kn + 1, from Corollary 2.8 we have that kn must be even. Then note that for k = 1, A admits nontrivial closed subarrays except when n = 2. Hence the thesis follows from Theorem 2.7.

Definition 5.4. Let A be an almost k-diagonal array of size n > k. We will say that A has width s if A can be obtained by adding an extra filled cell to a k-diagonal array with width s.

An almost k-diagonal array A of width s is said to be in standard form if the extra filled cell is $(1, \ell)$ with $2 \le \ell \le s + 1$, D_1 is a totally filled diagonal, D_n has at most (1, 2) as filled position (note that this happens only when $\ell = 2$). As before, without loss of generality we may assume that A is written in standard form.

Also in this section we will look for solutions with R = (1, ..., 1). As done in Section 4, by E we will denote the list of positions of -1's in C. Moreover by C_h we will mean the set of the cells (j, j) of D_1 such that $j \equiv h \pmod{d}$, where $d = \gcd(n, s + 1)$. **Lemma 5.5.** Let $k \ge 3$ be an odd integer and let A be an almost k-diagonal array of size n > k written in standard form and with width s, whose extra filled position is $(1, \ell)$. Then the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, where the positions of each -1 in C are described by E, and with $c_{\ell} = 1$, are a solution of P(A) if and only if:

- 1) the list E covers all the congruence classes modulo d if $\ell \equiv 1 \pmod{d}$, where $d = \gcd(n, k-1)$, otherwise the list E covers all the congruence classes modulo d except at most one among the classes of 1 and of ℓ ;
- 2) the list L(1,1) covers all the positions of $\{(e,e)|e \in E\}$.

Proof. Let $D_{i_1} = D_1, D_{i_2}, \ldots, D_{i_k}$ with $i_j < i_{j+1}$, for any $j = 1, \ldots, k-1$, be the k totally filled diagonals of A.

If R and C are a solution, the list L(1, 1) covers all the filled positions of A and hence all the positions of $\{(e, e) | e \in E\} \subseteq D_1$. Let us suppose, for the sake of contradiction, that R and C are a solution and that one of the following two conditions holds:

- a) there exists a congruence class [f] modulo d, with $f \not\equiv 1, \ell \pmod{d}$, that is not covered by the list E;
- b) the congruence classes of 1 and of ℓ modulo d are not covered by the list E.

Firstly suppose we are in case a). We note that, for $(i,i) \in D_1$ with $i + (s+1) \notin E$ and $i \not\equiv 1 \pmod{d}$, we have $S_{R,C}((i,i)) = (i + (s+1), i + (s+1)) \in D_1$ where $i + (s+1) \equiv i \pmod{d}$. Therefore the list L(f, f) is contained in D_1 , which is a contradiction.

Now suppose we are in case b). Note that $S_{R,C}((1,1)) = (\ell,\ell) \in C_{\ell}$ and that if $(j,j) \in C_{\ell}$ then $S_{R,C}((j,j)) \in C_{\ell}$ except when $j = n + \ell - (s+1)$, in fact $S_{R,C}((n + \ell - (s+1), n + \ell - (s+1)) = (1,\ell)$. Also, since A is in standard form, $S_{R,C}((1,\ell)) = (s+2, s+2) \in C_1$ and if $(j,j) \in C_1$ then $S_{R,C}((j,j)) \in C_1$ except when j = 1. Hence if $\ell \not\equiv 1 \pmod{d}$, then $L(1,1) \subseteq C_1 \cup C_\ell \cup (1,\ell)$, otherwise $L(1,1) \subseteq C_1 \cup (1,\ell)$. In both cases we obtain a contradiction.

Conversely, let us suppose, for the sake of contradiction, that conditions 1) and 2) of the statement are satisfied, but L(1,1) does not cover F(A), and let $(f,g) \in F(A) \setminus L(1,1)$. Let us consider the list L(f,g), obviously $L(1,1) \cap L(f,g) = \emptyset$, hence by 2) we have that

$$\{(e,e)|e \in E\} \cap L(f,g) = \emptyset. \tag{(*)}$$

We note that, given a position (i, j), $S_{R,C}((i, j)) \in D_1$ only in the following four cases:

- A) if $(i,j) = (i,i) \in D_1$ and $i \notin \{e (s+1) | e \in E\} \cup \{1, \ell (s+1)\}$, then $S_{R,C}((i,i)) = (i + (s+1), i + (s+1));$
- B) if (i, j) = (1, 1), then $S_{R,C}((1, 1)) = (\ell, \ell)$;
- C) if $(i, j) = (1, \ell)$ and $s + 2 \notin E$, then $S_{R,C}((1, \ell)) = (s + 2, s + 2);$

D) if
$$(i, j) \in D_{i_3}$$
 and $S_{R,C}((i, j)) = (e, e)$ for some $e \in E$.

If $f \not\equiv 1, \ell \pmod{d}$, from A) and hypothesis 1), it follows that $L(f, g) \not\subseteq D_1 \cup \{(1, \ell)\}$. Now let $f \equiv 1, \ell \pmod{d}$ and assume $L(f, g) \subseteq D_1 \cup \{(1, \ell)\}$. From A), B) and C) this implies that $L(f, g) = C_1 \cup C_\ell \cup \{(1, \ell)\}$. But this happens if and only if the list E does not cover the congruence classes of 1 and of ℓ modulo d, which is in contradiction with hypothesis 1). Hence, in any case, $L(f, g) \not\subseteq D_1 \cup \{(1, \ell)\}$.

Suppose now that there exists $(i', j') \in L(f, g) \setminus \{D_1 \cup \{(1, \ell)\}\}$ such that $S_{R,C}((i', j')) \in D_1 \cup \{(1, \ell)\}$. We note that $S_{R,C}((i', j')) = (1, \ell)$ if and only if $(i', j') = (n + \ell - (s + 1), n + \ell - (s + 1)) \in D_1$, but $(i', j') \notin D_1$. From D) it follows that $S_{R,C}((i', j')) \in \{(e, e) | e \in E\}$, but obviously $S_{R,C}((i', j')) \in L(f, g)$, which is in contradiction with (*). Hence $L(f, g) \cap (D_1 \cup \{(1, \ell)\}) = \emptyset$.

Set $\delta_h = s + 1$ if D_{i_h-1} is empty and $\delta_h = 1$ otherwise. We note that, given $(i, j) \in D_{i_h}$ with $h = 2, \ldots, k$, we have $S_{R,C}((i, j)) = (i + \delta_h, j + \delta_h) \in D_{i_h}$ if $j + \delta_h \notin E$ and $S_{R,C}((i, j)) \in D_{i_{h-2}}$ if $j + \delta_h \in E$ where the subscripts of the diagonals are considered modulo k. Since k is odd it means that the list L(f, g) will reach every totally filled diagonal and hence also the diagonal D_1 , but this is a contradiction because $L(f, g) \cap \{D_1 \cup \{(1, \ell)\}\} = \emptyset$.

As an immediate consequence we have the following result.

Proposition 5.6. Let $k \ge 3$ be an odd integer and let A be an almost k-diagonal array of size n > k and with width s, whose extra filled position is $(1, \ell)$ with ℓ even. If gcd(n, s + 1) = 2, then the vectors R := (1, ..., 1) and C := (-1, 1, ..., 1) are a solution of P(A).

We have an analogous partial result in the case ℓ odd only for almost k-diagonals arrays in which the k totally filled diagonals are consecutive. Hence we introduce the following concept.

Definition 5.7. An almost k-diagonal array A whose k totally filled diagonals are consecutive will be said a cyclically almost k-diagonal array.

Note that in this case if we write the array in standard form the extra filled position is $(1, \ell)$ with $2 \le \ell \le n - k + 1$ and that the width is s = n - k.

Example 5.8. This is a cyclically almost 2-diagonal array of size 5 with extra filled position (1, 4) written in standard form.

•			•	•
٠	•			
	•	•		
		•	•	
			•	٠

From now on we will consider a *cyclically* almost k-diagonal array written in standard form.

Proposition 5.9. Let $k \equiv 3 \pmod{4}$ and let A be a cyclically almost k-diagonal array of size n > k with the extra filled position $(1, \ell)$ with ℓ odd. If gcd(n, k-1) = 2, then the vectors $R := (1, \ldots, 1)$ and $C := (1, -1, 1, \ldots, 1, c_{\ell} = -1, 1, \ldots, 1)$ are a solution of P(A).

Proof. Let R and C be as in the statement. For any i = 2, ..., k we set $D_i^1 = \{(x, y) \in D_i \mid 2 \le y \le \ell - 1\}$ and $D_i^2 = D_i \setminus D_i^1$. Since $k \equiv 3 \pmod{4}$, it is easy to see that

$$L(1,1) = ((1,1), D_k^2, D_{k-2}^1, D_{k-4}^2, D_{k-6}^1, \dots, D_3^2, (2,2), (n+3-k, n+3-k), \dots, (k,k)).$$

Set $B = ((2, 2), \ldots, (k, k)) \subseteq L(1, 1)$. Note that $(2, 2) \in D_1$ hence at least one of (k+1, k+1) and $(\ell+k-1, \ell+k-1)$ is in B, or $B \subseteq D_1$. It is not hard to see that $B \subseteq D_1$ implies that $B = (b_1, b_2, \ldots, b_{|B|})$ where $b_\alpha = (2 - (\alpha - 1)(k - 1), 2 - (\alpha - 1)(k - 1))$ for $\alpha = 1, \ldots, |B|$. But $2 - (\alpha - 1)(k - 1) \neq k$ since k is odd. So $B \not\subseteq D_1$, hence at least one of (k + 1, k + 1) and $(\ell + k - 1, \ell + k - 1)$ is in B. Since $\ell + k - 1$ is odd while k + 1 is even, we have that $B = ((2, 2), \ldots, (k + 1, k + 1), \ldots, (k, k))$ with $((2, 2), \ldots, (k + 1, k + 1)) \subseteq D_1$ which implies:

$$L(1,1) = ((1,1), D_k^2, D_{k-2}^1, D_{k-4}^2, D_{k-6}^1, \dots, D_3^2, (2,2), \dots, (k+1, k+1),$$
$$D_{k-1}^1, D_{k-3}^2, \dots, D_4^2, D_2^1, (1,\ell), (n-k+2, n-k+2), \dots, (k,k)).$$

Set $G = ((n-k+2, n-k+2), \ldots, (k, k)) \subseteq L(1, 1)$. Note that $(n-k+2, n-k+2) \in D_1$ hence either $(\ell + k - 1, \ell + k - 1) \in G$, or $G \subseteq D_1$. It is not hard to see that $G \subseteq D_1$ implies that $G = (g_1, g_2, \ldots, g_{|G|})$ where $g_\alpha = (n-k+2-(\alpha-1)(k-1), n-k+2-(\alpha-1)(k-1))$ for $\alpha = 1, \ldots, |G|$. But $n-k+2-(\alpha-1)(k-1) = n+1-\alpha(k-1) \equiv k$ (mod n) if and only if $\alpha \equiv n/2 - 1 \pmod{n/2}$ that is G coincides with the odd elements of D_1 except (1, 1). So we can suppose that $(\ell + k - 1, \ell + k - 1) \in G$ which implies:

$$L(1,1) = ((1,1), D_k^2, D_{k-2}^1, D_{k-4}^2, D_{k-6}^1, \dots, D_3^2, (2,2), \dots, (k+1,k+1),$$

$$D_{k-1}^1, D_{k-3}^2, \dots, D_4^2, D_2^1, (1,\ell), (n-k+2, n-k+2), \dots, (\ell+k-1,\ell+k-1),$$

$$D_{k-1}^2, D_{k-3}^1, \dots, D_4^1, D_2^2, D_k^1, D_{k-2}^2, \dots, D_5^2, D_3^1, (\ell,\ell), \dots, (k,k)).$$

It is clear that $D_i \subseteq L(1,1)$ for any $i \neq 1$. Suppose by way of contradiction that there exists an element $(j,j) \in D_1$, with $(j,j) \notin L(1,1)$. Hence $L(j,j) \cap L(1,1) = \emptyset$, which implies $L(j,j) \subseteq D_1$ which is a contradiction since gcd(n, k-1) = 2 and $2, \ell \in E$.

Example 5.10. Consider a cyclically almost 7-diagonal array of size 14 in standard form with the extra filled position (1, 5). We show in the table (5.1) the tour obtained

	\downarrow	\uparrow	\downarrow	\downarrow	\uparrow	\downarrow									
\rightarrow	0				50				5	59	88	43	23	77	
\rightarrow	78	26								6	60	89	44	24	
\rightarrow	25	47	52								7	61	90	45	
\rightarrow	46	93	48	28								8	62	91	
\rightarrow	92	65	94	49	96								9	63	
\rightarrow	64	12	66	95	68	30								10	
\rightarrow	11	33	13	67	15	69	98								(5.1)
\rightarrow		79	34	14	36	16	70	32							
\rightarrow			80	35	82	37	17	71	51						
\rightarrow				81	54	83	38	18	72	27					
\rightarrow					1	55	84	39	19	73	53				
\rightarrow						2	56	85	40	20	74	29			
\rightarrow							3	57	86	41	21	75	97		
\rightarrow								4	58	87	42	22	76	31	

in the proof of Proposition 5.9, where cells are tagged as in previous examples.

Now we investigate the case k even. We recall that, in this case, in order to have a solution the size n of an almost k-diagonal array can be both even and odd.

Proposition 5.11. Let A be a cyclically almost 2-diagonal array of size $n \geq 3$ with extra filled cell $(1, \ell)$. The vectors $R := (1, \ldots, 1)$ and $C := (1, \ldots, 1, c_{\ell} = -1, 1, \ldots, 1)$ are a solution of P(A).

Proof. It is sufficient to note that $L(1,1) = ((1,1), (\ell+1,\ell), (\ell+2,\ell+1), (\ell+3,\ell+2), \dots, (n,n-1), (1,n), (2,1), (3,2), \dots, (\ell,\ell-1), (1,\ell), (n,n), (n-1,n-1), \dots, (2,2)).$

For $k \geq 4$, we start by deriving some properties that must be satisfied by any solution.

Proposition 5.12. Let $k \ge 4$ be an even integer and let A be a cyclically almost k-diagonal array of size n > k whose extra filled position is $(1, \ell)$. If the vectors R := (1, ..., 1) and $C \in \{-1, 1\}^n$, where the positions of each -1 in C are described by E, are a solution of P(A) then

- 1) $c_{\ell} = -1;$
- 2) $gcd(|E|, \frac{k}{2}) = 1;$
- 3) the list E covers all the congruence classes modulo d, where $d = \gcd(n, k 1)$, if $\ell \equiv 1 \pmod{d}$, otherwise the list E covers all the congruence classes modulo d except at most the class of 1.

Proof. 1) If, by way of contradiction, we suppose that $c_{\ell} = 1$, then $L(1,1) \subseteq D_1 \cup D_3 \cup \ldots \cup D_{k-1} \cup \{(1,\ell)\}$, hence L(1,1) does not cover F(A).

2) Note that since $k \ge 4$, $(\ell + 2, \ell - 1) \in D_4$ is a filled position hence we can consider the list $L(\ell + 2, \ell - 1)$. If $(1, \ell) \notin L(\ell + 2, \ell - 1)$ then R and C are not a

solution. So, suppose that $(1,\ell) \in L(\ell+2,\ell-1)$. Since, by 1), $c_{\ell} = -1$ we have that $S_{R,C}((i,j)) = (1,\ell)$ if and only if $(i,j) = (\ell,\ell-1)$. Also, note that applying $S_{R,C}$ to a filled cell of the diagonal D_{2a} and of the *j*-th column we obtain a cell of the diagonal D_{2a-2} (modulo k) and of the (j+1)-th column (modulo n) or a cell of D_{2a} . Hence we have $S_{R,C}^{\alpha}((\ell+2,\ell-1)) = (\ell,\ell-1) \in D_2$ implies that $\alpha = \lambda n$ for a suitable λ . We also note that given $(i,j) \neq (\ell,\ell-1)$ and $(i,j) \in D_h$ with $h = 2, \ldots, k, S_{R,C}((i,j)) \in D_h$ if $j+1 \notin E$, otherwise $S_{R,C}((i,j)) \in D_{h-2}$ where the subscripts are taken modulo k. Hence, after having applied $\alpha = \lambda n$ times $S_{R,C}$ to $(\ell+2,\ell-1)$ we are in the diagonal of index $4-2\lambda|E| \pmod{k}$, but we have already observed that $S_{R,C}^{\alpha}((\ell+2,\ell-1)) = (\ell,\ell-1) \in D_2$. So we obtain $4-2\lambda|E| \equiv 2$ (mod k), that is $\lambda|E| \equiv 1 \pmod{\frac{k}{2}}$, which implies that $\gcd(|E|,\frac{k}{2}) = 1$.

3) By way of contradiction, we suppose that there exists a congruence class [i] modulo d, with $i \not\equiv 1 \pmod{d}$ which is not covered by E. Let $(f, f) \in \mathcal{C}_i$ and note that $S_{R,C}((f, f)) \in \mathcal{C}_i$ since $(f, f) \neq (1, 1)$ because $i \not\equiv 1 \pmod{d}$ and since E does not cover the congruence class of i. Hence $L(f, f) \subseteq \mathcal{C}_i$. So L(1, 1) does not cover all the filled position of A.

Proposition 5.13. Let $k \ge 4$ be an even integer and let A be a cyclically almost kdiagonal array of size n > k whose the extra filled position is $(1, \ell)$. If gcd(n, k-1) =1 then the vectors R := (1, ..., 1) and $C := (1, ..., 1, c_{\ell} = -1, 1, ..., 1)$ are a solution of P(A).

Proof. Let R and C be as in the statement. For any i = 2, ..., k we set $D_i^1 = \{(x, y) \in D_i \mid 1 \le y \le \ell - 1\}$ and $D_i^2 = D_i \setminus D_i^1$. It is easy to see that

 $L(1,1) = ((1,1), D_k^2, D_k^1, D_{k-2}^2, D_{k-2}^1, \dots, D_2^2, D_2^1, (1,\ell), (n-k+2, n-k+2), \dots, (k,k)).$

Set $B = ((n-k+2, n-k+2), \dots, (k, k)) \subseteq L(1, 1)$. Note that $(n-k+2, n-k+2) \in D_1$ hence either $(\ell + k - 1, \ell + k - 1) \in B$ or $B \subseteq D_1$. It is not hard to see that $B \subseteq D_1$ implies that $B = (b_1, b_2, \dots, b_{|B|})$ where $b_{\alpha} = (n - k + 2 - (\alpha - 1)(k - 1), n - k + 2 - (\alpha - 1)(k - 1))$ for $\alpha = 1, \dots, |B|$. But $n - k + 2 - (\alpha - 1)(k - 1) \equiv k \pmod{n}$ if and only if $\alpha \equiv n - 1 \pmod{n}$, which implies |B| = n - 1 that is $B = D_1 \setminus (1, 1)$. Since $2 \leq \ell \leq n - k + 1$ we have $\ell + k - 1 \not\equiv 1 \pmod{n}$. So we can suppose that $(\ell + k - 1, \ell + k - 1) \in B$ which implies that

$$L(1,1) = ((1,1), D_k^2, D_k^1, D_{k-2}^2, \dots, D_2^2, D_2^1, (1,\ell), (n-k+2, n-k+2), \dots, (n-k+2), \dots, (n-k+2, n-k+2), \dots, (n-k+2), \dots, (n-$$

$$(\ell + k - 1, \ell + k - 1), D_{k-1}^2, D_{k-1}^1, D_{k-3}^2, \dots, D_3^2, D_3^1, (\ell, \ell), \dots, (k, k)).$$

It is clear that $D_i \subseteq L(1,1)$ for any $i \neq 1$. Suppose, by way of contradiction, that there exists an element $(j,j) \in D_1$, with $(j,j) \notin L(1,1)$. Hence $L(j,j) \cap L(1,1) = \emptyset$ implying $L(j,j) \subseteq D_1$, which is a contradiction since gcd(n, k-1) = 1. \Box

6 A recursive construction

We have already seen in Theorem 2.7 that if A admits nontrivial closed subarrays then P(A) has no solution. Note that if A has exactly t minimal closed subarrays,

adding to A less than t - 1 extra filled cells, we get again an array with nontrivial closed subarrays. Hence in order to obtain a solution of the problem for "arrays with almost closed subarrays" it is necessary to add at least t - 1 filled cells. We start investigating the case in which t = 2 and the arrays are placed diagonally.

Proposition 6.1. Let B_1 and B_2 be two arrays with no empty row and no empty column. Let A be an array of the form

$$A = \begin{array}{c|c} B_1 & E_1 \\ \hline E_2 & B_2 \end{array}$$

where $E_1 \cup E_2$ has exactly one filled position. Then P(A) has a solution if and only if both $P(B_1)$ and $P(B_2)$ do.

Proof. Let (i, j) be the unique filled position of $E_1 \cup E_2$. Firstly, we suppose that $(i, j) \in E_2$.

Let R_t and C_t be a solution of $P(B_t)$ for t = 1, 2. We will show that $R := (R_1, R_2)$ and $C := (C_1, C_2)$ are a solution of P(A). Note that the element (i', j') such that $S_{R,C}((i', j')) = (i, j)$ belongs to B_1 and that $S_{R,C}((i, j)) \in B_2$. Hence the list L(i, j) has elements both in B_1 and in B_2 . Suppose now that there exists $(f, g) \in F(A) \setminus L(i, j)$ and that L(f, g) contains elements both of B_1 and B_2 . Since for any $(i_1, j_1) \in B_1$, $S_{R,C}((i_1, j_1)) \notin B_2$, $(i, j) \in L(f, g)$ which is a contradiction. Hence, L(f, g) contains only elements of B_t for some $t \in \{1, 2\}$. Since R and C restricted to B_t are nothing but R_t and C_t , which solve $P(B_t)$, the list L(f, g) has to contain all the elements of B_t . We deduce that $L(f, g) \cap L(i, j) \neq \emptyset$, which is a contradiction.

Suppose now that R and C are a solution of P(A), namely that L(i, j) covers all the filled positions of A. Since, as already remark, $S_{R,C}((i, j)) \in B_2$ and the element (i', j') such that $S_{R,C}((i', j')) = (i, j)$ belongs to B_1 , we have that L(i, j) = $((i, j), \bar{B}_2, \bar{B}_1)$ where \bar{B}_t is a list of filled cells of B_t , for t = 1, 2. The list L(i, j) covers all the filled positions of A; hence $\bar{B}_t = F(B_t)$ for t = 1, 2, so the restrictions of Rand C to B_t give a solution of $P(B_t)$.

If $(i, j) \in E_1$ the proof can be done in a similar way.

		\uparrow	\uparrow	\downarrow	\downarrow	\downarrow	\uparrow	\uparrow	\downarrow	\downarrow	\downarrow	\uparrow	
	\leftarrow	31					25	28					
	\rightarrow	14	27					30					
	\rightarrow	26	21	13									
	\rightarrow		15	22	20								
A :=	\rightarrow			16	23	33							
	\rightarrow				17	24	19						
	\rightarrow					18	29	32					
	\rightarrow								9	2	7	4	
	\rightarrow			0					5	10	3	12	
	\rightarrow								1	6	11	8	

We have to point out that in Proposition 6.1 it is not necessary to require that the arrays B_1 and B_2 are placed diagonally. In fact the same proof holds also when the rows and the columns of B_1 (resp. B_2) are not consecutive in A, as shown in the following example.

Example 6.3. Let $B_1 = \mathcal{R}_1 \cap \mathcal{C}_1$ where $\mathcal{R}_1 = (R_1, R_2, R_3, R_4, R_5, R_7, R_8)$ and $\mathcal{C}_1 = (C_1, C_2, C_3, C_4, C_5, C_8, C_9)$ and set $B_2 = \mathcal{R}_2 \cap \mathcal{C}_2$ where $\mathcal{R}_2 = (R_6, R_9, R_{10})$ and $\mathcal{C}_2 = (C_6, C_7, C_{10}, C_{11})$. Note that B_1 and B_2 are the same arrays of Example 6.2, but now their rows and columns are not consecutive in A. Also in this case starting from the solutions of $P(B_1)$ and $P(B_2)$ it is possible to obtain a solution of P(A) as shown in table (6.2):

		\uparrow	\uparrow	\downarrow	\downarrow	\downarrow	\downarrow	\downarrow	\uparrow	\uparrow	\downarrow	\uparrow
	\leftarrow	31							25	28		
	\rightarrow	14	27							30		
	\rightarrow	26	21	13								
	\rightarrow		15	22	20							
A :=	\rightarrow			16	23	33						
	\rightarrow						9	2			7	4
	\rightarrow				17	24			19			
	\rightarrow					18			29	32		
	\rightarrow			0			5	10			3	12
	\rightarrow						1	6			11	8

Remark 6.4. Let A be an array with t minimal closed subarrays B_i , for i = 1, ..., t, such that $P(B_t)$ has a solution. Iterating the previous procedure, it is possible to add t-1 extra filled positions (in a suitable way) to A in order to obtain a new array A' such that P(A') has a solution.

(6.1)

(6.2)

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