# On various (strong) rainbow connection numbers of graphs 

Lin Chen Xueliang Li*<br>Center for Combinatorics and LPMC<br>Nankai University, Tianjin 300071<br>China<br>chenlin1120120012@126.com lxl@nankai.edu.cn<br>Henry Liu<br>School of Mathematics<br>Sun Yat-sen University, Guangzhou 510275<br>China<br>liaozhx5@mail.sysu.edu.cn<br>Jinfeng Liu ${ }^{\dagger}$<br>Center for Combinatorics and LPMC<br>Nankai University, Tianjin 300071<br>China<br>ljinfeng709@163.com


#### Abstract

An edge-coloured path is rainbow if all of its edges have distinct colours. For a connected graph $G$, the rainbow connection number $\operatorname{rc}(G)$ of $G$ is the minimum number of colours in an edge-colouring of $G$ such that, any two vertices are connected by a rainbow path. Similarly, the strong rainbow connection number $\operatorname{src}(G)$ of $G$ is the minimum number of colours in an edge-colouring of $G$ such that, any two vertices are connected by a rainbow geodesic (i.e., a path of shortest length). These two concepts of connectivity in graphs were introduced by Chartrand et al. in 2008. Subsequently, vertex-coloured versions of both parameters, $\operatorname{rvc}(G)$ and $\operatorname{srvc}(G)$, and a total-coloured version of the rainbow connection number, $\operatorname{trc}(G)$, were introduced. In this paper we introduce the strong total


[^0]rainbow connection number $\operatorname{strc}(G)$, which is the version of the strong rainbow connection number using total-colourings. Among our results, we will determine the strong total rainbow connection numbers of some special graphs. We will also compare the six parameters, by considering how close and how far apart they can be from one another. In particular, we will characterise all pairs of positive integers $a$ and $b$ such that, there exists a graph $G$ with $\operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b$, and similarly for the parameters rvc and srvc.

## 1 Introduction

In this paper, all graphs under consideration are finite and simple. For notation and terminology not defined here, we refer to [3].

In 2008, Chartrand et al. [6] introduced the concept of rainbow connection of graphs. An edge-coloured path is rainbow if all of its edges have distinct colours. Let $G$ be a non-trivial connected graph. An edge-colouring of $G$ is rainbow connected if any two vertices of $G$ are connected by a rainbow path. The minimum number of colours in a rainbow connected edge-colouring of $G$ is the rainbow connection number of $G$, denoted by $\operatorname{rc}(G)$. The topic of rainbow connection is an active area of research and numerous relevant papers have been published. In addition, the concept of strong rainbow connection was introduced by the same authors. For two vertices $u$ and $v$ of $G$, a $u-v$ geodesic is a $u-v$ path of length $d(u, v)$, where $d(u, v)$ is the distance between $u$ and $v$. An edge-colouring of $G$ is strongly rainbow connected if for any two vertices $u$ and $v$ of $G$, there is a rainbow $u-v$ geodesic. The minimum number of colours in a strongly rainbow connected edge-colouring of $G$ is the strong rainbow connection number of $G$, denoted by $\operatorname{src}(G)$. The investigation of $\operatorname{src}(G)$ is slightly more challenging than that of $\operatorname{rc}(G)$, and fewer results have been obtained on it. For details, see $[6,9,17,24]$.

As a natural counterpart of rainbow connection, Krivelevich and Yuster [14] proposed the concept of rainbow vertex-connection. A vertex-coloured path is vertexrainbow if all of its internal vertices have distinct colours. A vertex-colouring of $G$ is rainbow vertex-connected if any two vertices of $G$ are connected by a vertexrainbow path. The minimum number of colours in a rainbow vertex-connected vertex-colouring of $G$ is the rainbow vertex-connection number of $G$, denoted by $\operatorname{rvc}(G)$. Corresponding to the strong rainbow connection, Li et al. [20] introduced the notion of strong rainbow vertex-connection. A vertex-colouring of $G$ is strongly rainbow vertex-connected if for any two vertices $u$ and $v$ of $G$, there is a vertexrainbow $u-v$ geodesic. The minimum number of colours in a strongly rainbow vertex-connected vertex-colouring of $G$ is the strong rainbow vertex-connection number of $G$, denoted by $\operatorname{srvc}(G)$. For more results on rainbow vertex-connection, we refer to [21, 25].

It was also shown that computing the rainbow connection number and rainbow vertex-connection number of an arbitrary graph is NP-hard [4, 5, 7, 8, 12, 18]. For more results on the rainbow connection subject, we refer to the survey [22] and the
book [23].
Subsequently, Liu et al. [26] proposed the concept of total rainbow connection. A total-coloured path is total-rainbow if its edges and internal vertices have distinct colours. A total-colouring of $G$ is total rainbow connected if any two vertices of $G$ are connected by a total-rainbow path. The minimum number of colours in a total rainbow connected total-colouring of $G$ is the total rainbow connection number of $G$, denoted by $\operatorname{trc}(G)$. For more results on the total rainbow connection number, see [13, 27, 28, 29]. Inspired by the concept of strong rainbow (vertex-)connection, a natural idea is to introduce the strong total rainbow connection number. A totalcolouring of $G$ is strongly total rainbow connected if for any two vertices $u$ and $v$ of $G$, there is a total-rainbow $u-v$ geodesic. The minimum number of colours in a strongly total rainbow connected total-colouring of $G$ is the strong total rainbow connection number of $G$, denoted by $\operatorname{strc}(G)$.

Very recently, Dorbec et al. [10] initiated the study of rainbow connection in digraphs. Subsequently, versions of the other five parameters for digraphs were considered. For more details, see $[1,2,11,15,16]$.

This paper will be organised as follows. In Section 2, we will present results for all six rainbow connection parameters for general graphs. In Section 3, we will determine the strong total rainbow connection number of some specific graphs, including cycles, wheels and complete bipartite and multipartite graphs. Finally in Section 4, we will compare the six parameters, by considering how close and how far apart they can be from one another. In particular, we will characterise all pairs of integers $a$ and $b$ such that, there exists a connected graph $G$ with $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$, and similarly for the parameters trc and strc.

We mention a few more words on terminology and notation. For a graph $G$, its vertex and edge sets are denoted by $V(G)$ and $E(G)$, and its diameter is denoted by $\operatorname{diam}(G)$. Let $K_{n}$ and $C_{n}$ denote the complete graph and cycle of order $n$ (where $n \geq 3$ for $C_{n}$ ), and $K_{m, n}$ denote the complete bipartite graph with class sizes $m$ and $n$. For two graphs $G$ and $H$, and a vertex $u \in V(G)$, we define $G_{u \rightarrow H}$ to be the graph obtained by replacing $u$ with $H$, and replacing the edges of $G$ at $u$ with all edges between $H$ and the neighbours of $u$ in $G$. We say that $G_{u \rightarrow H}$ is obtained from $G$ by expanding $u$ into $H$. Note that the graph obtained from $G$ by expanding every vertex of $G$ into $H$ is also known as the lexicographic product $G \circ H$.

## 2 Remarks and results for general graphs

In this section, we present some results about the six rainbow connection parameters $\operatorname{rc}(G), \operatorname{src}(G), \operatorname{rvc}(G), \operatorname{srvc}(G), \operatorname{trc}(G)$ and $\operatorname{strc}(G)$, for general graphs $G$. Let $G$ be a non-trivial connected graph with $m$ edges and $n$ vertices, where $q$ vertices are non-pendent (i.e., with degree at least 2). We have the following inequalities.

$$
\begin{gather*}
\operatorname{diam}(G) \leq \operatorname{rc}(G) \leq \operatorname{src}(G) \leq m  \tag{1}\\
\operatorname{diam}(G)-1 \leq \operatorname{rvc}(G) \leq \operatorname{srvc}(G) \leq \min (n-2, q)  \tag{2}\\
2 \operatorname{diam}(G)-1 \leq \operatorname{trc}(G) \leq \operatorname{strc}(G) \leq \operatorname{srvc}(G)+m \leq \min (m+n-2, m+q) \tag{3}
\end{gather*}
$$

To see the last inequality of (2), the inequality $\operatorname{srvc}(G) \leq n-2$ is a result of Li et al. [20]. We also have $\operatorname{srvc}(G) \leq q$, since any vertex-colouring of $G$ where all $q$ non-pendent vertices are given distinct colours, is strongly rainbow vertex-connected. To see the third inequality of (3), we may take a strongly rainbow vertex-connected colouring of $G$ with $\operatorname{srvc}(G)$ colours, and then colour the edges with $m$ further distinct colours. This gives a strongly total rainbow connected colouring of $G$ with $\operatorname{srvc}(G)+$ $m$ colours. The last inequality of (3) then follows from (2). All remaining inequalities are trivial.

Also, the following upper bound is obvious.

$$
\operatorname{strc}(G) \leq \operatorname{src}(G)+q .
$$

Indeed, a strongly total rainbow connected colouring of $G$ can be obtained from a strongly rainbow connected colouring with $\operatorname{src}(G)$ colours, and then colouring the non-pendent vertices of $G$ with $q$ further distinct colours. Similarly, for graphs with diameter 2, we have the following proposition which will be very helpful later.

Proposition 2.1. Let $G$ be a graph with diameter 2. Then $\operatorname{strc}(G) \leq \operatorname{src}(G)+1$.
Proof. By definition, we may give $G$ a strongly rainbow connected colouring, using $\operatorname{src}(G)$ colours. Since $\operatorname{diam}(G)=2$, any two non-adjacent vertices $x, y \in V(G)$ are connected by a rainbow $x-y$ geodesic of length 2 . Now, colour all vertices of $G$ with a new colour. Then clearly, the resulting total-colouring uses $\operatorname{src}(G)+1$ colours, and is a strongly total rainbow connected colouring. Thus, $\operatorname{strc}(G) \leq \operatorname{src}(G)+1$.

For the parameters $\operatorname{rc}(G)$ and $\operatorname{trc}(G)$, we have the following upper bounds which are better than those of (1) and (3).

$$
\operatorname{rc}(G) \leq n-1, \quad \text { and } \quad \operatorname{trc}(G) \leq \min (2 n-3, n-1+q)
$$

Indeed, we may take a spanning tree $T$ of $G$, which has $n-1$ edges and at most $\min (n-2, q)$ non-pendent vertices. We can assign distinct colours to all edges of $T$, and to all edges and non-pendent vertices of $T$, to obtain, respectively, the above two upper bounds.

As for alternative lower bounds instead of those involving the diameter, we note that for any total rainbow connected colouring of $G$, the colours of the bridges and cut-vertices must be pairwise distinct. Similar observations hold for rainbow connected and rainbow vertex-connected colourings, where respectively, the colours of the bridges, and the colours of the cut-vertices, must be pairwise distinct. Hence, the following result holds.

Proposition 2.2. Let $G$ be a connected graph. Suppose that $B$ is the set of all bridges, and $C$ is the set of all cut-vertices. Denote $b=|B|$ and $c=|C|$, respectively. Then

$$
\begin{gathered}
\operatorname{src}(G) \geq \operatorname{rc}(G) \geq b \\
\operatorname{srvc}(G) \geq \operatorname{rvc}(G) \geq c \\
\operatorname{strc}(G) \geq \operatorname{trc}(G) \geq b+c .
\end{gathered}
$$

In the next result, we give equivalences and implications when the rainbow connection parameters are small.

Theorem 2.3. Let $G$ be a non-trivial connected graph.
(a) The following are equivalent.
(i) $G$ is a complete graph.
(ii) $\operatorname{diam}(G)=1$.
(iii) $\operatorname{rc}(G)=1$.
(iv) $\operatorname{src}(G)=1$.
(v) $\operatorname{rvc}(G)=0$.
(vi) $\operatorname{srvc}(G)=0$.
(vii) $\operatorname{trc}(G)=1$.
(viii) $\operatorname{strc}(G)=1$.
(b) $\operatorname{strc}(G) \geq \operatorname{trc}(G) \geq 3$ if and only if $G$ is not a complete graph.
(c) (i) $\operatorname{rc}(G)=2$ if and only if $\operatorname{src}(G)=2$.
(ii) $\operatorname{rvc}(G)=1$, if and only if $\operatorname{srvc}(G)=1$, if and only if $\operatorname{diam}(G)=2$.
(iii) $\operatorname{rvc}(G)=2$ if and only if $\operatorname{srvc}(G)=2$.
(iv) $\operatorname{trc}(G)=3$ if and only if $\operatorname{strc}(G)=3$.
(v) $\operatorname{trc}(G)=4$ if and only if $\operatorname{strc}(G)=4$.

Moreover, any of the conditions in (i) implies any of the conditions in (iv), and any of the conditions in (i), (iv) and (v) implies any of the conditions in (ii).

Proof. Although parts of this result can be found in [6, 20], we provide a proof for the sake of completeness.
(a) Clearly we have (ii) $\Rightarrow$ (i) $\Rightarrow$ (iv). Using (1), we can easily obtain (iv) $\Rightarrow$ (iii) $\Rightarrow$ (ii). Similarly, using (2) and (3), we have (ii) $\Rightarrow$ (i) $\Rightarrow$ (vi) $\Rightarrow$ (v) $\Rightarrow$ (ii), and (ii) $\Rightarrow$ (i) $\Rightarrow$ (viii) $\Rightarrow$ (vii) $\Rightarrow$ (ii).
(b) If $G$ is not a complete graph, then $\operatorname{diam}(G) \geq 2$, and $\operatorname{strc}(G) \geq \operatorname{trc}(G) \geq 3$ follows from (3). The converse clearly holds by (a).
(c) We first prove (i). Suppose first that $\operatorname{src}(G)=2$. Then by (a), we have $\operatorname{diam}(G) \geq 2$. By (1), we have $2 \leq \operatorname{rc}(G) \leq \operatorname{src}(G)=2$, and hence $\operatorname{rc}(G)=2$. Conversely, suppose that $\operatorname{rc}(G)=2$. Then (a) and (1) imply that $\operatorname{src}(G) \geq 2$ and $\operatorname{diam}(G)=2$. Also, there exists a rainbow connected colouring for $G$, using $\operatorname{rc}(G)=2$ colours. In such an edge-colouring, for any $x, y \in V(G)$, either $x y \in E(G)$, or $x y \notin E(G)$ and there is a rainbow $x-y$ path of length 2 , which is also a rainbow $x-y$ geodesic. Thus $\operatorname{src}(G) \leq 2$, and $\operatorname{src}(G)=2$ as required.

By similar arguments using (a), (2) and (3) we can prove (iv); that the first two conditions of (ii) are equivalent; and that the first condition of (v) implies the second.

Now we complete the proof of (ii). If $\operatorname{rvc}(G)=1$, then we can easily use (a) and (2) to obtain $\operatorname{diam}(G)=2$. If $\operatorname{diam}(G)=2$, then (2) gives $\operatorname{rvc}(G) \geq 1$. Clearly, the vertex-colouring of $G$ where every vertex is given the same colour is rainbow vertex-connected, and thus $\operatorname{rvc}(G) \leq 1$. Therefore (ii) holds.

Next, we prove (iii). Suppose first that $\operatorname{srvc}(G)=2$. Then $\operatorname{rvc}(G) \leq 2$ by (2). Clearly $\operatorname{rvc}(G) \neq 0$ by (a), and $\operatorname{rvc}(G) \neq 1$ by (c)(ii). Thus $\operatorname{rvc}(G)=2$. Conversely, suppose that $\operatorname{rvc}(G)=2$. Then by $(2)$, we have $\operatorname{srvc}(G) \geq 2$ and $\operatorname{diam}(G) \leq 3$. We may take a rainbow vertex-connected colouring of $G$, using at most $\operatorname{rvc}(G)=2$ colours. Let $x, y \in V(G)$. If $d(x, y) \in\{1,2\}$, then any $x-y$ geodesic is clearly vertex-rainbow. If $d(x, y)=3$, then since any $x-y$ path of length at least 4 cannot be vertex-rainbow, there must exist a vertex-rainbow $x-y$ path of length 3 , which is also an $x-y$ geodesic. Thus, the colouring is also strongly rainbow vertex-connected. We have $\operatorname{srvc}(G) \leq 2$, so that $\operatorname{srvc}(G)=2$, and (iii) holds.

Next, we complete the proof of (v). Suppose that $\operatorname{strc}(G)=4$. By (a) and (b), we have $3 \leq \operatorname{trc}(G) \leq \operatorname{strc}(G)=4$. By (iv), we have $\operatorname{trc}(G) \neq 3$, so that $\operatorname{trc}(G)=4$. Thus (v) holds.

Finally, we prove the last part of (c). Firstly, suppose that either condition in (i) holds, so that $\operatorname{rc}(G)=2$. Then (a) and (b) imply $\operatorname{trc}(G) \geq 3$. Moreover, there exists a rainbow connected edge-colouring for $G$, using $\operatorname{rc}(G)=2$ colours. Clearly by colouring all vertices of $G$ with a third colour, we have a total rainbow connected colouring for $G$, using 3 colours. Thus, $\operatorname{trc}(G) \leq 3$. We have $\operatorname{trc}(G)=3$, and thus both conditions of (iv) hold. Secondly, suppose that any of the conditions in (i), (iv) or (v) holds. It is easy to use (a), and (1) or (3), to obtain $\operatorname{diam}(G)=2$. Thus, the three conditions of (ii) also hold.

Remark. We remark that in Theorem 2.3(c), no other implication exists between the conditions of (i) to (v). Obviously, no implication exists between the conditions of (ii) and those of (iii). Thus by the last part of (c), no implication exists between the conditions of (iii) and those of (i), (iv) and (v). Similarly, no implication exists between the conditions of (iv) and those of (v), and thus no implication exists between the conditions of (i) and those of (v), since the conditions of (i) imply those of (iv). Clearly, the example of the stars $K_{1, n}$ shows that there are infinitely many graphs where the conditions of (ii) hold, but those of (i), (iv) and (v) do not hold. Indeed, for $n \geq 2$, we have $\operatorname{rvc}\left(K_{1, n}\right)=1$, while $\operatorname{rc}\left(K_{1, n}\right)=n$ and $\operatorname{trc}\left(K_{1, n}\right)=n+1$. Now, there are infinitely many graphs $G$ such that the conditions of (iv) hold, but those of (i) do not hold. For example, let $u$ be a vertex of the cycle $C_{5}$, and let $G$ be a graph obtained by expanding $u$ into a clique $K$. That is, $G=\left(C_{5}\right)_{u \rightarrow K}$. It was remarked in [26] (and also easy to show) that for any such graph $G$, we have $\operatorname{trc}(G)=\operatorname{rc}(G)=3$.

Now, it is easy to see that if $H$ is a spanning connected subgraph of a connected graph $G$, then we have

$$
\operatorname{rc}(G) \leq \operatorname{rc}(H), \quad \operatorname{rvc}(G) \leq \operatorname{rvc}(H), \quad \text { and } \quad \operatorname{trc}(G) \leq \operatorname{trc}(H)
$$

However, the following lemma shows that the same inequalities do not hold for the strong rainbow connection parameters.
Lemma 2.4. There exist connected graphs $G$ and $H$ such that $H$ is a spanning subgraph of $G$, and $\operatorname{src}(G)>\operatorname{src}(H)$. Similar statements hold for the parameters srvc and strc.
Proof. We construct graphs $G_{i}$ and $H_{i}$, for $i=1,2,3$, as follows. Let $H_{1}$ (respectively $H_{2}, H_{3}$ ) be the graph as shown in Figure 1(a) (respectively (b), (c)) consisting of the solid edges, and $G_{1}$ (respectively $G_{2}, G_{3}$ ) be the graph obtained by adding the dotted edge.


Figure 1. The graphs in Lemma 2.4
We will prove that

$$
\begin{equation*}
\operatorname{src}\left(G_{1}\right)>\operatorname{src}\left(H_{1}\right), \quad \operatorname{srvc}\left(G_{2}\right)>\operatorname{srvc}\left(H_{2}\right), \quad \text { and } \quad \operatorname{strc}\left(G_{3}\right)>\operatorname{strc}\left(H_{3}\right) \tag{4}
\end{equation*}
$$

Firstly, it is easy to see that the edge-colouring of $H_{1}$ as shown is strongly rainbow connected, and thus $\operatorname{src}\left(H_{1}\right) \leq 4$. In fact, we have $\operatorname{src}\left(H_{1}\right)=4$, since $\operatorname{src}\left(H_{1}\right) \geq \operatorname{diam}\left(H_{1}\right)=4$. Now, suppose that there exists a strongly rainbow connected colouring of $G_{1}$, using at most four colours. Note that the four pendent edges of $G_{1}$ must receive distinct colours, say colours $1,2,3,4$. The dotted edge has colour $1,2,3$ or 4 , and in each case, we can easily find two vertices that are not connected by a rainbow geodesic in $G_{1}$. We have a contradiction, and thus $\operatorname{src}\left(G_{1}\right) \geq 5>4=\operatorname{src}\left(H_{1}\right)$.

We can similarly prove the remaining two inequalities of (4). We have a strongly rainbow vertex-connected colouring of $H_{2}$ as shown, and since diam $\left(H_{2}\right)=7$, we have $\operatorname{srvc}\left(H_{2}\right)=6$. Suppose that there exists a strongly rainbow vertex-connected colouring of $G_{2}$, using at most six colours. Then, the six cut-vertices of $G_{2}$ must receive distinct colours, say colours $1,2,3,4,5,6$. The vertex $x$ has colour $1,2,3,4,5$ or 6 , and in each case, we can find two vertices that are not connected by a vertexrainbow geodesic in $G_{2}$. We have a contradiction, and thus $\operatorname{srvc}\left(G_{2}\right) \geq 7>6=$ $\operatorname{srvc}\left(H_{2}\right)$. Likewise, we have a strongly total rainbow connected colouring of $H_{3}$ as shown, and thus $\operatorname{strc}\left(H_{3}\right) \leq 14$. Suppose that there exists a strongly total rainbow connected colouring of $G_{3}$, using at most 14 colours. Then, the eight bridges and six cut-vertices of $G_{3}$ must receive distinct colours, say colours $1,2, \ldots, 14$. The dotted edge has colour $1,2, \ldots, 13$ or 14 , and in each case, we can find two vertices that are not connected by a total-rainbow geodesic in $G_{3}$. Again we have a contradiction, and thus $\operatorname{strc}\left(G_{3}\right) \geq 15>14 \geq \operatorname{strc}\left(H_{3}\right)$.

Li et al. [20] provided a similar example of graphs $G$ and $H$ which gave $\operatorname{srvc}(G)=$ $9>8=\operatorname{srvc}(H)$. However in their example, $H$ was not a spanning subgraph of $G$, although this could be easily corrected. Chartrand et al. [6] had conjectured that $\operatorname{src}(G) \leq \operatorname{src}(H)$ whenever $G$ and $H$ are connected graphs, with $H$ a spanning subgraph of $G$. They observed that if this conjecture was true, then we have $\operatorname{src}(G) \leq$ $n-1$ if $G$ is a connected graph of order $n$. However, Lemma 2.4 shows that the conjecture is false. The latter claim may still be true, and we propose this as an open problem, as well as the total-coloured analogue.

Problem 2.5. Let $G$ be a connected graph of order $n$ with $q$ non-pendent vertices. Then, are the following inequalities true?

$$
\operatorname{src}(G) \leq n-1, \quad \text { and } \quad \operatorname{strc}(G) \leq \min (2 n-3, n-1+q)
$$

## 3 Strong total rainbow connection numbers of some graphs

In this section, we consider the strong total rainbow connection numbers of some specific graphs, namely, trees, cycles, wheels, and complete bipartite and multipartite graphs. The remaining five rainbow connection parameters for these graphs have previously been considered by various authors, and we shall recall these previous results along the way.

First, let $T$ be a tree of order $n$, with $q$ non-pendent vertices. Note that, since any two vertices of $T$ are connected by a unique path, we have $\operatorname{rc}(T)=\operatorname{src}(T)$, $\operatorname{rvc}(T)=\operatorname{srvc}(T)$, and $\operatorname{trc}(T)=\operatorname{strc}(T)$. From Chartrand et al. [6], and Liu et al. [25, 26], we have $\operatorname{rc}(T)=\operatorname{src}(T)=n-1, \operatorname{rvc}(T)=q$, and $\operatorname{trc}(T)=n-1+q$. Moreover, it is well known that if $n \geq 3$, then $1 \leq q \leq n-2$; and that $q=1$ if and only if $T$ is a star, and $q=n-2$ if and only if $T$ is a path. Thus, we have the following result.

Proposition 3.1. Let $T$ be a tree with order n, and $q$ non-pendent vertices.
(a) $\operatorname{rvc}(T)=\operatorname{srvc}(T)=q$. In particular, for $n \geq 2, \operatorname{rvc}(T)=\operatorname{srvc}(T)=n-2$ if and only if $T$ is a path; and for $n \geq 3, \operatorname{rvc}(T)=\operatorname{srvc}(T)=1$ if and only if $T$ is a star.
(b) $\operatorname{trc}(T)=\operatorname{strc}(T)=n-1+q$. In particular, for $n \geq 2, \operatorname{trc}(T)=\operatorname{strc}(T)=2 n-3$ if and only if $T$ is a path; and for $n \geq 3, \operatorname{trc}(T)=\operatorname{strc}(T)=n$ if and only if $T$ is a star.

Our next task is to consider cycles. Recall that $C_{n}$ denotes the cycle of order $n \geq 3$. The values of $\operatorname{rc}\left(C_{n}\right)$ and $\operatorname{src}\left(C_{n}\right)$ were determined by Chartrand et al. [6], while the values of $\operatorname{rvc}\left(C_{n}\right), \operatorname{srvc}\left(C_{n}\right)$ and $\operatorname{trc}\left(C_{n}\right)$ were determined by Li and Liu [19], Lei et al. [15], and Liu et al. [26], respectively. We may summerise these results as follows.

Theorem 3.2. $[6,15,19,26]$
(a) $\operatorname{rc}\left(C_{3}\right)=\operatorname{src}\left(C_{3}\right)=1$, and $\operatorname{rc}\left(C_{n}\right)=\operatorname{src}\left(C_{n}\right)=\left\lceil\frac{n}{2}\right\rceil$ for $n \geq 4$.
(b) For $3 \leq n \leq 15$, the values of $\operatorname{rvc}\left(C_{n}\right)$ and $\operatorname{srvc}\left(C_{n}\right)$ are given in the following table.

| $n$ | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\operatorname{rvc}\left(C_{n}\right)$ | 0 | 1 | 1 | 2 | 3 | 3 | 3 | 4 | 5 | 5 | 6 | 7 | 7 |
| $\operatorname{srvc}\left(C_{n}\right)$ | 0 | 1 | 1 | 2 | 3 | 3 | 3 | 4 | 6 | 5 | 7 | 7 | 8 |

For $n \geq 16$, we have $\operatorname{rvc}\left(C_{n}\right)=\operatorname{srvc}\left(C_{n}\right)=\left\lceil\frac{n}{2}\right\rceil$.
(c) For $3 \leq n \leq 12$, the values of $\operatorname{trc}\left(C_{n}\right)$ are given in the following table.

| $n$ | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\operatorname{trc}\left(C_{n}\right)$ | 1 | 3 | 3 | 5 | 6 | 7 | 8 | 9 | 11 | 11 |

For $n \geq 13$, we have $\operatorname{trc}\left(C_{n}\right)=n$.
Note that we have the slightly surprising facts that $\operatorname{rc}\left(C_{n}\right)=\operatorname{src}\left(C_{n}\right)$, but $\operatorname{rvc}\left(C_{n}\right)$ $=\operatorname{srvc}\left(C_{n}\right)$ except for $n=11,13,15$; and that $\operatorname{srvc}\left(C_{11}\right)>\operatorname{srvc}\left(C_{12}\right)$. By taking advantage of the fact that $\operatorname{strc}\left(C_{n}\right) \geq \operatorname{trc}\left(C_{n}\right)$ and the proof of part (c) in [26], we have the following result for $\operatorname{strc}\left(C_{n}\right)$.

Theorem 3.3. For $n \geq 3$, we have $\operatorname{strc}\left(C_{n}\right)=\operatorname{trc}\left(C_{n}\right)$. That is, for $3 \leq n \leq 12$, the values of $\operatorname{strc}\left(C_{n}\right)$ are given in the table in Theorem 3.2(c). For $n \geq 13$, we have $\operatorname{strc}\left(C_{n}\right)=n$.

Proof. One can easily check that $\operatorname{strc}\left(C_{3}\right)=1, \operatorname{strc}\left(C_{4}\right)=3$, and $\operatorname{strc}\left(C_{5}\right)=3$. Now, let $n \geq 6$. We need to prove that $\operatorname{strc}\left(C_{n}\right) \leq \operatorname{trc}\left(C_{n}\right)$. Thus by Theorem 3.2(c), we need to prove that $\operatorname{strc}\left(C_{n}\right) \leq n-1$ for $6 \leq n \leq 10$ and $n=12$, and $\operatorname{strc}\left(C_{n}\right) \leq n$ for $n=11$ and $n \geq 13$. The following facts were shown in the proof of Theorem 3.2(c) in [26].

- For $6 \leq n \leq 10$ and $n=12$, there is a total-colouring of $C_{n}$, using $n-1$ colours, such that every path of length $\left\lceil\frac{n}{2}\right\rceil-1$ is total-rainbow, and when $n$ is even, any two opposite vertices of $C_{n}$ are connected by a total-rainbow path.
- For $n=11$ and $n \geq 13$, there is a total-colouring of $C_{n}$, using $n$ colours, such that every path of length $\left\lceil\frac{n}{2}\right\rceil$ is total-rainbow.

With these total-colourings, it is easy to see that any two vertices $x$ and $y$ of $C_{n}$ are connected by a total-rainbow $x-y$ path of length at most $\left\lfloor\frac{n}{2}\right\rfloor$, which must also be a total-rainbow $x-y$ geodesic. Thus the total-colourings are also strong total rainbow connected colourings, and the upper bound $\operatorname{strc}\left(C_{n}\right) \leq \operatorname{trc}\left(C_{n}\right)$ follows.

Next, we consider wheel graphs. The wheel $W_{n}$ of order $n+1 \geq 4$ is the graph obtained from the cycle $C_{n}$ by joining a new vertex $v$ to every vertex of $C_{n}$. The vertex $v$ is the centre of $W_{n}$. Trivially, we have $\operatorname{rvc}\left(W_{3}\right)=\operatorname{srvc}\left(W_{3}\right)=0$, and $\operatorname{rvc}\left(W_{n}\right)=$ $\operatorname{srvc}\left(W_{n}\right)=1$ for $n \geq 4$. The values of $\operatorname{rc}\left(W_{n}\right)$ and $\operatorname{src}\left(W_{n}\right)$ were determined by Chartrand et al. [6], while the values of $\operatorname{trc}\left(W_{n}\right)$ were determined by Liu et al. [26].

Theorem 3.4. [6, 26]
(a) $\operatorname{rc}\left(W_{3}\right)=1, \operatorname{rc}\left(W_{n}\right)=2$ for $4 \leq n \leq 6$, and $\operatorname{rc}\left(W_{n}\right)=3$ for $n \geq 7$.
(b) $\operatorname{src}\left(W_{n}\right)=\left\lceil\frac{n}{3}\right\rceil$ for $n \geq 3$.
(c) $\operatorname{trc}\left(W_{3}\right)=1, \operatorname{trc}\left(W_{n}\right)=3$ for $4 \leq n \leq 6, \operatorname{trc}\left(W_{n}\right)=4$ for $7 \leq n \leq 9$, and $\operatorname{trc}\left(W_{n}\right)=5$ for $n \geq 10$.

In the next result, we determine the values of $\operatorname{strc}\left(W_{n}\right)$. The proof is partially based on the fact that $\operatorname{strc}\left(W_{n}\right) \geq \operatorname{trc}\left(W_{n}\right)$.

Theorem 3.5. $\operatorname{strc}\left(W_{3}\right)=1$, and $\operatorname{strc}\left(W_{n}\right)=\left\lceil\frac{n}{3}\right\rceil+1$ for $n \geq 4$.
Proof. Let $v$ be the centre of $W_{n}$, and $v_{0}, v_{1}, \ldots, v_{n-1}$ be the vertices of $W_{n}$ in the cycle $C_{n}$. Since $W_{3}$ is precisely the complete graph $K_{4}$, we have $\operatorname{strc}\left(W_{3}\right)=1$.

Now, let $n \geq 4$. Since $\operatorname{diam}\left(W_{n}\right)=2$, by Proposition 2.1 and Theorem 3.4(b), we have $\operatorname{strc}\left(W_{n}\right) \leq \operatorname{src}\left(W_{n}\right)+1=\left\lceil\frac{n}{3}\right\rceil+1$. Also, by Theorem 3.4(c), we have $\operatorname{strc}\left(W_{n}\right) \geq \operatorname{trc}\left(W_{n}\right)=3=\left\lceil\frac{n}{3}\right\rceil+1$ for $4 \leq n \leq 6$. It remains to show that $\operatorname{strc}\left(W_{n}\right) \geq\left\lceil\frac{n}{3}\right\rceil+1$ for $n \geq 7$. Assume the contrary, and suppose that there is a strongly total rainbow connected colouring $c$ of $W_{n}$, using at most $\left\lceil\frac{n}{3}\right\rceil$ colours. Since $n \geq 7$, for each vertex $v_{i}$, there exists at least one vertex $v_{j}$ with $j \neq i$ such that the unique $v_{i}-v_{j}$ geodesic of length 2 passes through the centre $v$. Thus, $c(v) \neq c\left(v v_{i}\right)$ for $0 \leq i \leq n-1$. Therefore, the $n$ edges $v v_{i}$ use at most $\left\lceil\frac{n}{3}\right\rceil-1<\frac{n}{3}$ different colours. One can deduce that there exist at least four different edges, say $v v_{i}, v v_{j}$, $v v_{k}, v v_{\ell}$, such that $c\left(v v_{i}\right)=c\left(v v_{j}\right)=c\left(v v_{k}\right)=c\left(v v_{\ell}\right)$. Again, since $n \geq 7$, we may assume that the unique $v_{i}-v_{j}$ geodesic is precisely the path $v_{i} v v_{j}$. So, there is no total-rainbow $v_{i}-v_{j}$ geodesic, a contradiction. Consequently, $\operatorname{strc}\left(W_{n}\right) \geq\left\lceil\frac{n}{3}\right\rceil+1$ for $n \geq 7$.

Our next aim is to consider complete bipartite graphs $K_{m, n}$. Clearly we have $\operatorname{rc}\left(K_{1, n}\right)=\operatorname{src}\left(K_{1, n}\right)=n ; \operatorname{rvc}\left(K_{1,1}\right)=\operatorname{srvc}\left(K_{1,1}\right)=0$ and $\operatorname{rvc}\left(K_{m, n}\right)=\operatorname{srvc}\left(K_{m, n}\right)=$ 1 for $(m, n) \neq(1,1)$; and $\operatorname{trc}\left(K_{1,1}\right)=\operatorname{strc}\left(K_{1,1}\right)=1$ and $\operatorname{trc}\left(K_{1, n}\right)=\operatorname{strc}\left(K_{1, n}\right)=$ $n+1$ for $n \geq 2$. For $2 \leq m \leq n$, the values of $\operatorname{rc}\left(K_{m, n}\right)$ and $\operatorname{src}\left(K_{m, n}\right)$ were determined by Chartrand et al. [6], and the values of $\operatorname{trc}\left(K_{m, n}\right)$ were determined by Liu et al. [26].

Theorem 3.6. [6, 26] Let $2 \leq m \leq n$. We have the following.
(a) $\operatorname{rc}\left(K_{m, n}\right)=\min (\lceil\sqrt[m]{n}\rceil, 4)$.
(b) $\operatorname{src}\left(K_{m, n}\right)=\lceil\sqrt[m]{n}\rceil$.
(c) $\operatorname{trc}\left(K_{m, n}\right)=\min (\lceil\sqrt[m]{n}\rceil+1,7)$.

In the next result, we will determine $\operatorname{strc}\left(K_{m, n}\right)$ for $2 \leq m \leq n$.
Theorem 3.7. For $2 \leq m \leq n$, we have $\operatorname{strc}\left(K_{m, n}\right)=\lceil\sqrt[m]{n}\rceil+1$.

Proof. Since $\operatorname{diam}\left(K_{m, n}\right)=2$, we have $\operatorname{strc}\left(K_{m, n}\right) \leq \operatorname{src}\left(K_{m, n}\right)+1=\lceil\sqrt[m]{n}\rceil+1$ by Proposition 2.1 and Theorem 3.6(b).

Now we prove the lower bound $\operatorname{strc}\left(K_{m, n}\right) \geq\lceil\sqrt[m]{n}\rceil+1$. This proof will be a slight modification of the proof of the lower bound of Theorem 3.6(c) in [26], but we provide it for the sake of clarity. Let the classes of $K_{m, n}$ be $U=\left\{u_{1}, \ldots, u_{m}\right\}$ and $V$, where $|V|=n$. Let $b=\lceil\sqrt[m]{n}\rceil \geq 2$. If $m \leq n \leq 2^{m}$, then $\operatorname{strc}\left(K_{m, n}\right) \geq 3=b+1$. Now let $n>2^{m}$, so that $b \geq 3$. We have $(b-1)^{m}<n \leq b^{m}$. Let $c$ be a totalcolouring of $K_{m, n}$, using colours from $\{1, \ldots, b\}$. For $v \in V$, assign $v$ with the vector $\vec{v}$ of length $m$, where $\vec{v}_{i}=c\left(u_{i} v\right)$ for $1 \leq i \leq m$. For two partitions $\mathcal{P}$ and $\mathcal{P}^{\prime}$ of $V$, we say that $\mathcal{P}$ refines $\mathcal{P}^{\prime}$, written $\mathcal{P}^{\prime} \prec \mathcal{P}$, if for all $A \in \mathcal{P}$, we have $A \subseteq B$ for some $B \in \mathcal{P}^{\prime}$. In other words, $\mathcal{P}$ can be obtained from $\mathcal{P}^{\prime}$ by partitioning some of the sets of $\mathcal{P}^{\prime}$. We define a sequence of refining partitions $\mathcal{P}_{0} \prec \mathcal{P}_{1} \prec \cdots \prec \mathcal{P}_{m}$ of $V$, with $\left|\mathcal{P}_{i}\right| \leq(b-1)^{i}$ for $0 \leq i \leq m$, as follows. Initially, set $\mathcal{P}_{0}=\{V\}$. Now, for $1 \leq i \leq m$, suppose that we have defined $\mathcal{P}_{i-1}$ with $\left|\mathcal{P}_{i-1}\right| \leq(b-1)^{i-1}$. Let $\mathcal{P}_{i-1}=\left\{A_{1}, \ldots, A_{\ell}\right\}$, where $\ell \leq(b-1)^{i-1}$. Define $\mathcal{P}_{i}$ as follows. For $1 \leq q \leq \ell$ and $A_{q} \in \mathcal{P}_{i-1}$, let

$$
\begin{aligned}
& B_{1}^{q}=\left\{v \in A_{q}: \vec{v}_{i}=c\left(u_{i}\right) \text { or } c\left(u_{i}\right)+1(\bmod b)\right\}, \\
& B_{r}^{q}=\left\{v \in A_{q}: \vec{v}_{i}=c\left(u_{i}\right)+r(\bmod b)\right\}, \text { for } 2 \leq r \leq b-1 .
\end{aligned}
$$

Let $\mathcal{P}_{i}=\left\{B_{r}^{q}: 1 \leq q \leq \ell, 1 \leq r \leq b-1\right.$ and $\left.B_{r}^{q} \neq \emptyset\right\}$, so that $\mathcal{P}_{i}$ is a partition of $V$ with $\left|\mathcal{P}_{i}\right| \leq(b-1)^{i}$ and $\mathcal{P}_{i-1} \prec \mathcal{P}_{i}$. Proceeding inductively, we obtain the partitions $\mathcal{P}_{0} \prec \mathcal{P}_{1} \prec \cdots \prec \mathcal{P}_{m}$ of $V$, with $\left|\mathcal{P}_{i}\right| \leq(b-1)^{i}$ for $0 \leq i \leq m$. Now, observe that for every $1 \leq i \leq m$, and any two vertices $y$ and $z$ in the same set in $\mathcal{P}_{i}$, the path $y u_{i} z$ is not total-rainbow, since $c\left(u_{i} y\right)=\vec{y}_{i}$ and $c\left(u_{i} z\right)=\vec{z}_{i}$ are either in $\left\{c\left(u_{i}\right), c\left(u_{i}\right)+1\right\}(\bmod b)$, or they are both $c\left(u_{i}\right)+r(\bmod b)$ for some $2 \leq r \leq b-1$. Since $n>(b-1)^{m} \geq\left|\mathcal{P}_{m}\right|$, there exists a set in $\mathcal{P}_{m}$ with at least two vertices $w$ and $x$, and since $\mathcal{P}_{1} \prec \cdots \prec \mathcal{P}_{m}$, this means that $w$ and $x$ are in the same set in $\mathcal{P}_{i}$ for every $1 \leq i \leq m$. Therefore, $w u_{i} x$ is not a total-rainbow path for every $1 \leq i \leq m$. Since the paths $w u_{i} x$ are all the possible $w-x$ geodesics (with length 2) in $K_{m, n}$, it follows that there does not exist a total-rainbow $w-x$ geodesic. Hence, $c$ is not a strongly total rainbow connected colouring of $K_{m, n}$, and $\operatorname{strc}\left(K_{m, n}\right) \geq b+1$.

To conclude this section, we consider complete multipartite graphs. Let $K_{n_{1}, \ldots, n_{t}}$ denote the complete multipartite graph with $t \geq 3$ classes, where $1 \leq n_{1} \leq \cdots \leq n_{t}$ are the class sizes. Clearly, we have $\operatorname{rvc}\left(K_{n_{1}, \ldots, n_{t}}\right)=\operatorname{srvc}\left(K_{n_{1}, \ldots, n_{t}}\right)=0$ (respectively 1) if $n_{t}=1$ (respectively $n_{t} \geq 2$ ). The values of $\operatorname{rc}\left(K_{n_{1}, \ldots, n_{t}}\right)$ and $\operatorname{src}\left(K_{n_{1}, \ldots, n_{t}}\right)$ were determined by Chartrand et al. [6], and the values of $\operatorname{trc}\left(K_{n_{1}, \ldots, n_{t}}\right)$ were determined by Liu et al. [26], as follows.
Theorem 3.8. $[6,26]$ Let $G=K_{n_{1}, \ldots, n_{t}}$, where $t \geq 3,1 \leq n_{1} \leq \cdots \leq n_{t}, m=$ $\sum_{i=1}^{t-1} n_{i}$ and $n_{t}=n$. Then, the values of $\operatorname{rc}(G), \operatorname{src}(G)$ and $\operatorname{trc}(G)$ are given in the following table.

|  | $n=1$ | $n \geq 2$ and $m>n$ | $m \leq n$ |
| :---: | :---: | :---: | :---: |
| $\operatorname{rc}(G)$ | 1 | 2 | $\min (\lceil\sqrt[m]{n}\rceil, 3)$ |
| $\operatorname{src}(G)$ | 1 | 2 | $\lceil\sqrt[m]{n}\rceil$ |
| $\operatorname{trc}(G)$ | 1 | 3 | $\min (\lceil\sqrt[m]{n}\rceil+1,5)$ |

Here, we determine the values of $\operatorname{strc}\left(K_{n_{1}, \ldots, n_{t}}\right)$ for $t \geq 3$.
Theorem 3.9. Let $t \geq 3,1 \leq n_{1} \leq \cdots \leq n_{t}, m=\sum_{i=1}^{t-1} n_{i}$ and $n_{t}=n$. Then,

$$
\operatorname{strc}\left(K_{n_{1}, \ldots, n_{t}}\right)= \begin{cases}1 & \text { if } n=1, \\ 3 & \text { if } n \geq 2 \text { and } m>n, \\ \lceil\sqrt[m]{n}\rceil+1 & \text { if } m \leq n .\end{cases}
$$

Proof. Write $G$ for $K_{n_{1}, \ldots, n_{t}}$, and let $V_{i}$ be the $i$ th class (with $n_{i}$ vertices) for $1 \leq i \leq t$. If $n=1$, then $G=K_{t}$ and $\operatorname{strc}(G)=1$. Now for $n \geq 2$, we have $\operatorname{strc}(G) \geq 3$. For the case $n \geq 2$ and $m>n$, we have $\operatorname{src}(G)=2$ by Theorem 3.8. Since $\operatorname{diam}(G)=2$, by Proposition 2.1, we have $\operatorname{strc}(G) \leq \operatorname{src}(G)+1=3$. Thus, $\operatorname{strc}(G)=3$.

Now, let $m \leq n$. For this case, we have $\operatorname{src}(G)=\lceil\sqrt[m]{n}\rceil$ by Theorem 3.8. Again by Proposition 2.1, we have the upper bound $\operatorname{strc}(G) \leq \operatorname{src}(G)+1=\lceil\sqrt[m]{n}\rceil+1$. It remains to prove the lower bound $\operatorname{strc}(G) \geq\lceil\sqrt[m]{n}\rceil+1$. Let $b=\lceil\sqrt[m]{n}\rceil \geq 2$. If $m \leq n \leq 2^{m}$, then $\operatorname{strc}(G) \geq 3=b+1$. Now let $n>2^{m}$, so that $b \geq 3$. We have $(b-1)^{m}<n \leq b^{m}$. Suppose that we have a total-colouring $c$ of $G$, using at most $b$ colours. Note that $K_{m, n}$ is a spanning subgraph of $G$ with classes $U=V_{1} \cup \cdots \cup V_{t-1}$ and $V_{t}$. We can restrict the total-colouring $c$ to $K_{m, n}$ and apply the same argument involving the refining partitions as in Theorem 3.7. We have vertices $w, x \in V_{t}$ such that all of the paths $w u x$, for $u \in U$, are not total-rainbow. Since these paths are all the possible $w-x$ geodesics in $G$ (of length 2), it follows that there does not exist a total-rainbow $w-x$ geodesic in $G$. Therefore, $c$ is not a strongly total rainbow connected colouring of $G$, and $\operatorname{strc}(G) \geq b+1$.

## 4 Comparing the rainbow connection numbers

Our aim in this section is to compare the various rainbow connection parameters. In [14], Krivelevich and Yuster observed that for $\operatorname{rc}(G)$ and $\operatorname{rvc}(G)$, we cannot generally find an upper bound for one of the parameters in terms of the other. Indeed, let $s \geq 2$. By taking $G=K_{1, s}$, we have $\operatorname{rc}(G)=s$ and $\operatorname{rvc}(G)=1$. On the other hand, let the graph $G_{s}$ be constructed as follows. Take $s$ vertex-disjoint triangles and, by designating a vertex from each triangle, add a complete graph $K_{s}$ on the designated vertices. Then $\operatorname{rc}\left(G_{s}\right) \leq 4$ and $\operatorname{rvc}\left(G_{s}\right)=s$.

We may consider the analogous situation for the parameters $\operatorname{src}(G)$ and $\operatorname{srvc}(G)$. Again by taking $G=K_{1, s}$, we see that $\operatorname{src}(G)=s$ and $\operatorname{srvc}(G)=1$, so that $\operatorname{src}(G)$ can be arbitrarily larger than $\operatorname{srvc}(G)$. Rather surprisingly, unlike the situation for the values of $\operatorname{rvc}(G)$ and $\operatorname{rc}(G)$, we are uncertain if $\operatorname{srvc}(G)$ can also be arbitrarily larger than $\operatorname{src}(G)$. We propose the following problem.
Problem 4.1. Does there exist an infinite family of connected graphs $\mathcal{F}$ such that, $\operatorname{src}(G)$ is bounded on $\mathcal{F}$, while $\operatorname{srvc}(G)$ is unbounded?

When considering the total rainbow connection number in addition, we have the following trivial inequalities.

$$
\begin{align*}
\operatorname{trc}(G) & \geq \max (\operatorname{rc}(G), \operatorname{rvc}(G))  \tag{5}\\
\operatorname{strc}(G) & \geq \max (\operatorname{src}(G), \operatorname{srvc}(G)) \tag{6}
\end{align*}
$$

In [26], Liu et al. considered how close and how far apart the terms in the inequality (5) can be. They observed that by considering Krivelevich and Yuster's construction as described above, we have $\operatorname{trc}\left(G_{s}\right)=\operatorname{rvc}\left(G_{s}\right)=s$ for $s \geq 13$. Also, as mentioned in the remark after the proof of Theorem 2.3, if $G=\left(C_{5}\right)_{u \rightarrow K}$ is a graph obtained by expanding a vertex $u$ of the cycle $C_{5}$ into a clique $K$, then we have $\operatorname{trc}(G)=\operatorname{rc}(G)=3$. Thus, $\operatorname{trc}(G)$ can be equal to each of $\operatorname{rvc}(G)$ and $\operatorname{rc}(G)$ for infinitely many graphs $G$. On the other hand, Liu et al. also remarked that, given $1 \leq t<s$, there exists a graph $G$ such that $\operatorname{trc}(G) \geq s$ and $\operatorname{rvc}(G)=t$. Indeed, we can let $G=B_{s, t}$ be the graph obtained by taking the star $K_{1, s}$ and identifying the centre with one end-vertex of the path of length $t$ (this graph $B_{s, t}$ is a broom). Also, for $s \geq 13$, we can again consider the graphs $G_{s}$ and obtain $\operatorname{trc}\left(G_{s}\right)=s$ and $\operatorname{rc}\left(G_{s}\right) \leq 4$. Thus, $\operatorname{trc}(G)$ can also be arbitrarily larger than each of $\operatorname{rvc}(G)$ and $\operatorname{rc}(G)$. For the difference between the terms $\operatorname{trc}(G)$ and $\max (\operatorname{rc}(G), \operatorname{rvc}(G))$, one can consider $G$ to be the path of length $s$, and obtain $\operatorname{trc}(G)=2 s-1$ and $\max (\operatorname{rc}(G), \operatorname{rvc}(G))=s$, so that $\operatorname{trc}(G)-\max (\operatorname{rc}(G), \operatorname{rvc}(G))=s-1$ can be arbitrarily large. However, for this simple example, the term $\max (\operatorname{rc}(G), \operatorname{rvc}(G))$ is unbounded in $s$. In the final problem in [26], Liu et al. asked the question of whether there exists an infinite family of connected graphs $\mathcal{F}$ such that, $\max (\operatorname{rc}(G), \operatorname{rvc}(G))$ is bounded on $\mathcal{F}$, while $\operatorname{trc}(G)$ is unbounded. This open problem appears to be much more challenging.

Here, we consider the analogous situations for the terms in the inequality (6). From the previous remarks and results, we can easily obtain the following.

## Theorem 4.2.

(a) There exist infinitely many graphs $G$ with $\operatorname{strc}(G)=\operatorname{src}(G)=3$.
(b) Given $s \geq 13$, there exists a graph $G$ with $\operatorname{strc}(G)=\operatorname{srvc}(G)=s$.
(c) Given $1 \leq t<s$, there exists a graph $G$ such that $\operatorname{strc}(G) \geq s$ and $\operatorname{srvc}(G)=t$.

Proof. (a) Let $G=\left(C_{5}\right)_{u \rightarrow K}$ as described earlier. We have $\operatorname{trc}(G)=\operatorname{rc}(G)=3$. By Theorem 2.3(c), we have $\operatorname{strc}(G)=3$. Therefore by (1) and (6), we have $3=$ $\operatorname{strc}(G) \geq \operatorname{src}(G) \geq \operatorname{rc}(G)=3$, so that $\operatorname{strc}(G)=\operatorname{src}(G)=3$.
(b) We use the following construction which was given by Lei et al. [16]. For $s \geq 13$, let $H_{s}$ be the graph as follows. First, we take the graph $G_{s}$ from before, where $u_{0}, \ldots, u_{s-1}$ are the vertices of the $K_{s}$, and the remaining vertices are $v_{i}, w_{i}$, where $u_{i} v_{i} w_{i}$ is a triangle, for $0 \leq i \leq s-1$. We then add new vertices $z_{0}, \ldots, z_{s-1}$, and connect the edges $u_{i} z_{i}, u_{i+1} z_{i}, v_{i} z_{i}, w_{i} z_{i+4}$, for $0 \leq i \leq s-1$, where all indices are taken modulo $s$. In [16], Lei et al. proved that $\operatorname{strc}\left(H_{s}\right)=\operatorname{srvc}\left(H_{s}\right)=s$.
(c) Since the broom $G=B_{s, t}$ as described earlier is a tree, it is clear that $\operatorname{strc}(G)=$ $\operatorname{trc}(G) \geq s$ and $\operatorname{srvc}(G)=\operatorname{rvc}(G)=t$.

As before, if $G$ is the path of length $s$, then we have $\operatorname{strc}(G)-\operatorname{src}(G)=\operatorname{strc}(G)-$ $\max (\operatorname{src}(G), \operatorname{srvc}(G))=s-1$, so that the two differences can both be arbitrarily large. But the terms $\operatorname{src}(G)$ and $\max (\operatorname{src}(G), \operatorname{srvc}(G))$ are unbounded in $s$. Similar to the question of Liu et al. in [26] and Problem 4.1, we may ask the following question.

Problem 4.3. Does there exist an infinite family of connected graphs $\mathcal{F}$ such that, $\operatorname{src}(G)$ is bounded on $\mathcal{F}$, while $\operatorname{strc}(G)$ is unbounded? Similarly, does there exist an infinite family of connected graphs $\mathcal{F}$ such that, $\max (\operatorname{src}(G), \operatorname{srvc}(G))$ is bounded on $\mathcal{F}$, while $\operatorname{strc}(G)$ is unbounded?

Now, we proceed to the final part of this section. Recall that the following inequalities hold for a connected graph $G$.

$$
\operatorname{rc}(G) \leq \operatorname{src}(G), \quad \operatorname{rvc}(G) \leq \operatorname{srvc}(G), \quad \text { and } \quad \operatorname{trc}(G) \leq \operatorname{strc}(G)
$$

Chartrand et al. [6] considered the following question: Given positive integers $a \leq b$, does there exist a graph $G$ such that $\operatorname{rc}(G)=a$ and $\operatorname{src}(G)=b$ ? They gave positive answers for $a=b$, and $3 \leq a<b$ with $b \geq \frac{5 a-6}{3}$. Chern and Li [9] then improved this result as follows.

Theorem 4.4. [9] Let $a$ and $b$ be positive integers. Then there exists a connected graph $G$ such that $\operatorname{rc}(G)=a$ and $\operatorname{src}(G)=b$ if and only if $a=b \in\{1,2\}$ or $3 \leq a \leq b$.

Theorem 4.4 was an open problem of Chartrand et al., and it completely characterises all possible pairs $a$ and $b$ for the above question. Subsequently, Li et al. [20] studied the rainbow vertex-connection analogue, and they proved the following result.

Theorem 4.5. [20] Let $a$ and $b$ be integers with $a \geq 5$ and $b \geq \frac{7 a-8}{5}$. Then there exists a connected graph $G$ such that $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$.

Here, we will improve Theorem 4.5, and also study the total rainbow connection version of the problem. We will prove Theorems 4.6 and 4.7 below, where we will completely characterise all pairs of positive integers $a$ and $b$ such that, there exists a $\operatorname{graph} G$ with $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b(\operatorname{respectively} \operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b)$.

Theorem 4.6. Let $a$ and $b$ be positive integers. Then there exists a connected graph $G$ such that $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$ if and only if $a=b \in\{1,2\}$ or $3 \leq a \leq b$.

Theorem 4.7. Let $a$ and $b$ be positive integers. Then there exists a connected graph $G$ such that $\operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b$ if and only if $a=b \in\{1,3,4\}$ or $5 \leq a \leq b$.

To prove Theorems 4.6 and 4.7, we first prove three auxiliary lemmas.
Lemma 4.8. For every $b \geq 3$, there exists a graph $G$ such that $\operatorname{rvc}(G)=3$ and $\operatorname{srvc}(G)=b$.

Proof. We construct a graph $F_{b}$ as follows. We take a complete graph $K_{2 b}$, say with vertices $u_{0}, \ldots, u_{2 b-1}$, and further vertices $v_{0}, \ldots, v_{2 b-1}, w_{0}, \ldots, w_{2 b-1}$. For $0 \leq i \leq$ $2 b-1$, we add the edges $u_{i} v_{i}, u_{i} v_{i-1}, u_{i} w_{i}, w_{i} v_{i}, w_{i} v_{i-1}$. Throughout, the indices of the vertices $u_{i}, v_{i}, w_{i}$ are taken modulo $2 b$. We show that $\operatorname{rvc}\left(F_{b}\right)=3$ and $\operatorname{srvc}\left(F_{b}\right)=b$.

Suppose firstly that we have a vertex-colouring of $F_{b}$, using at most two colours. Since $2 b \geq 6$, we may assume that $u_{0}$ and $u_{\ell}$ have the same colour, for some $2 \leq$
$\ell \leq 2 b-2$. Then note that $w_{0} u_{0} u_{\ell} w_{\ell}$ is the unique $w_{0}-w_{\ell}$ geodesic, with length 3 . Thus, there does not exist a vertex-rainbow $w_{0}-w_{\ell}$ path, and we have $\operatorname{rvc}\left(F_{b}\right) \geq 3$. Now, we define a vertex-colouring $f$ of $F_{b}$ as follows. Let $f\left(u_{i}\right)=1$ if $i$ is odd, and $f\left(u_{i}\right)=2$ if $i$ is even. Let $f(z)=3$ for all other vertices $z$. It is easy to check that $f$ is a rainbow vertex-connected colouring for $F_{b}$. For example, to connect $w_{0}$ to $w_{i}$ with a vertex-rainbow path, where $2 \leq i \leq 2 b-2$, we may take $w_{0} u_{0} u_{i} w_{i}$ if $i$ is odd, and $w_{0} u_{0} u_{i-1} v_{i-1} w_{i}$ if $i$ is even. Thus $\operatorname{rvc}\left(F_{b}\right) \leq 3$, and we have $\operatorname{rvc}\left(F_{b}\right)=3$.

Next, suppose that we have a vertex-colouring of $F_{b}$, using fewer than $b$ colours. Then, three of the vertices $u_{i}$ have the same colour, so we may assume that $u_{0}$ and $u_{\ell}$ have the same colour, for some $2 \leq \ell \leq 2 b-2$. Note that $w_{0} u_{0} u_{\ell} w_{\ell}$ is the unique $w_{0}-w_{\ell}$ geodesic (with length 3). Thus, there does not exist a vertex-rainbow $w_{0}-w_{\ell}$ geodesic, and we have $\operatorname{srvc}\left(F_{b}\right) \geq b$. Now, we define a vertex-colouring $g$ of $F_{b}$ as follows. Let $g\left(u_{i}\right)=\left\lceil\frac{i+1}{2}\right\rceil$ for $0 \leq i \leq 2 b-1$, and $g(z)=1$ for all other vertices $z$. We show that $g$ is a strongly rainbow vertex-connected colouring for $F_{b}$. It is easy to see that each vertex $u_{i}$ is at distance at most 2 from every other vertex. Also, any two vertices of $\left\{v_{0}, \ldots, v_{2 b-1}, w_{0}, \ldots, w_{2 b-1}\right\}$ are either at distance at most 2 apart, or they are at distance 3 apart, and there is a vertex-rainbow geodesic connecting them. Indeed, if $d\left(v_{i}, v_{j}\right)=3$, then we can take the $v_{i}-v_{j}$ geodesic $v_{i} u_{i} u_{j} v_{j}$. Similarly, if $d\left(v_{i}, w_{j}\right)=3$, then we take $v_{i} u_{i} u_{j} w_{j}$ or $v_{i} u_{i+1} u_{j} w_{j}$. If $d\left(w_{i}, w_{j}\right)=3$, then we take $w_{i} u_{i} u_{j} w_{j}$. Thus $\operatorname{srvc}\left(F_{b}\right) \leq b$, and we have $\operatorname{srvc}\left(F_{b}\right)=b$.

Lemma 4.9. For every $a$ and $b, 4 \leq a \leq b$, there exists a connected graph $G$ such that $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$.

Proof. We construct a graph $F_{a, b}$ as follows. Let $n=2(b-1)(b-a+2) \geq 12$. We take a set of vertices $V=\left\{v_{1}, \ldots, v_{n}\right\}$, and another vertex $u$ and a path $u_{0} \cdots u_{a-3}$. We add the paths $u w_{i} v_{i}$ and $u_{a-3} x_{i} v_{i}$ for $1 \leq i \leq n$, and then the edges $v_{\ell} v_{\ell+1}, w_{\ell} w_{\ell+1}, x_{\ell} x_{\ell+1}$ for $1 \leq \ell<n$ with $\ell$ odd. Let $U=\left\{u_{0}, \ldots, u_{a-3}\right\}$, $W=\left\{w_{1}, \ldots, w_{n}\right\}$ and $X=\left\{x_{1}, \ldots, x_{n}\right\}$. Note that we have perfect matchings within the sets $V, W$ and $X$. We show that $\operatorname{rvc}\left(F_{a, b}\right)=a$ and $\operatorname{srvc}\left(F_{a, b}\right)=b$.

Clearly we have $\operatorname{rvc}\left(F_{a, b}\right) \geq \operatorname{diam}\left(F_{a, b}\right)-1=a$. Now, we define a vertex-colouring $c$ of $F_{a, b}$ as follows. Let $c\left(u_{j}\right)=j$ for $1 \leq j \leq a-3$. For $1 \leq i \leq n$, let $c\left(w_{i+1}\right)=$ $c\left(x_{i}\right)=a-2$ if $i$ is odd, and $c\left(w_{i-1}\right)=c\left(x_{i}\right)=a-1$ if $i$ is even. Let $c(z)=a$ for all other vertices $z$. It is easy to check that $c$ is a rainbow vertex-connected colouring for $F_{a, b}$. For example, for $i \neq 2$, to connect $v_{1}$ to $v_{i}$ with a vertex-rainbow path, we may take $v_{1} x_{1} u_{a-3} x_{i} v_{i}$ if $i$ is even, and $v_{1} x_{1} u_{a-3} x_{i+1} v_{i+1} v_{i}$ if $i$ is odd, since $a \geq 4$. Thus $\operatorname{rvc}\left(F_{a, b}\right) \leq a$, and we have $\operatorname{rvc}\left(F_{a, b}\right)=a$.

Next, suppose that there exists a strongly rainbow vertex-connected colouring $f$ of $F_{a, b}$, using at most $b-1$ colours, say colours $1,2, \ldots, b-1$. Then note that for every $1 \leq i \leq n$, the unique $u_{0}-v_{i}$ geodesic is $u_{0} u_{1} \cdots u_{a-3} x_{i} v_{i}$. Thus we may assume that $f\left(u_{j}\right)=j$ for $1 \leq j \leq a-3$, so that $f\left(x_{i}\right) \in\{a-2, a-1, \ldots, b-1\}$ for $1 \leq i \leq n$. Also, we have $f\left(w_{i}\right), f(u) \in\{1, \ldots, b-1\}$ for $1 \leq i \leq n$. For $a-2 \leq p \leq b-1$ and $1 \leq q \leq b-2$, we define the set $A_{p, q} \subset V$ where

$$
\begin{aligned}
& A_{p, 1}=\left\{v_{i} \in V: f\left(x_{i}\right)=p \text { and } f\left(w_{i}\right)=f(u) \text { or } f(u)+1(\bmod b-1)\right\}, \\
& A_{p, q}=\left\{v_{i} \in V: f\left(x_{i}\right)=p \text { and } f\left(w_{i}\right)=f(u)+q(\bmod b-1)\right\}, \text { for } q \geq 2 .
\end{aligned}
$$

Note that $V=\bigcup_{p, q}\left\{A_{p, q}: A_{p, q} \neq \emptyset\right\}$ is a partition of $V$ with at most $(b-2)(b-a+2)$ parts. Since $n=2(b-1)(b-a+2)$, there exists a set $A_{r, s}$ with at least three vertices. Thus, we may assume that $v_{1}, v_{\ell} \in A_{r, s}$ with $\ell \neq 2$. Observe that the path $v_{1} x_{1} u_{a-3} x_{\ell} v_{\ell}$ is not vertex-rainbow, since $f\left(x_{1}\right)=f\left(x_{\ell}\right)=r$. Also, the path $v_{1} w_{1} u w_{\ell} v_{\ell}$ is not vertex-rainbow, since $f\left(w_{1}\right)$ and $f\left(w_{\ell}\right)$ are either in $\{f(u), f(u)+1\}$ $(\bmod b-1)$, or they are both $f(u)+s(\bmod b-1)$. Since these two paths are the only $v_{1}-v_{\ell}$ geodesics (with length 4), we have a contradiction. Thus, $\operatorname{srvc}\left(F_{a, b}\right) \geq b$.

Finally, we define a vertex-colouring $g$ of $F_{a, b}$, using colours $1,2, \ldots, b$, as follows. Let $g\left(u_{j}\right)=j$ for $1 \leq j \leq a-3$, and $g(u)=g\left(u_{0}\right)=g\left(v_{i}\right)=b$ for $1 \leq i \leq n$. Now, note that there are $(b-1)(b-a+2)$ pairs $\left\{v_{\ell}, v_{\ell+1}\right\}$ with $\ell$ odd, and also $(b-1)(b-a+2)$ distinct vectors of length 2 , whose first coordinate is in $\{a-2, \ldots, b-1\}$ and second coordinate is in $\{1, \ldots, b-1\}$. Thus we may assign these distinct vectors to all vertices of $V$ such that, both vertices of a pair $\left\{v_{\ell}, v_{\ell+1}\right\}$ with $\ell$ odd receive the same vector (so that every vector appears exactly twice). If $v_{\ell}$ and $v_{\ell+1}$ have been assigned the vector $\vec{v}$, then we set $g\left(x_{\ell}\right)=g\left(x_{\ell+1}\right)=\vec{v}_{1} \in\{a-2, \ldots, b-1\}$, and $g\left(w_{\ell}\right)=g\left(w_{\ell+1}\right)=\vec{v}_{2} \in\{1, \ldots, b-1\}$. We show that $g$ is a strongly rainbow vertexconnected colouring for $F_{a, b}$. We must show that for every $x, y \in V\left(F_{a, b}\right)$, there is a vertex-rainbow $x-y$ geodesic.

- If $x \in U$ and $y \neq u$, then it is easy to find a vertex-rainbow $x-y$ geodesic. For example, if $x=u_{j}$ and $y=w_{i}$, then we take $u_{j} \cdots u_{a-3} x_{i} v_{i} w_{i}$. If $x=u_{j}$ and $y=u$, then we take $u_{j} \cdots u_{a-3} x_{\ell} v_{\ell} w_{\ell} u$, where $v_{\ell}$ is assigned the vector $(a-2, b-1)$. Similarly, it is easy to deal with the case when $x=u$ and $y \in V \cup W \cup X$.
- Now we consider the case $x, y \in V \cup W \cup X$. Firstly, the cases $x, y \in W$ and $x, y \in X$ are clear, since $d(x, y) \leq 2$. Next, suppose that $x \in V$, say $x=v_{1}$. Then the case $y \in\left\{w_{1}, x_{1}, v_{2}, w_{2}, x_{2}\right\}$ is clear, since we have $d(x, y) \leq 2$. If $y=w_{\ell}$ (respectively $x_{\ell}$ ) for some $\ell \neq 2$, then we take $v_{1} w_{1} u w_{\ell}$ (respectively $v_{1} x_{1} u_{a-3} x_{\ell}$ ). If $y=v_{\ell}$ for some $\ell \neq 2$, then $x$ and $y$ are assigned different vectors, say $\vec{x} \neq \vec{y}$. If $\vec{x}_{1} \neq \vec{y}_{1}$, then we take $v_{1} x_{1} u_{a-3} x_{\ell} v_{\ell}$, and if $\vec{x}_{2} \neq \vec{y}_{2}$, then we take $v_{1} w_{1} u w_{\ell} v_{\ell}$. Finally, it remains to consider the case $x \in W$ and $y \in X$. We may assume that $x=w_{1}$ and $y=x_{\ell}$ for some $1 \leq \ell \leq n$. We take $w_{1} v_{1} x_{1}$ if $\ell=1 ; w_{1} v_{1} x_{1} x_{2}$ if $\ell=2$; and $w_{1} v_{1} x_{1} u_{a-3} x_{\ell}$ if $\ell \geq 3$.

We always have a vertex-rainbow $x-y$ geodesic, so that $g$ is a strongly rainbow vertex-connected colouring. Therefore $\operatorname{srvc}\left(F_{a, b}\right) \leq b$, and we have $\operatorname{srvc}\left(F_{a, b}\right)=b$.
Lemma 4.10. For every $a$ and $b, 5 \leq a<b$, there exists a connected graph $G$ such that $\operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b$.

Proof. We consider the complete multipartite graph $K_{1, \ldots, 1, n}$, where there are $m \geq 2$ singleton classes, say $\left\{u_{1}\right\}, \ldots,\left\{u_{m}\right\}$. Let $U=\left\{u_{1}, \ldots, u_{m}\right\}$, and $V$ be the class with $n$ vertices. Given $5 \leq a<b$, let $G_{a, b, m}$ be the graph constructed as follows. We take $K_{1, \ldots, 1, n}$, and set $n=(b-2)^{m}+1$. We then add $a-1 \geq 4$ pendent edges at $u_{1}$, say $W=\left\{w_{1}, \ldots, w_{a-1}\right\}$ is the set of pendent vertices. We claim that for sufficiently large $m$, we have $\operatorname{trc}\left(G_{a, b, m}\right)=a$ and $\operatorname{strc}\left(G_{a, b, m}\right)=b$.

Since the bridges of $G_{a, b, m}$ are the $a-1$ pendent edges, and the only cut-vertex is $u_{1}$, clearly we have $\operatorname{trc}\left(G_{a, b, m}\right) \geq a$ by Proposition 2.2. Now we define a totalcolouring $f$ of $G_{a, b, m}$ as follows. Let $f\left(u_{1} w_{\ell}\right)=\ell$ for $1 \leq \ell \leq a-1$. For every $v \in V$, let $f\left(u_{1} v\right)=1$, and $f\left(u_{i} v\right)=2$ for all $2 \leq i \leq m$. Let $f\left(u_{i} u_{j}\right)=4$ for all $1 \leq i<j \leq m$. Let $f\left(u_{1}\right)=a$, and $f(z)=3$ for all $z \in V\left(G_{a, b, m}\right) \backslash\left\{u_{1}\right\}$. We claim that $f$ is a total rainbow connected colouring for $G_{a, b, m}$. We need to show that for every $x, y \in V\left(G_{a, b, m}\right)$, there is a total-rainbow $x-y$ path. Since $u_{1}$ is connected to all other vertices, it suffices to consider $x, y \in V\left(G_{a, b, m}\right) \backslash\left\{u_{1}\right\}$. If $x, y \notin W$ and $x, y$ are not adjacent, then $x, y \in V$, in which case we take the path $x u_{1} u_{2} y$. Now suppose $x \in W$. Then we can take the path $x u_{1} y$, unless if $x=w_{1}$ and $y \in V$, in which case we take $x u_{1} u_{2} y$; or $x=w_{4}$ and $y \in U \backslash\left\{u_{1}\right\}$, in which case we take $x u_{1} v y$ for some $v \in V$. Thus $f$ is a total rainbow connected colouring for $G_{a, b, m}$, and $\operatorname{trc}\left(G_{a, b, m}\right) \leq a$. We have $\operatorname{trc}\left(G_{a, b, m}\right)=a$.

Now, suppose that we have a total-colouring of $G_{a, b, m}$, using fewer than $b$ colours. Note that $\lceil\sqrt[m]{n}\rceil+1=b$, so that by Theorem 3.9, for the copy of $K_{1, \ldots, 1, n}$, we have $\operatorname{strc}\left(K_{1, \ldots, 1, n}\right)=b$. It follows that when restricted to the $K_{1, \ldots, 1, n}$, there are two vertices $w, x$ that are not connected by a total-rainbow $w-x$ geodesic. This means that we have $w, x \in V$, and the paths $x u w$, for $u \in U$, are all not total-rainbow. Since these paths are also all the possible $w-x$ geodesics in $G_{a, b, m}$, we do not have a total-rainbow $w-x$ geodesic in $G_{a, b, m}$. Thus $\operatorname{strc}\left(G_{a, b, m}\right) \geq b$.

It remains to prove that $\operatorname{strc}\left(G_{a, b, m}\right) \leq b$. Let $m$ be sufficiently large so that $(b-1)^{m-1}>(b-2)^{m}$. This inequality holds if $m>\frac{\log (b-1)}{\log (b-1)-\log (b-2)}$. Thus, we have $(b-1)^{m-1} \geq n$. We define a total-colouring $g$ of $G_{a, b, m}$ as follows. Let $g\left(u_{1} w_{\ell}\right)=\ell$ for $1 \leq \ell \leq a-1$. Let $g\left(u_{1}\right)=a$, and $g\left(u_{1} v\right)=g\left(u_{i} u_{j}\right)=g(z)=b$ for all $v \in V, 1 \leq i<j \leq m$, and $z \in V\left(G_{a, b, m}\right) \backslash\left\{u_{1}\right\}$. Now since $(b-1)^{m-1} \geq n$, we may assign distinct vectors of length $m-1$ to the vertices of $V$, with entries from $\{1,2, \ldots, b-1\}$. Suppose that $v \in V$ has been assigned the vector $\vec{v}$. We let $g\left(u_{i+1} v\right)=\vec{v}_{i}$ for $1 \leq i \leq m-1$ and $v \in V$. We claim that $g$ is a strongly total rainbow connected colouring for $G_{a, b, m}$. Similar to before, it suffices to show that for all $x, y \in V\left(G_{a, b, m}\right) \backslash\left\{u_{1}\right\}$, there is a total-rainbow $x-y$ geodesic. If $x, y \notin W$ and $x, y$ are not adjacent, then $x, y \in V$. We have $\vec{x}_{i} \neq \vec{y}_{i}$ for some $1 \leq i \leq m-1$, so that we can take the geodesic $x u_{i+1} y$. If $x \in W$, then we can take the geodesic $x u_{1} y$. Thus $g$ is a strongly total rainbow connected colouring for $G_{a, b, m}$, and $\operatorname{strc}\left(G_{a, b, m}\right) \leq b$. We have $\operatorname{strc}\left(G_{a, b, m}\right)=b$.

We can now prove Theorems 4.6 and 4.7.
Proof of Theorem 4.6. Suppose that there exists a connected graph $G$ such that $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$. Then obviously we have $a \leq b$. If $a=1$ (respectively $a=2$ ), then Theorem 2.3(c)(ii) (respectively (c)(iii)) gives $b=1$ (respectively $b=2$ ). Therefore, we have either $a=b \in\{1,2\}$, or $3 \leq a \leq b$.

Conversely, given $a, b$ such that either $a=b \in\{1,2\}$ or $3 \leq a \leq b$, we show that there exists a connected graph $G$ with $\operatorname{rvc}(G)=a$ and $\operatorname{srvc}(G)=b$. Obviously if $a=b \geq 1$, then $\operatorname{rvc}(G)=\operatorname{srvc}(G)=a$ if $G$ is the path of length $a+1$. The
remaining cases satisfy $3 \leq a \leq b$, and these are covered by Lemmas 4.8 and 4.9. Thus Theorem 4.6 follows.

Proof of Theorem 4.7. Suppose that there exists a connected graph $G$ such that $\operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b$. Then obviously we have $a \leq b$. If $a=1$ (respectively $a=$ $3, a=4$ ), then Theorem 2.3(a) (respectively (c)(iv), (c)(v)) gives $b=1$ (respectively $b=3, b=4$ ). Theorem 2.3(a) and (b) also imply that $a, b \neq 2$. Therefore, we have either $a=b \in\{1,3,4\}$, or $5 \leq a \leq b$.

Conversely, given $a, b$ such that either $a=b \in\{1,3,4\}$ or $5 \leq a \leq b$, we show that there is a connected $\operatorname{graph} G$ with $\operatorname{trc}(G)=a$ and $\operatorname{strc}(G)=b$. Obviously, if $a=b=1$, then $\operatorname{trc}(G)=\operatorname{strc}(G)=1$ if $G$ is any non-trivial complete graph, and if $a=b \geq 3$, then $\operatorname{trc}(G)=\operatorname{strc}(G)=a$ if $G$ is the star of order $a$. The remaining cases satisfy $5 \leq a<b$, and these are covered by Lemma 4.10. Thus Theorem 4.7 follows.

## Acknowledgements

Lin Chen, Xueliang Li and Jinfeng Liu were supported by the National Science Foundation of China (Nos. 11371205 and 11531011). Henry Liu was partially supported by the Startup Fund of One Hundred Talent Program of SYSU. Henry Liu would also like to thank the Chern Institute of Mathematics, Nankai University, for their generous hospitality. He was able to carry out part of this research during his visit there.

The authors thank the anonymous referees for their careful reading of the manuscript.

## References

[1] J. Alva-Samos and J.J. Montellano-Ballesteros, Rainbow connection in some digraphs, Graphs Combin. 32 (2016), 2199-2209.
[2] J. Alva-Samos and J.J. Montellano-Ballesteros, Rainbow connectivity of cacti and of some infinite digraphs, Discuss. Math. Graph Theory 37 (2017), 301-313.
[3] B. Bollobás, Modern Graph Theory, Springer-Verlag, New York, 1998, xiv+394 pp.
[4] Y. Caro, A. Lev, Y. Roditty, Zs. Tuza and R. Yuster, On rainbow connection, Electron. J. Combin. 15 (2008), R57.
[5] S. Chakraborty, E. Fischer, A. Matsliah and R. Yuster, Hardness and algorithms for rainbow connectivity, 26th Int. Symposium on Theoret. Aspects of Comp. Sci., STACS 2009 (2009), 243-254. See also J. Comb. Optim. 21 (2011), 330-347.
[6] G. Chartrand, G.L. Johns, K.A. McKeon and P. Zhang, Rainbow connection in graphs, Math. Bohem. 133 (2008), 85-98.
[7] L. Chen, X. Li and H. Lian, Further hardness results on the rainbow vertex connection number of graphs, Theoret. Comput. Sci. 481 (2013), 18-23.
[8] L. Chen, X. Li and Y. Shi, The complexity of determining the rainbow vertexconnection of a graph, Theoret. Comput. Sci. 412 (2011), 4531-4535.
[9] X. Chern and X. Li, A solution to a conjecture on the rainbow connection number, Ars Combin. 104 (2012), 193-196.
[10] P. Dorbec, I. Schiermeyer, E. Sidorowicz and É. Sopena, Rainbow connection in oriented graphs, Discrete Appl. Math. 179 (2014), 69-78.
[11] R. Holliday, C. Magnant and P. Salehi Nowbandegani, Note on rainbow connection in oriented graphs with diameter 2, Theory Appl. Graphs 1(1) (2014), art. 2 .
[12] X. Huang, X. Li and Y. Shi, Note on the hardness of rainbow connections for planar and line graphs, Bull. Malays. Math. Sci. Soc. 38 (2015), 1235-1241.
[13] H. Jiang, X. Li and Y. Zhang, Upper bounds for the total rainbow connection of graphs, J. Comb. Optim. 32 (2016), 260-266.
[14] M. Krivelevich and R. Yuster, The rainbow connection of a graph is (at most) reciprocal to its minimum degree, J. Graph Theory 63 (2010), 185-191.
[15] H. Lei, S. Li, H. Liu and Y. Shi, Rainbow vertex connection of digraphs, J. Comb. Optim. (2017), https://doi.org/10.1007/s10878-017-0156-7.
[16] H. Lei, H. Liu, C. Magnant and Y. Shi, Total rainbow connection of digraphs, Discrete Appl. Math., (to appear). ArXiv preprint arXiv:1701.04283.
[17] H. Li, X. Li and S. Liu, The (strong) rainbow connection numbers of Cayley graphs on Abelian groups, Comput. Math. Appl. 62 (2011), 4082-4088.
[18] S. Li, X. Li and Y. Shi, Note on the complexity of deciding the rainbow (vertex-) connectedness for bipartite graphs, Appl. Math. Comput. 258 (2015), 155-161.
[19] X. Li and S. Liu, Tight upper bound of the rainbow vertex-connection number for 2-connected graphs, Discrete Appl. Math. 173 (2014), 62-69.
[20] X. Li, Y. Mao and Y. Shi, The strong rainbow vertex-connection of graphs, Util. Math. 93 (2014), 213-223.
[21] X. Li and Y. Shi, On the rainbow vertex-connection, Discuss. Math. Graph Theory 33 (2013), 307-313.
[22] X. Li, Y. Shi and Y. Sun, Rainbow connections of graphs: A survey, Graphs Combin. 29 (2013), 1-38.
[23] X. Li and Y. Sun, Rainbow Connections of Graphs, Springer Briefs in Math., Springer, New York, 2012, viii+103 pp.
[24] X. Li and Y. Sun, On the strong rainbow connection of a graph, Bull. Malays. Math. Sci. Soc. 36 (2013), 299-311.
[25] H. Liu, A. Mestre and T. Sousa, Rainbow vertex $k$-connection in graphs, Discrete Appl. Math. 161 (2013), 2549-2555.
[26] H. Liu, Â. Mestre and T. Sousa, Total rainbow $k$-connection in graphs, Discrete Appl. Math. 174 (2014), 92-101.
[27] Y. Ma, Total rainbow connection number and complementary graph, Results Math. 70 (2016), 173-182.
[28] Y. Sun, On rainbow total-coloring of a graph, Discrete Appl. Math. 194 (2015), 171-177.
[29] Y. Sun, Z. Jin and F. Li, On total rainbow $k$-connected graphs, Appl. Math. Comput. 311 (2017), 223-227.
(Received 18 Apr 2017; revised 5 Oct 2017)


[^0]:    * Also at School of Mathematics and Statistics, Qinghai Normal University, Xining, Qinghai 810008, China
    $\dagger$ Corresponding author

