Two invariants for adjointly equivalent graphs

F.M. Dong

Math and Math E, NIE Nanyang Technological University Singapore

K.L. Teo, C.H.C. Little and M.D. Hendy

Institute of Fundamental Sciences Massey University Palmerston North New Zealand

Abstract

Two graphs are defined to be adjointly equivalent if their complements are chromatically equivalent. We study the properties of two invariants under adjoint equivalence.

1 Introduction

In this paper, all graphs considered are simple graphs. For a graph G, let \overline{G} , V(G), E(G), v(G), e(G), t(G), c(G) and $P(G, \lambda)$, respectively, be the complement, vertex set, edge set, order, size, number of triangles, number of components and chromatic polynomial of G.

A partition $\{A_1, A_2, \dots, A_k\}$ of V(G), where k is a positive integer, is called a *k*-independent partition of a graph G if each A_i is a nonempty independent set of G. Let $\alpha(G, k)$ denote the number of k-independent partitions of G. Then

$$P(G,\lambda) = \sum_{k=1}^{v(G)} \alpha(G,k)(\lambda)_k, \tag{1}$$

where $(\lambda)_k = \lambda(\lambda - 1) \cdots (\lambda - k + 1)$. (See [13].)

Two graphs G and H are said to be *chromatically equivalent* if they have the same chromatic polynomial. In this case we write $G \sim H$. The equivalence class determined by a graph G is denoted by [G]. A graph G is said to be *chromatically unique* if $[G] = \{G\}$.

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The determination of [G] for a given graph G has received much attention in the literature (see [4, 5]). The adjoint polynomial of a graph is a useful tool for this study. We now proceed to define it.

Let G be a graph with order n. If H is a spanning subgraph of G and each component of H is complete, then H is called a *clique cover* [2] (or, by Liu [6], an *ideal subgraph*) of G. Two clique covers are considered to be different if they have different edge sets. For $k \ge 1$, let N(G, k) be the number of clique covers H in G with c(H) = k. The number N(G, k) is referred to as a *clique cover number*. It is clear that N(G, n) = 1 and N(G, k) = 0 for k > n. Define

$$h(G,\mu) = \begin{cases} \sum_{k=1}^{n} N(G,k)\mu^{k}, & \text{if } n \ge 1, \\ 1, & \text{if } n = 0. \end{cases}$$
(2)

The polynomial $h(G, \mu)$ is called the *adjoint polynomial* of G. Observe that $h(G, \mu) = h(G', \mu)$ if $G \cong G'$. Hence $h(G, \mu)$ is a well-defined graph-function. The notion of the adjoint polynomial of a graph was introduced by Liu [6]. Note that the adjoint polynomial is a special case of an F-polynomial [2].

Two graphs G and H are said to be *adjointly equivalent* if they have the same adjoint polynomial. In this case we write $G \sim_h H$. The equivalence class determined by a graph G is denoted by $[G]_h$. A graph G is said to be *adjointly unique* if $[G]_h = \{G\}$. Note that

$$\alpha(G,k) = N(\overline{G},k), \qquad k = 1, 2, \cdots, n.$$
(3)

It follows that

Theorem 1.1 (i) $G \sim H$ iff $\overline{G} \sim_h \overline{H}$; (ii) $[G] = \{H | \overline{H} \in [\overline{G}]_h\};$ (iii) G is chromatically unique if and only if \overline{G} is adjointly unique.

Hence the goal of determining [G] for a given graph G can be realised by determining $[\overline{G}]_h$. Thus, as has been observed in [6, 7, 8, 9, 10, 11, 12], if e(G) is very large, it may be easier to study $[\overline{G}]_h$ rather than [G].

Section 2 computes some clique cover numbers that are used to study two invariants for adjoint polynomials. These invariants, $R_1(G)$ and $R_2(G)$, are the subject matter of Sections 3 and 4 respectively. For a polynomial $f(x) = x^n + b_1 x^{n-1} + b_2 x^{n-2} + \cdots + b_n$, define

$$R_1(f) = \begin{cases} -\binom{b_1}{2} + b_1, & \text{if } n = 1, \\ b_2 - \binom{b_1}{2} + b_1, & \text{if } n \ge 2 \end{cases}$$
(4)

and

$$R_2(f) = b_3 - {b_1 \choose 3} - (b_1 - 2) \left(b_2 - {b_1 \choose 2}\right) - b_1,$$

where $b_k = 0$ for k > n. For any graph G, define

$$R_i(G) = R_i(h(G,\mu)) \tag{5}$$

for each $i \in \{1, 2\}$. It is clear that $R_i(G)$ is an invariant for adjointly equivalent graphs, since N(G, k) is an invariant for each positive integer k. The invariant $R_1(G)$ was introduced by Liu [6] and used by him and others to study adjoint uniqueness of graphs. In particular in [12] Liu and Zhao showed that $R_1(G) \leq 1$ for any connected graph G, and characterised the connected graphs G with $R_1(G) \geq 0$. They also established the chromatic uniqueness of certain dense graphs. In Section 3 we obtain a recursive formula and a sharper upper bound for $R_1(G)$. We also show for which graphs this upper bound is met. In Section 4 we obtain alternative formulae for $R_2(G)$ which enable us to compute $R_2(G)$ for some specific graphs. In a subsequent paper we use both $R_1(G)$ and $R_2(G)$ to determine adjoint equivalence classes of certain graphs and confirm a conjecture of Liu [9] that P_n is adjointly unique for each even $n \neq 4$.

2 Computation of some clique cover numbers

In this section we calculate the clique cover numbers N(G, n-k) for k = 0, 1, 2, 3 in order to obtain an expression for each $R_i(G)$, where i = 1, 2.

Theorem 2.1 [7] For any graph G with order n,

(i) N(G, n) = 1 if $n \ge 1$; (ii) N(G, n - 1) = e(G) if $n \ge 2$; (iii) $N(G, n - 2) = t(G) + {e(G) \choose 2} - \sum_{x \in V(G)} {d_G(x) \choose 2}$ if $n \ge 3$.

For $x \in V(G)$, let $\Delta_G(x)$ (or simply $\Delta(x)$) be the number of triangles in Gwhich include x. For any graphs G and Q, let $n_G(Q)$ (or simply n(Q)) denote the number of subgraphs in G which are isomorphic to Q. Thus $n_G(K_2) = e(G)$ and $n_G(K_3) = t(G)$. In particular, let $p_k(G) = n_G(P_k)$, i.e., the number of paths of order k in G.

The next result gives an expression for N(G, v(G) - 3).

Theorem 2.2 For any graph G with order n, we have

$$N(G, n-3) = \binom{e(G)}{3} + p_4(G) + 5t(G) + n(K_4) - \sum_{x \in V(G)} d(x) \Delta(x) + e(G) \left(t(G) - \sum_{x \in V(G)} \binom{d(x)}{2} \right) + 2 \sum_{x \in V(G)} \binom{d(x) + 1}{3}.$$
 (6)

Proof. By definition, N(G, n-3) is the number of clique covers H in G with c(H) = n - 3. Since v(H) = n, each component of H is of order at most 4, we find that H is one of the following types of graphs:

(i)
$$3K_2 \cup (n-6)K_1$$
,
(ii) $K_3 \cup K_2 \cup (n-5)K_1$,
(iii) $K_4 \cup (n-4)K_1$.

Thus

$$N(G, n-3) = n_G(3K_2) + n_G(K_3 \cup K_2) + n_G(K_4).$$

Observe that

$$n_G(K_3 \cup K_2) = \sum_{\Delta xyz \text{ in } G} (e(G) - d(x) - d(y) - d(z) + 3),$$

where the sum is taken over all triangles xyz in G. Hence

$$n_G(K_3 \cup K_2) = (e(G) + 3)t(G) - \sum_{x \in V(G)} d(x) \triangle(x).$$

Now consider the number $n_G(3K_2)$. The following figure shows all possible graphs with size 3 and no isolated vertices.

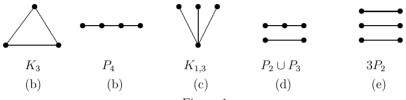


Figure 1

Observe that

$$n_G(K_{1,3}) = \sum_{x \in V(G)} \binom{d(x)}{3}$$

and

$$\sum_{x \in V(G)} {d(x) \choose 2} (e(G) - d(x)) = 3n_G(K_3) + 2n_G(P_4) + n_G(P_2 \cup P_3).$$

Thus

$$n_G(3K_2) = \binom{e(G)}{3} - n_G(K_3) - n_G(P_4) - n_G(K_{1,3}) - n_G(P_2 \cup P_3)$$

= $\binom{e(G)}{3} - \sum_{x \in V(G)} \binom{d(x)}{3} - \sum_{x \in V(G)} \binom{d(x)}{2} (e(G) - d(x))$
+ $2n_G(K_3) + n_G(P_4)$
= $\binom{e(G)}{3} + 2 \sum_{x \in V(G)} \binom{d(x) + 1}{3} - e(G) \sum_{x \in V(G)} \binom{d(x)}{2}$
+ $2n_G(K_3) + n_G(P_4).$

The result is then obtained.

3 The Invariant $R_1(G)$

By Theorem 2.1 and the definition of $R_1(G)$, we have

Lemma 3.1 For any graph G,

$$R_1(G) = t(G) + e(G) - \sum_{x \in V(G)} {\binom{d_G(x)}{2}}.$$
(7)

Corollary $R_1(G) = 0$ if e(G) = 0.

By Lemma 3.1, the next result is obtained.

Lemma 3.2 For any graph G with components G_1, G_2, \dots, G_k ,

$$R_1(G) = \sum_{i=1}^k R_1(G_i).$$
 (8)

If e(G) = 0, then $R_1(G) = 0$. We shall find a recursive expression for $R_1(G)$ when e(G) > 0. For $x, y \in V(G)$, let $N_G(x, y)$ (or simply N(x, y)) denote the set

$$(N(x) \cup N(y)) - \{x, y\}.$$

Observe that

$$|N_G(x,y)| = \begin{cases} d(x) + d(y) - |N(x) \cap N(y)|, & \text{if } xy \notin E(G) \\ d(x) + d(y) - |N(x) \cap N(y)| - 2, & \text{if } xy \in E(G) \end{cases}$$

Lemma 3.3 For any graph G and $xy \in E(G)$, we have

$$R_1(G) = R_1(G - xy) + 1 - |N_G(x, y)|.$$
(9)

Proof. By (7), we have

$$R_{1}(G) - R_{1}(G - xy) = t(G) - t(G - xy) + (e(G) - e(G - xy)) \\ - \left(\binom{d_{G}(x)}{2} - \binom{d_{G}(x) - 1}{2} \right) - \left(\binom{d_{G}(y)}{2} - \binom{d_{G}(y) - 1}{2} \right)$$

$$= |N_{G}(x) \cap N_{G}(y)| + 1 - (d_{G}(x) - 1) - (d_{G}(y) - 1) \\ = 1 - |N_{G}(x, y)|.$$

By Lemma 3.3, we find a sufficient condition for two graphs G and G' to satisfy $R_1(G) = R_1(G')$.

Lemma 3.4 Let xy be an edge in G with $N_G(x) \cap N_G(y) = \emptyset$. Let G' be any graph obtained from G by replacing the edge xy by a path containing no vertices of $V(G) - \{x, y\}$. Then

$$R_1(G) = R_1(G'). (10)$$

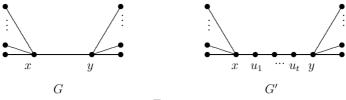


Figure 2

Proof. Let G' be the graph obtained from G by replacing the edge xy by the path with t+2 vertices, as shown in Figure 2. To prove the lemma, it suffices to show that $R_1(G') = R_1(G)$ for t = 1. Let t = 1. Assume that $d_G(x) = 1 + a$ and $d_G(y) = 1 + b$. By Lemma 3.3, we have

$$R_{1}(G') = R_{1}(G' - xu_{1}) + 1 - (1 + a)$$

$$= (R_{1}(G' - xu_{1} - u_{1}y) + 1 - b) - a$$

$$= R_{1}((G - xy) \cup K_{1}) + 1 - a - b$$

$$= R_{1}(G - xy) + 1 - a - b$$

$$= R_{1}(G).$$

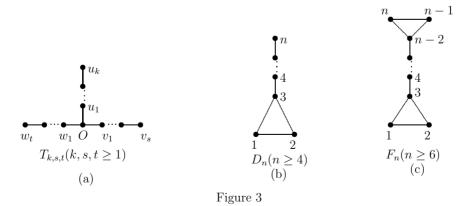
By using Lemmas 3.3 and 3.4, it is easy to compute $R_1(G)$ for some special graphs. Let $K_4 - e$ be the graph obtained from K_4 by deleting one edge.

Lemma 3.5 (i)
$$R_1(P_1) = 0$$
 and $R_1(P_t) = 1$ for $t \ge 2$.
(ii) $R_1(K_3) = 1$, $R_1(K_4) = -2$ and $R_1(K_4 - e) = -1$.
(iii) $R_1(C_k) = 0$ for $k \ge 4$.

For positive integers k, s and t, let $T_{k,s,t}$ be the graph in Figure 3(a). Let

$$\mathcal{T}' = \{ T_{k,s,t} | k \ge s \ge t \ge 1 \}.$$

Let D_n and F_n be the graphs shown in Figure 3 (b) and (c).



Theorem 3.1 [12] Let G be a connected graph. Then $R_1(G) \leq 1$ and (i) $R_1(G) = 1$ if and only if $G \in \{K_3\} \cup \{P_n | n \geq 2\}$, (ii) $R_1(G) = 0$ if and only if $G \in \{K_1\} \cup T' \cup \{C_n, D_n | n \geq 4\}$, and (iii) $R_1(G) = -1$ with $e(G) \geq v(G) + 1$ if and only if $G \in \{K_4 - e\} \cup \{F_n | n \geq 6\}$. \Box

From Theorem 3.1, we observe that for any connected graph G, if $G \not\cong K_3$ and $R(G) \geq -1$, then $e(G) + R_1(G) \leq v(G)$. We shall show that for any connected graph G, if $G \not\cong K_4$ and $R_1(G) \leq -2$, then $e(G) + R_1(G) \leq v(G) - 1$. First we establish the following result.

Theorem 3.2 For any connected graph G, if $G \not\cong K_4$, then

$$R_1(G) \le 2(v(G) - e(G)) + 1. \tag{11}$$

Proof. For any graph G, let

$$\phi(G) = R_1(G) - 2(v(G) - e(G)).$$

We have to show that for any connected graph G, if $G \not\cong K_4$, then

$$\phi(G) \le 1. \tag{12}$$

By Lemma 3.1, we have

$$\phi(G) = t(G) - 2v(G) + 3e(G) - \sum_{x \in V(G)} \binom{d(x)}{2} \\
= t(G) + \sum_{x \in V(G)} (3d(x)/2 - 2) - \frac{1}{2} \sum_{x \in V(G)} d(x)(d(x) - 1) \\
= t(G) - \frac{1}{2} \sum_{x \in V(G)} (d(x) - 2)^2.$$
(13)

It follows that $\phi(G) \leq 1$ if $t(G) \leq 1$. Hence (12) holds for connected graphs G with $e(G) \leq 4$. Note that $\phi(K_4) = 4 - \frac{1}{2} \times 4 = 2$.

Suppose that H is a connected graph with minimum size such that $H \not\cong K_4$ and $\phi(H) \geq 2$. We prove that such a graph H does not exist.

Claim 1: For any $x \in V(H)$, if $N_H(x) = \{y, z\}$, then $N_H(y) \cap N_H(z) \neq \{x\}$.

Suppose that $N_H(y) \cap N_H(z) = \{x\}$. Let H' be the graph $H \cdot xy$. Observe that H' has an edge which is not contained in any triangle, which implies that $H' \not\cong K_4$. Since e(H') < e(H), we have $\phi(H') \leq 1$. By Lemma 3.4, $R_1(H) = R_1(H')$. Since v(H) = v(H') + 1 and e(H) = e(H') + 1, we have $\phi(H) = \phi(H') \leq 1$, a contradiction. The claim holds.

Claim 2: $\delta(H) \ge 2$.

Suppose that $d_H(x) = 1$ and $N_H(x) = \{y\}$. Let H' = H - x. By (13),

$$\phi(H) - \phi(H') = -1/2 - 1/2(d_H(y) - 2)^2 + 1/2(d_H(y) - 3)^2 = 2 - d_H(y).$$

Since $H \neq K_2$ (as $\phi(K_2) = -1$) and $d_H(y) \neq 2$ by Claim 1, we have $d_H(y) \geq 3$. Hence $\phi(H) \leq \phi(H') - 1$. If $H' \ncong K_4$, then as H' is connected and e(H') < e(H), we have $\phi(H') \leq 1$; thus $\phi(H) \leq 0$. If $H' \cong K_4$, we have $\phi(H) \leq 1$. Both cases lead to a contradiction.

Claim 3: *H* does not contain a bridge.

Suppose that xy is a bridge of H. Let H_1 and H_2 be the two components of H - xy. By (13),

$$\phi(H) - \phi(H_1) - \phi(H_2) = 5 - d_H(x) - d_H(y)$$

Observe that H_1, H_2 are connected. Thus $\phi(H_i) \leq 1$ if $H_i \not\cong K_4$. By Claim 1 and 2, $d_H(x), d_H(y) \geq 3$. Let $x \in V(H_1)$ and $y \in V(H_2)$. Notice that $d_H(x) = 4$ if $H_1 \cong K_4$ and $d_H(y) = 4$ if $H_2 \cong K_4$. Hence $\phi(H) \leq 1$, a contradiction. The claim holds. **Claim 4**: For each $xy \in E(H), |N_H(x, y)| \leq 2$.

Suppose that $|N_H(x,y)| \ge 3$ for some $xy \in E(H)$. By Lemma 3.3,

$$R_1(H) \le R_1(H - xy) - 2.$$

By Claim 3, H-xy is connected. Since H-xy is not complete and e(H-xy) < e(H), we have $\phi(H-xy) \leq 1$. Hence by the definition of $\phi(H)$,

$$\phi(H) - \phi(H - xy) = R_1(H) - R_1(H - xy) + 2 \le 0,$$

which implies that $\phi(H) \leq 1$, a contradiction. The claim follows. Claim 5: t(H) = 0.

If *H* contains a subgraph isomorphic to $K_4 - e$, then v(H) = 4 by Claim 4, which implies that either $H \cong K_4$ or $H \cong K_4 - e$. But $\phi(K_4 - e) = 1$, a contradiction. Thus *H* does not contain any subgraph isomorphic to $K_4 - e$.

Suppose that xyz is a triangle in H. If d(x) = d(y) = d(z) = 2, then $H \cong K_3$ and $\phi(H) = 1$, a contradiction. By Claim 4, $d(x), d(y), d(z) \leq 3$. Now say d(x) = 3. Let $xw \in E(H)$, where $w \notin \{y, z\}$. Since $\delta(H) \geq 2$, we have $d(w) \geq 2$. Since Hdoes not contain any subgraph isomorphic to $K_4 - e$, we have $|N_H(x, w)| \geq 3$, which contradicts Claim 4. Hence Claim 5 holds.

Since t(H) = 0 by Claim 5, we have $\phi(H) \leq 0$ by (13), a contradiction. Hence H does not exist.

Recall from Theorem 3.1 that $R_1(G) \leq 1$. By Theorems 3.1 and 3.2, we have

Corollary 3.1 For any connected graph G with $G \notin \{K_3, K_4\}$, (i) if $-1 \leq R_1(G) \leq 1$, then $R_1(G) \leq v(G) - e(G)$ with equality if and only if

$$G \in \{K_4 - e\} \cup \{P_n, C_{n+1}, D_{n+2}, F_{n+4} | n \ge 2\}.$$

(ii) if
$$R_1(G) \leq -2$$
, then $R_1(G) \leq v(G) - e(G) - 1$.

Proof. The result of (i) follows from Theorem 3.1.

(ii) If $R_1(G) \leq -2$, then by Theorem 3.2,

$$v(G) - e(G) \ge R_1(G)/2 - 1/2 = R_1(G) - R_1(G)/2 - 1/2 \ge R_1(G) + 1 - 1/2,$$

which implies $v(G) - e(G) \ge R_1(G) + 1$. Thus (ii) holds.

4 The Invariant $R_2(G)$

Theorem 4.1 For any graph G,

$$R_{2}(G) = 2\sum_{x \in V(G)} {d(x) \choose 3} - \sum_{x \in V(G)} d(x) \triangle_{G}(x) -e(G) + p_{4}(G) + 7t(G) + n_{G}(K_{4}).$$
(14)

Proof. Let v(G) = n. For $f(\mu) = h(G, \mu)$, we have $b_i = N(G, n-i), i \ge 1$. Observe that

$$b_2 - {b_1 \choose 2} = t(G) - \sum_{x \in V(G)} {d(x) \choose 2},$$

and

$$\sum_{x \in V(G)} \binom{d(x) + 1}{3} = \sum_{x \in V(G)} \binom{d(x)}{3} + \sum_{x \in V(G)} \binom{d(x)}{2}.$$

The result is then obtained from (5) and (6).

The term $p_4(G)$ can be expressed in terms of $d_G(x)$ and t(G). Thus there is another expression for $R_2(G)$.

Theorem 4.2 For any graph G,

$$R_{2}(G) = 2 \sum_{x \in V(G)} {d(x) \choose 3} - \sum_{x \in V(G)} d(x) \Delta_{G}(x) - e(G) + 4t(G) + n_{G}(K_{4}) + \sum_{xy \in E(G)} (d_{G}(x) - 1)(d_{G}(y) - 1).$$
(15)

Proof. For $xy \in E(G)$, let $p_4(xy)$ be the number of paths of the form uxyv in G, where $u \neq v$. Observe that

$$p_4(xy) = (d(x) - 1)(d(y) - 1) - \triangle(xy),$$

where $\Delta(xy)$ is the number of triangles in G containing xy. Thus

$$p_4(G) = \sum_{xy \in E(G)} p_4(xy)$$

=
$$\sum_{xy \in E(G)} ((d(x) - 1)(d(y) - 1) - \triangle(xy))$$

=
$$\sum_{xy \in E(G)} (d(x) - 1)(d(y) - 1) - 3t(G).$$

The result is then obtained.

Corollary 4.1 If G is K_3 -free, then

$$R_2(G) = 2\sum_{x \in V(G)} {d(x) \choose 3} - e(G) + \sum_{xy \in E(G)} (d(x) - 1)(d(y) - 1).$$

Corollary 4.2 If G_1, G_2, \dots, G_k are the components of G, then

$$R_2(G) = \sum_{i=1}^k R_2(G_i).$$

Proof. It follows from Theorem 4.2.

Let Y_n denote the graph $T_{n-3,1,1}$, where $n \ge 4$. By applying Theorem 4.2, we have

Corollary 4.3 (i)
$$R_2(P_1) = 0$$
, $R_2(P_2) = -1$ and $R_2(P_n) = -2$ for $n \ge 3$;
(ii) $R_2(K_3) = -2$ and $R_2(C_n) = 0$ for $n \ge 4$;
(iii) $R_2(Y_4) = -1$ and $R_2(Y_n) = 0$ for $n \ge 5$;
(iv) $R_2(D_4) = 0$ and $R_2(D_n) = 1$ for $n \ge 5$;
(v) $R_2(F_6) = 5$ and $R_2(F_n) = 4$ for $n \ge 7$;
(vi) $R_2(K_4 - e) = 3$ and $R_2(K_4) = 7$.

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