On digraphs with unique walks of closed lengths between vertices

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Abstract

It is known that regular digraphs of degree d, diameter k and unique walks of length not smaller than h and not greater than k between all pairs of vertices ([h, k]-digraphs), exist only for h = k and h = k - 1, if $d \geq 2$. This paper deals with the problem of the enumeration of [k-1,k]-digraphs in the case of diameter k=2 or degree d=2. It is shown, using algebraic techniques, that the line digraph $L K_{d+1}$ of the complete digraph K_{d+1} is the only [1,2]-digraph of degree d, that is to say the only digraph —up to isomorphisms— whose adjacency matrix A fulfills the equation $A + A^2 = J$, where J denotes the all-one matrix. As a consequence, we deduce that there does not exist any other almost-Moore digraph of diameter k = 2 with all selfrepeat vertices apart from Kautz digraph. In addition, the cycle structure of a [k-1,k]-digraph is studied. Thus, a formula that provides the number of short cycles (cycles of length $\leq k$) of such a digraph is obtained. From this formula, using graphical arguments, the enumeration of [k-1,k]-digraphs of degree 2 and diameter not greater than 4 is concluded.

1 Introduction

Digraphs with unique walks of length in a fixed interval [h, k] between all pairs of vertices, were introduced by Plesník and Znám [19] and, since then, they have been extensively studied, because they are suitable models for dense interconnection networks. Some problems related to such a class of digraphs, denoted by [h, k]digraphs, have been solved by means of algebraic (spectral) techniques applied to the equation

$$A^h + A^{h+1} + \dots + A^k = J \tag{1}$$

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satisfied by their corresponding adjacency matrix A, where J denotes the all-one matrix. Thus, for instance, the regularity of a [h, k]-digraph and the computation of its order $n = d^{h} + d^{h+1} + \cdots + d^{k}$, where d is its degree, are easily derived from (1) (see Hoffman and McAndrew [12].) Furthermore, the problem of their existence was completely solved by Bosák [4], who proved that, for degree $d \geq 2$, there only exist [k, k]- and [k-1, k]-digraphs, where k is precisely their diameter. Among such digraphs there are Good-De Bruijn (see [16]) and Kautz (proposed in [13, 14]), which, in fact, can be seen as iterated line digraphs of the complete digraph of degree d with loops and without loops, respectively (see Fiol et al. [7, 8].) With regard to their structure and enumeration, several results concerning the class of [k, k]-digraphs¹ have already been obtained (see [9], [15], [16].) Thus, it is known that, in general, the family of Good-De Bruijn digraphs is not the only one whose adjacency matrix Afulfills the equation $A^k = J$. In particular, Mendelsohn [16] proved that for diameter k=2 and degree d=3 there are exactly six non-isomorphic digraphs of that type. Furthermore, Fiol et al. [9] described two direct methods of constructing digraphs of this kind by adequately modifying Good-De Bruijn digraphs.

In this paper, we focus our attention on the study of [k-1,k]-digraphs and, in particular, we face the problem of their enumeration in the case of diameter k = 2or degree d = 2. First, in Section 2, we introduce the concepts of (l,m)-reachable digraph and reachable digraph with a maximum delay m (an extension of the notions of l-reachable and equi-reachable digraph, respectively, introduced in [9]), we give a characterization of the class of reachable digraphs with a maximum delay m and we meet the class of [k-1,k]-digraphs when we consider (k, 1)-reachable digraphs with maximum order. In Section 3, we prove, by means of algebraic methods, that there is only one [1,2]-digraph of degree d, namely the line digraph $L K_{d+1}$ of the complete digraph K_{d+1} and, as a consequence, we deduce that there does not exist any other almost-Moore digraph of diameter k = 2 with all selfrepeat vertices apart from Kautz digraph (this solves a question formulated by Baskoro et al. in [1].)

Some results about the cycle structure of a [k-1, k]-digraph are presented in Section 4. Thus, using the spectrum of a [k-1, k]-digraph and the property that any of its closed walks of length $l \leq k$ is either a cycle or a repeated cycle, we obtain a formula for the computation of its short cycles, that is to say cycles of length not greater than k. Such a formula, which almost coincides with that given in [16] for [k, k]-digraphs, allow us to know the number of cycles of length a divisor of k-1 or k, which constitute a partition of the set of vertices of a [k-1, k]-digraph. Finally, in Section 5 we make a first attempt to study the enumeration of [k-1, k]-digraphs of degree d = 2 and diameter $k \geq 3$. Using the previous results about their cycle structure, together with other graphical arguments, we prove that for k = 3, 4 there is only one digraph of degree d = 2 of that type.

We conclude this introduction fixing the terminology used throughout the paper. Thus, a digraph G consists of a finite non-empty set V(G) of objects called vertices

¹The class of [k, k]-digraphs has been studied by Mendelsohn in [16] as "UPP digraphs" (digraphs with the unique path property of order k) and by Conway and Guy in [5] as "tight precisely k-steps digraphs".

and a set E(G) of ordered pairs of vertices called *arcs*. The *order* n of G is the cardinality of V(G), n = |V(G)|. If (u, v) is an arc, it is said that v is adjacent from u [u is adjacent to v] and also that a is incident from u to v [a is an out-arc of u and an *in-arc* of v. The set of vertices which are adjacent from [to] a given vertex v, also called successors of v, is denoted by $\Gamma^+(v)$ [$\Gamma^-(v)$] and its cardinality is the outdegree of v, $d^+(v) = |\Gamma^+(v)|$ [in-degree of v, $d^-(v) = |\Gamma^-(v)|$]. A vertex v is isolated if $d^+(v) = d^-(v) = 0$. A digraph is regular of degree d (d-regular) if, for any vertex $v, d^+(v) = d^-(v) = d$. A walk of length h from a vertex u to a vertex $v (u \to v \text{ walk})$ is a sequence of vertices $u = u_0, u_1, \ldots, u_{h-1}, u_h = v$ such that (u_{i-1}, u_i) is an arc. A *circuit* is a closed walk with all its arcs distinct. A cycle of length h > 0 (h-cycle) is a closed walk with h distinct vertices. A repeated cycle is a closed walk originated by the repetition of a cycle. The length of a shortest $u \to v$ walk is the distance from u to v. Its maximum value over all pairs of vertices is the diameter k of the digraph. A *p*-generalized cycle is a digraph G such that V(G) can be partitioned into *p*-parts V_i such that all adjacent vertices from vertices of V_i belong to V_{i+1} , for *i* modulo $p \ge 1$. The reader is referred to Chartrand and Lesniak [6] for additional graph concepts.

2 Reachable digraphs with a fixed maximum delay

Let G be a strongly connected digraph with adjacency matrix A and let m be a non-negative integer. Then, we say that G is (l, m)-reachable if $l \ge m$ is the smallest integer such that for each pair of vertices $u, v \in V(G)$, there is at least one $u \to v$ walk of length in the interval [l - m, l], that is to say $A^{l-m} + \cdots + A^l \ge J$ and $A^{l'-m} + \cdots + A^{l'} \ge J$, if l' < l. Moreover, if m is the smallest non-negative integer such that G is (l, m)-reachable for some $l \ge m$, then we say that G is reachable with a maximum delay m. Notice that in the case m = 0 the notions of *l*-reachable and equi-reachable digraph, presented in [9], are recovered. We also remark that if G is a [k-1, k]-digraph of degree $d \ge 2$, then G is (k, 1)- and (k + 1, 0)-reachable, since $A^{k-1} + A^k = J$ and $A^{k+1} = AJ - A^k \ge (d-1)J$, and, consequently, G is reachable with a maximum delay 0 (equi-reachable).

Now, we state a characterization of reachable digraphs with a maximum delay m, which is in fact a natural extension of the one given in [9] for the case m = 0, taking into account that any digraph G can be seen as a *p*-generalized cycle with p = 1.

Proposition 1. A strongly connected digraph G is reachable with a maximum delay m iff p = m + 1 is the greatest integer such that G is a p-generalized cycle.

Proof. We will prove that a strongly connected digraph G is reachable with a maximum delay $\leq m$ unless G is a p-generalized cycle with $p \geq m + 2$. Since any reachable digraph G with a maximum delay $\leq m$ has walks of lengths that belong to a set of cardinal m + 1 between all its pairs of vertices, the condition of not being a p-generalized cycle with a number of parts $p \geq m + 2$ is clearly necessary. The sufficiency of such a condition can be derived using the same reasoning detailed in [9] for m = 0.

From now on, we focus our attention on the study of largest (l,m)-reachable digraphs. Thus, if G is (l,m)-reachable with a maximum out-degree d, then its order n satisfies the following inequality

 $n \le d^{l-m} + \dots + d^l,$

since this bound is the maximum number of distinct walks of length not smaller than l-m and not greater than l from any vertex of G. To attain this bound there must be exactly one walk of length in the interval [l-m, l] between each pair of vertices. So, any digraph with such a property is an [l-m, l]-digraph and, consequently, it only exists for m = 0, 1, if $d \geq 2$.

3 Enumeration of [1, 2]-digraphs

In this section we present the main result of this paper. We prove that Kautz digraphs of diameter 2 are the unique [1,2]-digraphs —up to isomorphisms. This means that all $n \times n$ binary matrices that are solution of the equation $A + A^2 = J$ are of the form PAP^t , where A is the adjacency matrix of the line digraph $L K_{d+1}$, where $n = d + d^2$, and P is a permutation matrix.

We use, as an auxiliary result, the following characterization of regular line digraphs, which is based on Heuchenne's condition (see [11]). Such a condition says that a digraph G is a line digraph iff every pair of vertices u, v of G satisfy that their corresponding sets of successors $\Gamma^+(u), \Gamma^+(v)$ are either disjoint or equal. In matrix terms, a (0, 1)-matrix A represents the adjacency matrix of a line digraph G iff the rows of A are either mutually ortogonal or identical.

Lemma 1. A regular digraph G of degree $d \ge 1$ and order n is a line digraph iff the rank of its adjacency matrix A is equal to $\frac{n}{d}$.

Proof. Let us suppose that G is a regular digraph of degree d, order n = dn' and such that the rank of its adjacency matrix A is n'. Then, we take n' row vectors of A that constitute a basis of the row-space of A. Since the total number of entries of these (0, 1)-vectors that are equal to 1 is exactly n, we deduce that these vectors must be mutually ortogonal. Otherwise, A would have a null column, which is impossible since JA = dJ (G is d-regular.) Moreover, since each row has just d elements equal to 1, we deduce that any other row is identical to one of the basis. Hence, G satisfies Heuchenne's condition and, consequently, G is a line digraph.

Conversely, let G be the line digraph of a digraph G' of order n' without isolated vertices. Then, since G is regular, we have that G' is also a regular digraph with the same degree d and, consequently, the order of G is n = n'd. Therefore, using Heuchenne's condition, we deduce that the adjacency matrix A of G has a set of n' mutually ortogonal rows such that any other row is equal to one of those. Hence, rank $A = n' = \frac{n}{d}$.

From the relation between the rank of a square matrix A and the dimension of its null-space, we can reformulate the previous characterization, saying that a regular

digraph G of degree $d \ge 2$ and order n is a line digraph, iff 0 is an eigenvalue of G with geometric multiplicity equal to $n - \frac{n}{d}$.

Besides, from the construction of a line digraph it follows that if the adjaceny matrix A of a digraph G fulfills the equation P(A) = J, where P(x) is a polynomial, then the adjacency matrix A_L of LG satisfies the equation $A_LP(A_L) = J$. The converse is also true if we assume that G has no isolated vertices. Taking into account this property, the following result is derived.

Lemma 2. Let G be a digraph without isolated vertices. Then, LG is a [k-1,k]-digraph iff G is a [k-2, k-1]-digraph, and both of them have the same degree.

Now, we can characterize the class of [1, 2]-digraphs.

Theorem 1. There is only one [1, 2]-digraph of degree d, namely $L K_{d+1}$.

Proof. Let A be the adjacency matrix of a [1,2]-digraph G of degree d and order $n = d + d^2$. Then, A fulfills the equation $A + A^2 = J$ and, consequently, tr A = 0. From these relations, it follows that the characteristic polynomial of A is

$$\det(xI - A) = (x - d)x^{n - d - 1}(x + 1)^d.$$

Now, we will prove, using the previous lemmas, that G is the line digraph $L K_{d+1}$. In fact, for d = 1 the result is trivial. So, from here on, we assume $d \ge 2$.

Since 0 is an eigenvalue of A with algebraic multiplicity equal to n - d - 1, we have that rank $A \ge d + 1 = \frac{n}{d}$. Suppose that rank A > d + 1. Then, A would have at least d + 2 column vectors u_1, \ldots, u_{d+2} which are linearly independent. However, the relation $A + A^2 = (A + I)A = J$ implies that $(A + I)u_i = (1, \ldots, 1)^t$, from which we deduce that the vectors $u_i - u_1$ belong to the null-space of A + I. Then, since these vectors $u_i - u_1$ ($2 \le i \le d + 2$) are also linearly independent, we conclude that 0 must be an eigenvalue of A + I with algebraic multiplicity not smaller than d + 1. But, this is impossible because the characteristic polynomial of A + I is

$$\det(xI - (A + I)) = \det((x - 1)I - A) = (x - (d + 1))(x - 1)^{n - d - 1}x^{d}.$$

Therefore, since rank of A must be equal to d + 1, from Lemma 1 it turns out that G is a line digraph.

Let G' be a digraph —without isolated vertices— such that G = LG'. Then, from Lemma 2, we have that G' is a [0,1]-digraph of degree d. Hence, G' is the complete digraph K_{d+1} .

The previous result answers a question formulated by Baskoro et al. in [1] about the enumeration of almost-Moore digraphs (also called (d, k)-digraphs) of diameter k = 2 with all selfrepeat vertices. A (d, k)-digraph is a regular directed graph of degree d > 1, diameter k > 1 and order one less than the (unattainable) Moore bound. Every (d, k)-digraph G has the property that for each vertex $v \in V(G)$ there exists only one vertex, denoted by r(v) and called the *repeat* of v, such that there are exactly two $v \to r(v)$ walks of length less than or equal to k (one of them must be of length k.) If r(v) = v, which means that v is contained in exactly one k-cycle, v is called a selfrepeat of G. The map r, which assigns the vertex r(v) to each vertex $v \in V(G)$, is an automorphism of G and its associated permutation matrix P is related to the adjacency matrix A of G by means of the following equation $I + A + \cdots + A^k = J + P$ (see, for instance, [2], [3] and [17].) Therefore, the notion of (d, k)-digraph with all selfrepeat vertices (P = I) and the concept of [1, k]-digraph are equivalent. Hence, (d, k)-digraphs with all selfrepeat vertices do only exist for diameter k = 2 (see [2] and [4].) Moreover, in such a case, using Theorem 1, we have that $L K_{d+1}$ is the only (d, 2)-digraph with all selfrepeat vertices. Furthermore, using the results about the structure of almost-Moore digraphs with selfrepeat vertices of any (d, 2)-digraph of degree d, $3 \leq d \leq 12$, are selfrepeats. Hence, the enumeration of almost-Moore digraphs of diameter 2 and small degree is reduced to the resolution of the equation $A + A^2 = J$ (given in Theorem 1.)

4 Short cycles in a [k-1,k]-digraph

Several properties about the cycle structure of a [k, k]-digraph are already known. Thus, in [16] it is proved, using matrix techniques, that each closed walk of a [k, k]digraph of length $\leq k$ is either a cycle or a repeated cycle. We will extend such a result, by means of graphical techniques, and, as a consequence, we will deduce that any [k-1, k]-digraph shares this same property. From this fact and the knowledge of the spectrum of a [k-1, k]-digraph, we will obtain a formula for the number of its short cycles.

Lemma 3. If G is a digraph such that for each pair of (not necessarily different) vertices u, v there is at most one $u \to v$ walk of length k, then each closed walk of G of length $\leq k$ is either a cycle or a repeated cycle.

Proof. Let $C: u_0, u_1, \ldots, u_l$ be a closed walk of G of length $l \leq k$ such that is neither a cycle nor a repeated cycle. Let u_i and u_j be the two first repeated vertices of the sequence C, where $0 \leq i < j < l$. Then, $C_1: u_i, u_{i+1}, \ldots, u_j$ is a cycle of length j-iand $C_2: u_j, u_{j+1}, \ldots, u_l, u_1, \ldots, u_i$ is a closed walk of length l + i - j. Since C is not a repeated cycle, it can be seen that the two concatenation of these sequences, C_1C_2 and C_2C_1 , represent different $u_i \rightarrow u_i$ closed walks of length l. But, then, there would be two $u_i \rightarrow u_{i+r}$ walks of length k, where $k - l \equiv r \pmod{j-i}$, which is impossible. \Box

Corollary 1. Every closed walk of a [k-1, k]-digraph of length $\leq k$ is either a cycle or a repeated cycle.

Moreover, since the adjacency matrix A of a [k-1, k]-digraph G of degree d fulfills the equation $A^k + A^{k-1} = J$, we can deduce that the characteristic polynomial of A is $(x-d)x^{n-1-d}(x+1)^d$, where $n = d^{k-1} + d^k$. Therefore,

$$tr A^{l} = d^{l} + (-1)^{l} d.$$
(2)

Working with such identities and using the previous corollary, we obtain the following result for the computation of the number of distinct short cycles of a [k - 1, k]-digraph. We say that two cycles $C_1 : u_0, u_1, \ldots, u_l$ and $C_2 : u'_0, u'_1, \ldots, u'_l$ are equal iff one can be obtained from the other by means of a rotation, that is to say $C_2 : u_i, u_{i+1}, \ldots, u_{i+l}$, for some *i* modulo *l*. Although the number of such cycles depends on the degree *d* of the digraph as well as on the length *l*, we simply denote it by c(l).

Theorem 2. Let G be a [k-1,k]-digraph of degree d. Then, the number c(l) of distinct cycles of G of length l, $l \leq k$, is given by the formula

$$c(l) = \begin{cases} 0, & \text{if } l = 1, \\ \binom{d+1}{2}, & \text{if } l = 2, \\ \frac{1}{l} \sum_{m \mid l} \mu(\frac{l}{m}) d^m, & \text{if } 3 \le l \le k, \end{cases}$$
(3)

where $\mu(l)$ denotes the Möbius function.

Proof. From Corollary 1 and from identity (2), we have that

tr
$$A^{l} = d^{l} + (-1)^{l} d = \sum_{m|l} m \cdot c(m)$$
, if $l \le k$.

Therefore, applying Möbius's inversion formula [18], we deduce that

$$c(l) = \frac{1}{l} \sum_{m|l} \mu(\frac{l}{m}) d^m + \frac{1}{l} d \sum_{m|l} \mu(\frac{l}{m}) (-1)^m.$$
(4)

Now, we compute the auxiliary function

$$f(l) = \sum_{m|l} \mu(\frac{l}{m})(-1)^m = \sum_{m|l} \mu(m)(-1)^{\frac{l}{m}}.$$

If l is odd, then

$$\begin{aligned} f(l) &= -\sum_{m|l} \mu(m) = -\sum_{m|l} \operatorname{tr} \Phi_m(x) = -\operatorname{tr} \prod_{m|l} \Phi_m(x) \\ &= -\operatorname{tr} (x^l - 1) = \begin{cases} -1, & \text{if } l = 1, \\ 0, & \text{otherwise,} \end{cases} \end{aligned}$$

where $\Phi_m(x)$ denotes the *m*-th cyclotomic polynomial and tr $\Phi_m(x)$ represents the sum of all its complex roots.

Likewise, if $l = 2^{e}l'$, where $e \ge 1$ and l' is odd, then

$$f(l) = \sum_{m|l'} \mu(m)(-1)^{2^e \frac{l'}{m}} + \sum_{m|l'} \mu(2m)(-1)^{2^{e-1} \frac{l'}{m}} = \sum_{m|l'} \mu(m)(1-(-1)^{2^{e-1}}).$$

Clearly, f(l) = 0, if $l = 2^e l'$ and e > 1. Moreover, if l = 2l', then

$$f(l) = 2\sum_{m|l'} \mu(m) = \begin{cases} 2, \text{ if } l = 2\\ 0, \text{ otherwise.} \end{cases}$$

Hence, f(1) = -1, f(2) = 2 and f(l) = 0 if l > 2. The proof is concluded by substituting these values into (4).

We notice that the number c(l) of distinct cycles of length $l \leq k$ in a [k-1, k]digraph does not depend on k. Moreover, if $l \geq 3$, then c(l) turns out to be equal to the number of distinct cycles of the same length in a [k, k]-digraph with equal degree (see [16]). Furthermore, the number of 2-cycles of a [k-1, k]-digraph is equal to the number of loops plus the number of 2-cycles of a [k, k]-digraph. We also remark that if $d \geq 2$, then $c(l) \geq 1$, which implies that any [k-1, k]-digraph of degree $d \geq 2$ has cycles of each length $l, 2 \leq l \leq k$. The problem of the existence of longer cycles has been solved in the particular case of Kautz digraphs by Villar, who proved in [21] that any Kautz digraph has cycles of any length, except for 1 and n-1, where n is its order.

Taking into account that each vertex of a [k-1, k]-digraph is included in exactly one closed walk of length k-1 or k, we can also derive some other properties about the cycle structure of such a digraph.

Corollary 2. If G is a [k-1,k]-digraph of degree d, then the following statements hold.

(i) There exists a partition of the set of vertices of G into $\binom{d+1}{2}$ cycles of length 2 and $\frac{1}{l} \sum_{m|l} \mu(\frac{l}{m}) d^m$ cycles of length l, for each $l \ge 3$ a divisor of (k-1) or k. Moreover, the total number $\mathcal{N}(k)$ of these cycles is given by the expression

$$\mathcal{N}(k) = \frac{1}{k} \sum_{l|k} \phi(\frac{k}{l}) d^{l} + \frac{1}{k-1} \sum_{l|(k-1)} \phi(\frac{k-1}{l}) d^{l} - d,$$
(5)

where $\phi(i)$ stands for the Euler function.

(ii) Each arc not contained in a cycle of length a divisor of k belongs to a unique closed walk of length k + 1, which is either a cycle, a repeated cycle or the concatenation of two arc-disjoint cycles.

Proof. From the definition of a [k-1,k]-digraph G, the existence of a partition of V(G) into cycles of length a divisor of k-1 or k is derived. The number c(l) of such cycles of each length l is given by Theorem 2 and its total number $\mathcal{N}(k)$ can be deduced as follows. Since gcd(k-1,k) = 1 and c(1) = 0, we have that $\mathcal{N}(k) = \sum_{l|k} c(l) + \sum_{l|(k-1)} c(l)$. Therefore, using Theorem 2, we obtain that

$$\mathcal{N}(k) = \sum_{l|k} \frac{1}{l} \sum_{m|l} \mu(\frac{l}{m}) d^m + \sum_{l|(k-1)} \frac{1}{l} \sum_{m|l} \mu(\frac{l}{m}) d^m - d.$$

Then, taking into account that

$$\sum_{l|n} \frac{1}{l} \sum_{m|l} \mu(\frac{l}{m}) d^m = \sum_{m|n} d^m \sum_{\substack{m|l \\ l|n}} \frac{1}{l} \mu(\frac{l}{m}) = \frac{1}{n} \sum_{m|n} d^m \sum_{\substack{l'|\frac{n}{m} \\ l'|\frac{n}{m}}} \frac{\frac{n}{m}}{l'} \mu(l')$$

and using the identity $\sum_{d|n} \frac{n}{d} \mu(d) = \phi(n)$ (see [18]), we can deduce (5).

Besides, given an arc uv of G, and since there exists exactly one $v \to u$ walk of length k - 1 or k, we have that uv is included in exactly one closed walk of length k or k + 1. Therefore, using Corollary 1, we deduce that if uv is not included in a cycle of length a divisor of k, then uv belongs to a unique closed walk of length k + 1. Let $C : u_0, u_1, \ldots, u_{k+1}$ be such a closed walk and let us assume that C is neither a cycle nor a repeated cycle. Then, if u_i and u_j are the first two repeated vertices of the sequence C, we have that $C_1 : u_i, u_{i+1}, \ldots, u_j$ is a cycle and $C_2 : u_j, u_{j+1}, \ldots, u_{k+1}, u_0, \ldots, u_i$ is a cycle or the repetition of a cycle C'_2 distinct from C_1 . But such a repetition is impossible because, otherwise, we would have more than one walk of length k between u_i and u_{i-1} . Hence, C is equal to the concatenation of the two cycles C_1 and C_2 , which are arc-disjoint.

We point out that the computation of the number of *Kautz necklaces*, that is to say cycles of a Kautz digraph of length a divisor of its diameter k or k - 1, was previously proved by Tvrdík in [20], by using combinatorial techniques. Here, we have extended such a result (5) to the class of [k - 1, k]-digraphs.

Thus, if G is a [k-1,k]-digraph of degree d and diameter $k \leq 5$, then G has a vertex partition into c(l) distinct *l*-cycles, where such numbers are shown in Table 1. For k = 2, since $L K_{d+1}$ is the only [1,2]-digraph, we have that the number of *l*-cycles equals the number of circuits of K_{d+1} of the same length, which represents the number of closed sequences of length *l* of (d + 1)-ary digits: $0, \ldots, d$ such that two consecutive digits are different and subsequences of length two are all different. It can be verified that these computations turn out to be equal to c(l) for $l \leq 4$ but not for l = 5, that is to say the expression of c(l), given in Theorem 2 for $l \leq k$, can be extended, in the case k = 2, for l = k + 1, k + 2. We do not know if such an extension can be generalized, neither do we know if there exists a formula for the computation of long cycles (cycles of length > k) of a [k - 1, k]-digraph.

k	num. 2-cycles	num. 3-cycles	num. 4-cycles	num. 5-cycles
3	$\frac{d^2+d}{2}$	$\frac{d^3-d}{3}$		
4	$\frac{d^2+d}{2}$	$\frac{d^3-d}{3}$	$\frac{d^4-d^2}{4}$	
5	$\frac{d^2+d}{2}$	$\frac{d^3-d}{3}$	$\frac{d^4-d^2}{4}$	$\frac{d^5-d}{5}$

Table 1: Number of short cycles of a [k-1, k]-digraph of diameter $k \leq 5$.

5 The case d = 2

In this section, we will illustrate how the previous results about the cycle structure of a [k-1,k]-digraph can be used in order to find the enumeration of such digraphs in the case of degree d = 2 and small diameter k.

Applying Corollary 2, in the case k = 3, we derive the following result.

Lemma 4. If G is a [2,3]-digraph of degree d, then the following statements hold.

(i) There is a partition

$$\mathcal{P} = \{C_1^2, C_2^2, \dots, C_{\frac{d^2+d}{2}}^2, C_1^3, C_2^3, \dots, C_{\frac{d^3-d}{3}}^3\}$$

of the set of vertices of G, where $C_i^j = \{v_1^i, v_2^i, \ldots, v_j^i\} \subset V(G)$ is such that $v_1^i, v_2^i, \ldots, v_j^i, v_1^i$ represents a *j*-cycle of G.

(ii) Each arc of G not contained in a cycle of length 2 or 3 belongs to a unique cycle of G of length 4.

In particular, every [2, 3]-digraph G of degree d = 2 must contain the subdigraph $G_{\mathcal{P}}$ shown in Figure 1. Moreover, the strongly connected components of the subdigraph $\overline{G}_{\mathcal{P}}$ of G induced by the remaining arcs —one incident arc from each vertex are cycle digraphs of order 4. The following technical lemma says how the arcs of $\overline{G}_{\mathcal{P}}$ have to be placed. Using this lemma, we will show that there is only one way of constructing (up to isomorphisms) a [2,3]-digraph of degree 2. In order to simplify the notation, from now on each part C_i^j will be identified by its associated cycle.

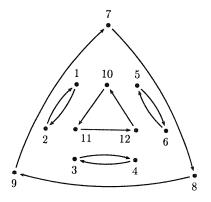


Figure 1: The subdigraph $G_{\mathcal{P}}$.

Lemma 5. Let G be a [2,3]-digraph of degree 2. Let $G_{\mathcal{P}}$ be the subdigraph of G induced by its 2-cycles (C_1^2, C_2^2, C_3^2) and its 3-cycles (C_1^3, C_2^3) , and let $\overline{G}_{\mathcal{P}}$ the subdigraph of G induced by its remaining arcs. Then, the arcs of $\overline{G}_{\mathcal{P}}$ satisfy the following properties:

- (i) Each arc of \$\overline{G}_{P}\$ is incident from a vertex of a 2-cycle [3-cycle] to a vertex of a 3-cycle [2-cycle]. Moreover, the arcs of \$\overline{G}_{P}\$ incident from vertices of \$G\$ included in the same 2-cycle [3-cycle] are incident to vertices of \$G\$ included in distinct 3-cycles [2-cycles]. (See Figure 2 (I).)
- (ii) Let u₁, u₂, u₃ be a walk of G_P. If u₁ is included in a 2-cycle [3-cycle] of G, then u₃ is contained in a distinct 2-cycle [3-cycle] of G. (See Figure 2 (II).)

(iii) Let u_1v_1 $[v_1u_1]$ be an arc of $\overline{G}_{\mathcal{P}}$ such that u_1 is included in a 2-cycle C_i^2 : u_1, u_2, u_1 and v_1 is included in a 3-cycle C_j^3 : v_1, v_2, v_3, v_1 of G. Then, v_3u_2 $[u_2v_2]$ is an arc of $\overline{G}_{\mathcal{P}}$. (See Figure 2 (III).)

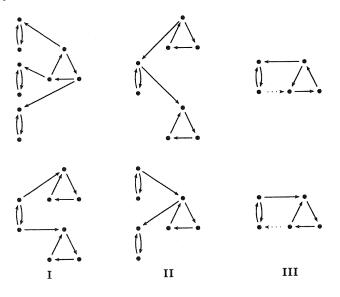


Figure 2: Conditions about the arcs of $\overline{G}_{\mathcal{P}}$.

Proof. Taking into account the definition of a [k-1,k]-digraph G, it can be seen that, given a vertex v_0 included in a cycle $C: v_0, v_1, \ldots, v_l$ of G of length l a divisor of k-1, then each out-arc of v_0 , except for v_0v_1 , is an in-arc of a vertex included in a cycle of length a divisor of k. Since the converse digraph of G (derived from Gby changing the orientation of its arcs) is also a [k-1,k]-digraph, every property satisfied by the out-arcs of a vertex v of G is also fulfilled by its in-arcs. In particular, if k = 3, then each arc of $G_{\mathcal{P}}$ that is incident from [to] a vertex included in a 2-cycle must also be incident to [from] a vertex included in a cycle of length 3. Moreover, since in the case d = 2 there are equal numbers of vertices included in 2-cycles and in 3-cycles, we have that each arc of $\overline{G}_{\mathcal{P}}$ joins two vertices included in cycles of $G_{\mathcal{P}}$ of distinct length. Now, let us suppose that $\overline{G}_{\mathcal{P}}$ has two arcs u_1v_1 and u_2v_2 , where $C^2: u_1, u_2, u_1$ and $C^3: v_1, v_2, v_3, v_1$ are cycles of $G_{\mathcal{P}}$ of length 2 and 3, respectively. Then, the sequences u_1, v_1, v_2, v_3 and u_1, u_2, v_2, v_3 are two different $u_1 \rightarrow v_3$ walks of length 3, which is impossible. The proof of (ii) is quite similar and property (iii) is a consequence of the first two. Π

The application of the previous lemma to the subdigraph shown in Figure 1 allows us to conclude the enumeration of [2,3]-digraphs of degree 2.

Proposition 2. There is only one [2,3]-digraph of degree 2, namely $L^2 K_3$.

Proof. Let G be a [2,3]-digraph of degree 2. We will prove, using Heuchenne's condition, that G is a line digraph, from which we will deduce, taking into account Lemma 2 and Theorem 1, that G is the Kautz digraph of diameter k = 3 and degree d = 2.

Let u_1 and v_1 be two vertices of G such that $\Gamma^+(u_1) \cap \Gamma^+(v_1) \neq \emptyset$. From property (i) of Lemma 5, we have that u_1 and v_1 must belong to cycles of $G_{\mathcal{P}}$ of distinct length. Let us assume that u_1 belongs to the cycle $C^2 : u_1, u_2, u_1$ and that v_1 is included in the 3-cycle $C^3 : v_1, v_2, v_3, v_1$. Therefore, since u_1v_2 $[v_1u_2]$ is an arc of $\overline{G}_{\mathcal{P}}$, we deduce, using property (iii) of the previous lemma, that $\Gamma^+(u_1) = \Gamma^+(v_1)$ and, consequently, G is a line digraph. Then, from Lemma 2, we have that G = L G'where G' is a [1,2]-digraph of degree 2. Hence, using Theorem 1, we obtain that Gis $L^2 K_3$. (See Figure 3.)

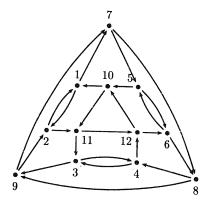


Figure 3: The digraph $L^2 K_3$.

We notice that while Good-De Bruijn digraph of degree 2 and diameter 3 is one of the three non-isomorphic [3, 3]-digraphs of such a degree (see [9]), Kautz digraph with equal parameters is the only [2, 3]-digraph of degree 2. This may strengthen the idea that, as the order becomes closer to the Moore bound, there are fewer digraphs. With a more detailed reasoning it can be seen that $L^3 K_3$ is the only [3, 4]-digraph of degree 2. However, because of the ad hoc nature of these techniques, it may be worthwhile to find other approaches to the problem of the enumeration of [k-1, k]digraphs for $k \geq 3$.

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