On permutation patterns with constrained gap sizes

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Abstract

We consider avoidance of permutation patterns with designated gap sizes between pairs of consecutive letters. We call the patterns having such constraints distant patterns (DPs), and we show their relation to other pattern notions investigated in the past. New results on DPs with 2 and 3 letters are obtained. Furthermore, we show how one can use DPs to prove two former conjectures of Kuszmaul without a computer. In addition, we deduce a surprising relation between the sets of permutations avoiding the classical patterns 123 and 132 by looking at a class of DPs with tight gap constraints. Some interesting analogues of the Stanley–Wilf former conjecture for DPs are also discussed.

1 Introduction

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The notion of patterns in permutations has applications when solving a variety of enumeration problems in different areas including algorithms (sortable permutations), algebraic geometry (Kazhdan–Lusztig polynomials and Schubert varieties), statistical mechanics (generalizations of gas models), computational biology and even chemistry.

The work of Simion and Schmidt [36] was the first systematic study of permutation pattern avoidance. Before that, in 1968, Knuth [27] showed that the permutation π can be sorted by a stack if and only if π avoids 231, and that the stack-sortable permutations are enumerated by the Catalan numbers.

This work is related to the idea of considering permutation patterns for which we have additional constraints for the distance between some pairs of consecutive letters in the pattern's occurrences. For example, when we want two consecutive letters to have a positive gap size in any occurrence of a pattern, we will write \square between these two letters. For instance, a permutation $\pi = \pi_1 \pi_2 \cdots \pi_n$ avoids $1\square 2$ if there is no $1 \le i < j \le n$ with $j \ge i + 2$, such that $\pi_i < \pi_j$. In this paper, we obtain some interesting results related to these types of patterns, and show its usefulness when interpreting other results combinatorially.

1.1 Definitions

Most of the definitions and terms we use can be found in [6]. An *n*-permutation or a permutation of size n is a sequence of the numbers in $[n] := \{1, 2, \dots, n\}$ in which each number occurs exactly once. We will usually use the one-line notation to write permutations, i.e., we will just write the corresponding sequences explicitly. The formulation and the proof of Theorem 3.2 is the only place, where we write permutations as lists of their cycles, i.e., there we use the cycle notation for writing permutations. A sequence of distinct numbers will be just called *sequence*. An *inversion* in a permutation $p_1 p_2 \cdots p_n$ is a pair of numbers (p_i, p_j) , such that i < j and $p_i > p_j$. An occurrence of the classical pattern p in a permutation π is a subsequence in π whose letters are in the same relative order as those in p. We use the std operator, with its name coming from "standardization", to define this formally. Given a sequence of distinct real numbers $u_1u_2\cdots u_k$, define $\operatorname{std}(u_1u_2\cdots u_k)$, which stands for "standardization", to be the permutation $q \in S_k$ such that $u_i < u_j$ if and only if $q_i < q_j$. A permutation is uniquely defined by the set of its inversions, so this condition uniquely defines q. A permutation $\pi \in S_n$ contains a pattern $p = p_1 p_2 \cdots p_k \in S_k$ for $k \leq n$ if there exists $i_1 < i_2 < \cdots < i_k$ such that $\operatorname{std}(\pi_{i_1} \pi_{i_2} \cdots \pi_{i_k}) = p$. Otherwise, π avoids p. For instance, the permutation 32514 has two occurrences of the pattern 231, namely the subsequences 351 and 251 and it avoids the pattern 123.

Vincular patterns allow the requirement that some of the letters that are at consecutive positions in the pattern must correspond to numbers at consecutive positions in the permutation. We indicate this requirement by underlining the letters in the pattern corresponding to numbers at consecutive positions. For instance, if the pattern $3\underline{12}$ occurs in a permutation π , then the letters in π that correspond to 1 and 2 are adjacent. For example, the permutation 621543 has only one occurrence of the pattern $3\underline{12}$, namely the subsequence 615. Vincular patterns were first studied systematically by Babson and Steingrímsson in [3], under the name "generalized patterns". When all the entries of a pattern are required to occur in adjacent positions, we will call the pattern consecutive.

We will denote by S_n the set of permutations of size n and by $\operatorname{Av}_n(\Pi)$ the set of permutations of size n avoiding each of the patterns in the set Π . To refer to the size of this set, we will write $av_n(\Pi) := |\operatorname{Av}_n(\Pi)|$. When Π contains a single pattern p, we will write $\operatorname{Av}_n(p)$ and $av_n(p)$, instead of $\operatorname{Av}_n(\{p\})$ and $av_n(\{p\})$. The set of all permutations, of any size, avoiding the pattern p or each of the patterns in the set Π will be denoted by $\operatorname{Av}(p)$ and $\operatorname{Av}(\Pi)$, respectively. If each number in a sequence σ_1 is greater than each number in a sequence σ_2 , then we write $\sigma_1 > \sigma_2$. Otherwise, we write $\sigma_1 \not> \sigma_2$. Finally, we write $\sigma[< x]$ (respectively, $\sigma[> x]$) for the subsequence of elements of σ that are less than (respectively, greater than) x.

A permutation class is a set C of permutations such that every pattern occurring in a permutation in C is also in C. Certainly, $\operatorname{Av}(\Pi)$ is a permutation class for any set of permutations Π . Two permutations $\sigma, \tau \in S_k$ are Wilf-equivalent if, for each n, $av_n(\sigma)$ equals $av_n(\tau)$. Whenever we say generating function, we will mean ordinary generating function. We will write \Box^r to denote a gap with at least r letters, with $\Box := \Box^1$. Therefore avoiding the pattern $12\Box^23$ will be the same as avoiding occurrences xyz of the classical pattern 123 for which there are at least two other letters between y and z. Similarly, avoiding \Box^2123 (respectively, $123\Box^2$) means that we do not have such occurrences of xyz with at least 2 other letters before x (respectively, after z) in the permutation. The patterns containing \Box^r symbols will be called $distant\ patterns$ (DPs) and we will use gap and $gap\ size$ for the space between two consecutive letters of a pattern and for its size. Any distant pattern can be written in the form

$$\square^{r_0}q_1\square^{r_1}q_2\square^{r_2}\cdots\square^{r_{k-1}}q_k\square^{r_k},$$

where each r_i is a non-negative integer and $q_1q_2\cdots q_k\in S_k$. We will also consider tight constraints and we will underline the corresponding part of the pattern in case of a tight constraint as, for example, in $1\square^4 23$ to denote that we want to avoid the pattern 123 with gap size exactly 4 between the letters 1 and 2. DPs without any tight constraints are classical distant patterns while DPs having at least one tight constraint are vincular distant patterns. If we take a classical pattern q and require the minimal gap size to be the same number r for all pairs of consecutive letters, we will write $\operatorname{dist}_r(q)$ and we will call these uniform distant patterns. For example, $\operatorname{dist}_3(312) = 3\square^3 1\square^3 2$. Note that DPs generalize classical patterns since $q = \operatorname{dist}_0(q)$ for any classical pattern q. DPs generalize vincular patterns, as well, since one can write any vincular pattern as a vincular distant pattern. Finally, when we say that a distant pattern has size n, we mean that the number of its non-square letters is n. For example $2\square 1\square^2 3$ is a distant pattern of size 3.

1.2 Related work and motivation

The idea of arbitrary constraints for the gap sizes between any two consecutive pattern letters is not new, even though not much has been written on the subject and the topic seems to be not so well explored yet. The thesis of Ghassan Firro [13] defines a more general concept of permutation patterns with gap constraints unifying many popular pattern notations. He calls them distanced patterns or d-patterns and uses a different notation. The distanced patterns described there also allow requiring a gap size to be at most some given number r. The thesis itself enumerates the patterns of the kind $xy \square z$ using both a direct bijection and an analytical approach. We have included this result in Section 4.1. The paper of Hopkins and Weiler [18] describes the concept of uniform distant patterns under the name of gap patterns and obtains an important result related to them, as a corollary of their work on pattern avoidance over posets. We state this corollary in Section 4.

In his book dedicated to pattern avoidance [25], Kitaev mentions patterns containing the \square -symbol, where the work of Hou and Mansour on the so-called Horse permutations [19] is listed. There, the authors proved that the permutations avoiding both the classical pattern 132 and the pattern $1\square 23$ are in one to one correspondence with the so-called Horse paths.

In [22], Kitaev introduces partially ordered patterns (POPs) and partially ordered

generalized patterns (POGPs), which further generalize classical DPs (respectively, vincular DPs). While in classical patterns, all of the letters form one totally ordered set (e.g. in 123, 1 < 2 < 3), in POPs this set is partially ordered. In an occurrence of a distant pattern, an element at the place of a \square is incomparable to any other element, which shows us that POPs (and POGPs that allow tight constraints) are indeed generalizations. If we have a classical distant pattern or a vincular distant pattern, we could easily write it as a POP (respectively, POGP) by replacing each square with a letter in its own group. POPs were studied in the context of permutations, words and compositions in a series of papers [17, 20, 21, 22, 23, 24, 25] including a recent work [16] of Gao and Kitaev where a systematic search of connections between sequences in the Online Encyclopedia of Integer Sequences (OEIS) and permutations avoiding POPs of size 4 and 5 was conducted. Two other works [7] and [8] study avoidance of non-consecutive occurrence of a pattern and this has connections with both POPs and DPs. Another generalization of the DPs are the so-called place-difference-value patterns [26].

1.3 Organization of the results

Section 2 contains proofs of two basic facts about DPs, which set the stage for the later work.

Section 3 considers the avoidance of classical DPs of size 2, namely $2\Box 1$ and the more general case $2\Box^r 1$. While $av_n(2\Box 1) = av_n(1\Box 2)$ was shown to be given by the Fibonacci numbers F_{n+1} in many ways in the past, we use this old result and a technique that can be generally applied to any classical DP to obtain a summation formula for F_{n+1} , for which we believe is new (see Section 3.1). In addition, we establish a bijection between the permutations in $Av_n(2\Box^r 1)$ and the permutations in S_n for which any two elements in a cycle differ by at most r.

Section 4 considers the DPs of size 3. Previous results by Firro and Mansour[13, 15], as well as by Hopkins and Weiler[18] are first listed. Then, we describe briefly our approach towards the enumeration of $Av_n(1\square 3\square 2)$. The main tool that was used was the block-decomposition approach initiated by Mansour and Vainshtein [32]. All the details in the proposed approach will be described in a forthcoming article.

In Section 5.1, we classify the patterns of the form $\underline{ab}\Box c$ and $\underline{a}\Box\underline{bc}$. Recursive formulas for $av_n(1\Box 23)$ and $av_n(1\Box 32)$ are obtained, which help us to show that $1\Box 23$ is the only pattern among $1\Box 23$, $12\Box 3$, $1\Box 2\Box 3$ and 123 that is avoided by fewer permutations of size n, compared to the same pattern after we switch the places of 2 and 3. This is somewhat surprising since, as we know, $av_n(123) = av_n(132)$ (see [36]). We study consecutive DPs in Section 5.2. We will show a simple but surprising relation between these patterns and the question of avoidance of arithmetic progressions in permutations for which still not much is known.

In Section 6, we present direct solutions, without using a computer, for two former conjectures of Kuzsmaul [28], which give the generating functions for the permutations avoiding two big sets of size four patterns. These solutions use formulations of the conjectures in terms of distant patterns.

Section 7 is dedicated to some analogues of the Stanley–Wilf former conjecture for distant patterns. We conclude with Section 8, which lists some open problems.

2 Two basic facts about distant patterns

Avoidance of classical distant patterns can be formulated as a statement about simultaneous avoidance of classical patterns. For example, avoiding $1\square 2$ is equivalent to the simultaneous avoidance of the 3-letter classical patterns $\{123, 132, 213\}$. In the general case, we have the fact below, where

$$x^{(y)} := x(x-1)\cdots(x-y+1) = \frac{x!}{(x-y)!}$$

denotes the falling factorial.

Theorem 2.1. The avoidance of $q = \Box^{r_0} q_1 \Box^{r_1} q_2 \Box^{r_2} \cdots \Box^{r_{k-1}} q_k \Box^{r_k}$, where $\sum_{j=0}^k r_j = S$, is equivalent to the simultaneous avoidance of $(S+k)^{(S)}$ classical patterns of size S+k.

Proof. If a permutation π avoids q, then π avoids all classical patterns of the kind

$$q' = p_{0,1} \cdots p_{0,r_0} q'_1 p_{1,1} \cdots p_{1,r_1} q'_2 \cdots q'_{k-1} p_{k-1,1} \cdots p_{k-1,r_{k-1}} q'_k p_{k,1} \cdots p_{k,r_k},$$

where $p_{i,j}$ are S distinct numbers from 1 to S+k and q'_1, q'_2, \dots, q'_k are the remaining k numbers from 1 to S+k and they are in the same relative order as the numbers $q_1q_2\cdots q_k$, i.e., $\operatorname{std}(q'_1q'_2\cdots q'_k)=q_1q_2\cdots q_k$. Indeed, any occurrence of such classical pattern would be an occurrence of q. The number of such classical patterns is $\frac{(S+k)!}{k!}$ since the relative order of exactly k of the positions is fixed. Conversely, if π avoids all the listed classical patterns, then it does not have an occurrence of q. Assume the opposite and take one such occurrence of q: $\pi_{x_1}\pi_{x_2}\cdots\pi_{x_k}$. We know that $x_1>r_0$, $x_2-x_1>r_1$, etc. Select any r_0 letters of π , appearing before π_{x_1} , any r_1 letters between π_{x_1} and π_{x_2} and so forth. You will obtain a subsequence $s=\pi_{y_1}\pi_{y_2}\cdots\pi_{y_{S+k}}$ of π and $\operatorname{std}(s)$ must be one of the already listed classical patterns of size S+k, which is a contradiction.

In fact, we can restrict our attention to DPs without leading or trailing squares.

Theorem 2.2. For any $r_1, r_2 > 0$ and a distant pattern q, we have

$$av_n(\Box^{r_1}q\Box^{r_2}) = n^{(r)}av_{n-r}(q),$$

where $r := r_1 + r_2$.

Proof. If $\sigma = \sigma_1 \sigma_2 \cdots \sigma_n \in \operatorname{Av}_n(\square^{r_1} q \square^{r_2})$, then

$$\operatorname{std}(\sigma_{r_1+1}\cdots\sigma_{n-r_2})\in\operatorname{Av}_{n-r}(q).$$

Conversely, any $\sigma = \sigma_1 \sigma_2 \cdots \sigma_n$ for which $\sigma_1, \ldots, \sigma_{r_1}, \sigma_{n-r_2}, \ldots, \sigma_n$ are any r numbers in [n] and for which $\operatorname{std}(\sigma_{r_1+1} \cdots \sigma_{n-r_2}) \in \operatorname{Av}_{n-r}(q)$, would be such that $\sigma \in \operatorname{Av}_n(\square^{r_1}q\square^{r_2})$, since any possible occurrence of q in σ would have either fewer than r_1 other elements in front of it or fewer than r_2 elements after it.

3 Classical DPs of size 2

By Theorem 2.2, it suffices to consider $1\square^r 2$ and $2\square^r 1$ as the only DPs of size 2. When r = 0, we have $av_n(1\square^0 2) = av_n(2\square^0 1) = av_n(21) = 1$. When r = 1, one can obviously see that $av_n(2\square 1) = av_n(1\square 2)$, by applying the reverse map. The seminal paper of Simion and Schmidt [36] finds this number of avoiders. There, they showed that $av_n(123, 132, 213) = F_{n+1}$ (see Theorem 2.1). Below, we present a proof in the language of DPs.

Theorem 3.1. For $n \geq 3$,

$$av_n(2\square 1) = F_{n+1},$$

i.e., the (n+1)st Fibonacci number.

Proof. If $p_1p_2\cdots p_n\in \operatorname{Av}_n(2\square 1)$, then either $p_n=n$ or $p_{n-1}=n$ since otherwise n will participate in an inversion that is not of consecutive letters. If $p_n=n$ then $p_1p_2\cdots p_{n-1}$ must be in $\operatorname{Av}_{n-1}(2\square 1)$, and for each permutation in $\operatorname{Av}_{n-1}(2\square 1)$, we obtain a permutation in $\operatorname{Av}_n(2\square 1)$ after appending the letter n at the end. Thus, we have $av_{n-1}(2\square 1)$ permutations in $\operatorname{Av}_n(2\square 1)$ ending with n. If $p_{n-1}=n$, then we must have $p_n=n-1$ to prevent n-1 from forming a prohibited inversion with p_n . Thus, in this second case we must have $p_{n-1}=n$ and $p_n=n-1$ and the prefix $p_1p_2\cdots p_{n-2}$ must be a permutation in $\operatorname{Av}_{n-2}(2\square 1)$. For each such permutation in $\operatorname{Av}_{n-2}(2\square 1)$, we obtain a new one in $\operatorname{Av}_n(2\square 1)$ by appending n and then n-1 to it. Therefore $av_n(2\square 1)=av_{n-1}(2\square 1)+av_{n-2}(2\square 1)$. It remains to note that $av_3(2\square 1)=3$ and $av_4(2\square 1)=5$.

Theorem 3.2 below gives insight about the general case, when we have an arbitrary value of r and it can help us to find $av_n(2\square^r 1)$, when r > 1. Note that $av_n(2\square^r 1) = n!$, when $r \ge n - 1$.

Theorem 3.2. The permutations in $Av_n(2\square^r 1)$ are in one-to-one correspondence with the permutations in S_n for which, when written in a cycle notation, any two elements in a cycle differ by at most r.

Proof. We give a bijection that is a variation of Foata's fundamental bijection between permutations written in cycle notation and permutations written in one-line notation [38, Chapter 1.3]. In particular, if Foata's map is denoted by F, our bijection maps a permutation π to $rc(F^{-1}(rc(\pi^{-1})))$, where rc denotes the reverse-complement map. We believe that the description of the map given below is more easily comprehensible.

Let $X = \operatorname{Av}_n(2\square^r 1)$ be the set of all permutations in S_n that do not have inversions at distance greater than r in their one-line notation. Let Y be the set of those permutations in S_n for which any two elements in the same cycle differ by at most r. We will describe a bijective map $f: X \to Y$. Consider $p = p_1 p_2 \cdots p_n \in X$. We will show how to obtain the standard form of f(p) written in cycle notation, i.e., the minimal element of every cycle is at its first position and the cycles are ordered in increasing order of their minimal elements. Below is the description of f.

The number p_1 is at position 1, so let us look at the set of positions of all numbers with which p_1 is in inversion: $1, 2, \ldots, p_1 - 1$. Denote their positions with $i_1, i_2, \ldots, i_{p_1-1}$, respectively. These positions are not bigger than r+1, since $p \in X$. Then take $(1, i_{p_1-1}, i_{p_1-2}, \ldots, i_1)$ to be the first cycle in the standard form of the cycle decomposition for f(p). Then, let j be the minimal number that is not already used in this cycle decomposition, and let the numbers $p_j - 1, \ldots, p_1 + 1$ be at positions $j_{p_j-1}, j_{p_j-2}, \ldots, j_{p_1+1}$. Take $(j, j_{p_j-1}, j_{p_j-2}, \ldots, j_{p_1+1})$ as the next cycle in the standard form of the cycle decomposition for f(p) and continue in the same way afterwards. Note that the length of some of those cycles might be 1.

Here are two examples:

- If n = 9, r = 3 and p = 352149867, then f(352149867) = (134)(25)(6798).
- If n = 8, r = 4 and p = 41352867, then f(41352867) = (1352)(4)(687).

Obviously, f maps each $\sigma \in X$ to a permutation $f(\sigma)$ such that any two numbers in the same cycle of $f(\sigma)$ differ by at most r, since these two numbers correspond to indices of two numbers, in the one-line notation of σ , which are in inversion in $\sigma \in X$. To prove that f is indeed a bijection, we will describe its inverse. Consider $\pi \in Y$ in its standard cycle decomposition form. If the first cycle of π is $(\pi_1\pi_2\cdots\pi_{i_1})$, then put the number i_1 in the first place, then i_1-1 at position π_2 , i_1-2 at position π_3 and so on. The number 1 will be placed at position π_{i_1} . Note also that π_1 is always 1. Next, go to the next cycle $(\pi_{j_1}\pi_{j_2}\cdots\pi_{j_{i_2}})$. We will determine the positions of the next i_2 numbers: $i_1+1, i_1+2, \ldots, i_1+i_2$. We can see that π_{j_1} must be the least integer not occurring in the first given cycle. We will place at this position, the number i_1+i_2 . Then, i_1+i_2-1 should be placed at position π_{j_2} , i_1+i_2-2 at position π_{j_3} and so on. One can use the two given examples above for verification.

The sequences of the numbers $av_n(2\square^r 1)$, for different fixed values of n and r are respectively rows and columns of the table described in [34, A276837]. For example, in the case r=2, if we denote $a_n=av_n(2\square^2 1)$, then it is not hard to prove the recurrence

$$a_n = a_{n-1} + a_{n-2} + 3a_{n-3} + a_{n-4}.$$

3.1 Recurrence formula for $F_{n+1} = av_n(1\square 2)$

While trying to obtain a general enumeration approach when dealing with DPs, we used a technique which helped us to obtain a recurrence formula for the number of permutations avoiding the distant pattern $1\square 2$. To the best of our knowledge this is a new recurrence formula for the Fibonacci numbers (see Theorem 3.1). We describe this result in the current section. The same technique can be used when trying to enumerate the set of avoiders for other DPs.

The idea is that almost all permutations containing a given distant pattern can be obtained by first taking a permutation containing the corresponding classical pattern and then inserting additional numbers between some of the letters (where we have

the \square symbol) for a certain occurrence of this classical pattern. Let us describe this more concretely with the following algorithm that we will use for the pattern $1\square 2$. Recall that $C_m(p) := S_m \setminus \operatorname{Av}_m(p)$, for any pattern p.

Algorithm 1

- 1. For a given $n \geq 3$ and $j \in [n]$, take any $\pi \in C_{n-1}(12)$.
- 2. Find the leftmost 1 that is part of a classical 12-pattern and insert the number j immediately after it.
- 3. Increase by 1 the numbers $j, j+1, \ldots, n-1$, except the j that we just inserted (unless $j=n, \pi$ contains another j).

This algorithm defines a map $g: A_{n-1} \to B_n$, where $A_{n-1} := C_{n-1}(12) \times [n]$ and $B_n := C_n(1\square 2)$.

Example 3.3. g(3412, 2) = 42513. The leftmost occurrence of the pattern 12 in 3412 is by the first two letters, 3 and 4. Therefore we insert j = 2 immediately after the letter 3 and then increase the 2, 3 and 4 in the original permutation. Note that the added number j always keeps its value in the final image.

We will first need to show that this map is not far from being injective.

Theorem 3.4. No permutation in $B_n = C_n(1\square 2)$, the codomain of the map g, is the image of more than two different elements of $A_{n-1} = C_{n-1}(12) \times [n]$.

Proof. Assume the opposite. Let $\pi = g(\pi^{(1)}, j_1) = g(\pi^{(2)}, j_2) = g(\pi^{(3)}, j_3)$ for three different tuples $(\pi^{(1)}, j_1), (\pi^{(2)}, j_2), (\pi^{(3)}, j_3) \in A_n = C_{n-1}(12) \times [n]$. We can see that j_1, j_2 and j_3 must be different since if two of them, say j_1 and j_2 , are equal then obviously $\pi^{(1)} = \pi^{(2)}$ and we will not have different tuples. Now, we know that without loss of generality $1 < c_{j_1} < c_{j_2} < c_{j_3}$ are three different positions for the three different numbers j_1, j_2, j_3 in the final permutation π . By step 2 of Algorithm 1, after removing j_1 from π , the first occurrence of the classical pattern 12, should be some $\pi_x \pi_y$, where $x = c_{j_1} - 1$. Similarly, after removing j_2 , the first such occurrence should begin at position $c_{j_2} - 1 > c_{j_1} - 1$, but this is only possible if the position $y = c_{j_2}$ since if this is not the case then $\pi_x \pi_y$ would be an occurrence of 12 that begins before position $c_{j_2} - 1$. However, after removing j_3 from π (note that $c_{j_3} > c_{j_2}$), the first occurrence of 12 should begin at position $c_{j_3} - 1 > c_{j_1} - 1$. This is a contradiction. \square

There are many permutations in B_n that are the image of g for two different elements of A_{n-1} . For example, g(312,4) = g(231,1) = 3142. The next fact that we will need gives the number of these permutations.

Theorem 3.5. The number of permutations in B_n , which are an image of exactly two different elements of A_{n-1} , after applying the map g, is given by the sum

$$\sum_{j=3}^{n-1} (j-2)(n-j)(n-j)! = \sum_{k=0}^{n-2} k! - n + 1.$$

Proof. Let $\pi = g(\pi^{(1)}, x) = g(\pi^{(2)}, y)$ for $\pi^{(1)}, \pi^{(2)} \in C_{n-1}(12)$ and $x, y \in [n]$, where $(\pi^{(1)}, x) \neq (\pi^{(2)}, y)$ are different. We saw in the proof of Theorem 3.4 that x and y must be different. Let us denote the positions of x and y in π by i and j, respectively. Without loss of generality, let i < j. We know that after removing $y = \pi_j$ from π , then $\pi_{j-1}\pi_{j+k}$, for some $k \geq 1$, is the first occurrence of the classical pattern 12. Note that some of the two letters, π_{j-1} and π_{j+k} , may have decreased in value by one. Therefore, we should have $\pi_1 > \pi_2 > \cdots > \pi_{j-1}$. If we remove $x = \pi_i$ from π , then $\pi_{i-1}\pi_j$ must be the first occurrence of 12, it follows that we must have $\pi_1 > \pi_2 > \cdots > \pi_{i-2} > \pi_j > \pi_{i-1} > \cdots > \pi_{j-1}$. In other words, the number π_j is between π_{i-2} and π_{i-1} . Otherwise, we would have a 12-occurrence ending at π_j that starts before position i-1. We also have that $\pi_{j-t} > \pi_{j+t}$, for any $t=2,\ldots,j-1$ and $t=1,\ldots,n-j$, because otherwise when removing π_j from π , a 12-occurrence starting before π_{j-1} will be present.

In order to determine π completely, we must known the relations between the n-j+1 numbers $\pi_{j-1}, \pi_{j+1}, \pi_{j+2}, \ldots, \pi_n$. The only constraint that we have is that π_{j-1} is not the largest among them. Thus, when i and j are fixed, we always have (n-j+1)!-(n-j)! possible ways to write π . Therefore, the number of different permutations $\pi \in S_n$ that are an image for two different tuples is

$$\sum_{j=3}^{n-1} \sum_{i=2}^{j-1} [(n-j+1)! - (n-j)!] = \sum_{j=3}^{n-1} (j-2)(n-j)(n-j)!.$$
 (1)

Each term in the latter sum gives us the number of permutations in $\pi \in B_n$, where $\pi = g(\pi^{(1)}, x) = g(\pi^{(2)}, y)$ for some $\pi^{(1)}, \pi^{(2)} \in C_{n-1}(12)$ and $x, y \in [n]$, where x < y and y is at position j in π .

Finally, it is easy to see that the right-hand side of Equation (1) can be written as

$$1!((n-2)! - (n-3)!) + 2((n-3)! - (n-4)!) + 3((n-4)! - (n-5)!) + \cdots + (n-3)!(2!-1!)$$

$$= (n-2)! + (n-3)! + \cdots + 2! - (n-3) = \sum_{k=0}^{n-2} k! - n + 1.$$

As we have seen that no permutation in B_n is counted more than two times, it remains to obtain the number of permutations in B_n that are not an image of g for any permutation in the set of tuples A_{n-1} .

Theorem 3.6. The number of permutations in the set $B_n \setminus g(A_{n-1})$ is:

$$\sum_{k=3}^{n-2} (F_{n-k+1} - 1)k(k-2)(k-2)!,$$

where F_i denotes the i-th Fibonacci number.

Proof. An example of a permutation in B_n that cannot be obtained as an output of the function g (i.e., with Algorithm 1) for any input in A_{n-1} is the permutation 45132. The reason is that before the first occurrence of the distant pattern $1\square 2$, there exist an occurrence of the classical pattern 12. We want to obtain a formula for all permutations in B_n having this property. Consider one such permutation $\pi =$ $\pi_1\pi_2\cdots\pi_n$ and let the first occurrence of $1\square 2$ be $\pi_j\pi_{j+v}$ for some $j\geq 1,\ v\geq 2$, and $j+v \leq n$. Since this is the first such occurrence, observe that $\pi_i > \pi_{i+d}$, for any i < jand $d \geq 0$. Otherwise we would have another occurrence of $1\square 2$, preceding $\pi_i \pi_{i+v}$. Thus $\pi' = \pi_1 \cdots \pi_{j-1}$ must be a permutation of the j-1 numbers in [n-j+2,n]and $\pi'' = \pi_j \cdots \pi_n$ is simply a permutation of [1, n-j+1] for which $\pi_j < \pi_{j+v}$ for some $v \in [1, n-j]$. This means that π' avoids $1\square 2$, but contains the classical pattern 12. Therefore the number of possibilities for π' is $av_{j-1}(1\square 2) - av_{j-1}(12) = F_j - 1$ since $av_{i-1}(1\square 2) = F_i$ (see Theorem 3.1), $av_{i-1}(12) = 1$ and $A_{i-1}(12) \subset A_{i-1}(1\square 2)$. Now, let us denote k = n - j + 1, for clarity. For π'' , we can see that it could be any k-permutation except that it could not start with k or with (k-1)k since if this is the case, then π'' will not start with an occurrence of the $1\square 2$ pattern. The latter means that the possible values for π'' are exactly k! - (k-1)! - (k-2)! = k(k-2)(k-2)!. Summing over k, we obtain the given formula.

Now, we are ready to derive the recurrence formula that we want. Note that $|A_{n-1}| = |C_{n-1}(12) \times [n]| = ((n-1)!-1) \cdot n = n!-n$ gives the number of permutations in B_n that are the image of the map g for at least one tuple in A_{n-1} . Theorems 3.5 and 3.6 give the number of permutations being the image of g for 2 and 0 tuples in A_{n-1} , respectively. We also know that $|B_n| = |C_n(1\square 2)| = n! - F_{n+1}$. Thus using inclusion-exclusion we have:

$$n! - F_{n+1} = (n! - n) - \sum_{j=3}^{n-1} (j-2)(n-j)(n-j)! + \sum_{k=3}^{n-2} (F_{n-k+1} - 1)k(k-2)((k-2)!).$$

After simplifying, we obtain the following recurrence formula for the Fibonacci numbers and respectively for the number of permutations avoiding the distant pattern $1\square 2$ (or $2\square 1$):

$$av_n(1\square 2) = F_{n+1} = n + \sum_{k=1}^{n-3} (n - (k+2)F_{n-(k+1)}) \cdot k \cdot k!.$$

4 Classical DPs of size 3

As we can infer from Theorem 2.1, finding a closed formula for the avoidance set of a distant pattern becomes more complicated as its size increases, because the number of classical patterns that must be simultaneously avoided increases, as well. In this section, we describe some already established results on the DPs of size 3 with one square $(xy\Box z)$ and two squares $(x\Box y\Box z)$. Then, we discuss an approach that we have used to obtain the generating function for $av_n(1\Box 3\Box 2)$, which represents one of the two different Wilf-equivalent classes for patterns of the latter kind.

4.1 Patterns of the kind $xy\Box z$

Consider the patterns $xy \square z$ and $x \square yz$, for some permutation $xyz \in S_3$. The thesis of Firro [13] and two related works [14, 15] give the formula

$$av_n(12\Box 3) = \sum_{k>0} \frac{1}{n-k} \binom{2n-2k}{n-1-2k} \binom{n-k}{k}.$$
 (2)

The same thesis gives two bijections between $12\square 3$ -avoiding permutations and odd-dissections of a given (n+2)-gon, which are dissections with non-crossing diagonals so that no 2m-gons (m>1) appear [34, A049124]. In fact, it turns out that this is the cardinality of the avoidance set for any pattern of the kind $xy\square z$ or $x\square yz$ [13]. We know that all classical patterns in S_3 are avoided by the same number of permutations, namely the Catalan numbers. One might suspect that whenever two classical patterns $p, q \in S_k$ are Wilf-equivalent, then inserting a square at the same place in p and q will produce two Wilf-equivalent distant patterns. The computer simulations suggest that this is true for the Wilf-equivalent patterns $\{1234, 1243, 2143\}$. We have formulated this conjecture in Section 8.

It was shown in [13] that if $xyz \in S_3$, then inserting a square between x and y or between y and z always gives us two Wilf-equivalent patterns. It is worth noting that we do not have a similar fact when considering patterns of size four. For example, $av_7(1\square 234) = 3612 \neq 3614 = av_7(12\square 34)$.

4.2 Patterns of the kind $x \Box y \Box z$

The reverse map and the complement map give us at most two Wilf-equivalent permutation classes: $\{\text{dist}_1(p) \mid p = 132, 231, 213, 312\}$ and $\{\text{dist}_1(p) \mid p = 123, 321\}$. Unlike the case of classical patterns in which these are, in fact, one class [36], here, these classes are different.

Theorem 4.1. (Hopkins and Weiler [18]) For n > 5,

$$av_n(\text{dist}_1(123)) > av_n(\text{dist}_1(132)).$$

The theorem above is a special case of a result of Hopkins and Weiler [18, Theorem 3]. In that work they extend the result of Simion and Schmidt that $av_n(123) = av_n(132)$ from permutations on a totally ordered set to a similar result for pattern avoidance in permutations on partially ordered sets. In particular, they show that $av_{P,n}(132) \leq av_{P,n}(123)$ for any poset P, where $av_{P,n}(q)$ is the number of n-permutations on the poset P avoiding the pattern q. Furthermore, they classify the posets for which equality holds. Here, we state the corollary of their result generalizing Theorem 4.1, as formulated by the authors.

Theorem 4.2. ([18]) For r > 0 and n > 1, we have

$$av_n(\operatorname{dist}_r(123)) \ge av_n(\operatorname{dist}_r(132)),$$

with strict inequality if and only if $r \ge 1$ and $n \ge 2r + 4$.

Note that in the case n = 2r + 3, $av_{2r+3}(\operatorname{dist}_r(123)) = av_{2r+3}(\operatorname{dist}_r(132))$ since there is only one triple of positions where each of these two patterns can occur in a (2r+3)-permutation, namely the positions 1, r+2 and 2r+3. So for each such occurrence, we can exchange the elements at positions r+2 and 2r+3 to get an occurrence of the other pattern. A similar statement about consecutive patterns was first proved in [10] with a simple injection. It states that $av_n(\underline{123}) > av_n(\underline{132})$ for every $n \geq 4$. The listed facts imply that the monotonic pattern 123 is avoided more frequently than 132 when we have two gaps of size exactly 0 between the letters in each occurrence of the two patterns, or when the minimal constraint for each gap is some fixed positive number. However, when patterns with all possible gap sizes must be avoided, we have an equality since $av_n(123) = av_n(132)$. We address this surprising fact in the next section.

Along those lines is another work of Elizalde [12] on consecutive patterns, where he generalizes [10] by proving that the number of permutations avoiding the monotone consecutive pattern $\underline{12\cdots m}$ is asymptotically larger than the number of permutations avoiding any other consecutive pattern of size m. He also proved there that $av_n(\underline{12\dots(m-2)m(m-1)})$ is asymptotically smaller than the number of permutations avoiding any other consecutive pattern of the same size. Similar conjectures can be formulated for distant patterns (see Section 8).

4.2.1 The pattern $1\square 3\square 2$

In this subsection, we will describe an approach that one can use to find the generating function $G(x) = \sum_{n \geq 0} av_n(q)x^n$, where $q = 1 \square 3 \square 2 = \mathrm{dist}_1(132)$. The complete proof is rather technical and will be a focus of a future article. The main ideas in it are sketched in the current section. We will need to define the following sets of permutations:

 $\mathbb{H}_1 := \{\pi \mid \pi \in \operatorname{Av}(q), |\pi| \geq 1 \text{ and } \pi \text{ has no occurrence of } 1 \square \underline{32} \text{ ending at its last position} \},$

 $\mathbb{H}_2 := \{\pi \mid \pi \in \operatorname{Av}(q), |\pi| \geq 1 \text{ and } \pi \text{ has no occurrence of } \underline{13}\square 2 \text{ beginning at its first position} \}.$

Let us also denote the corresponding generating functions with

$$H_i(x) = \sum_{k=1}^{\infty} h_i(k) x^k,$$

where $h_i(n)$ is the number of permutations of size n in \mathbb{H}_i . Now, we can describe a useful decomposition for the permutations in $\operatorname{Av}_n(q)$ which is similar, but more complicated, to the one given in [13] for the permutations in $\operatorname{Av}_n(13\square 2)$.

Theorem 4.3. For all $n \geq 1$, $\pi = \alpha n \beta \in Av_n(q)$ if and only if:

(i)
$$\alpha > \beta$$
, $\alpha, \beta \in Av(q)$

- (ii) $\alpha \geqslant \beta$, but $\alpha' > \beta'$, where $\alpha = \alpha' t_1$ and $\beta = t_2 \beta'$ for some $t_1, t_2 \in [n-1]$. and one of the following holds:
 - 1. $t_1 > \beta'$, $t_2 < \alpha'$, $t_1 < t_2$ and $\alpha', \beta' \in Av(q)$.
 - 2. $t_1 > \beta'$, $t_2 \not< \alpha'$, $\beta' \in Av(q)$ and $\sigma = \alpha' t_1 t_2 \in \mathbb{H}_1$ with t_2 not being the smallest element in σ and not being the second smallest, after t_1 .
 - 3. $t_1 \not> \beta'$, $t_2 < \alpha'$, $\alpha' \in Av(q)$ and $\sigma = t_1t_2\beta' \in \mathbb{H}_2$ with t_1 not being the biggest element in σ and not being the second biggest, after t_2 .
 - 4. $t_1 \not> \beta'$, $t_2 \not< \alpha'$, $\sigma_1 = \alpha' t_2 \in \mathbb{H}_1$ with t_2 not being the smallest element in σ_1 and $\beta' = x\beta''$, where $x > t_1 > \beta''$ and $\beta'' \in Av(q)$.

The proof of this theorem is not included here. The described decomposition gives us the next result almost directly.

Theorem 4.4.

$$G(x) = 1 + G(x)(xH_1(x) + xH_2(x) + x^3H_1(x)) + G^2(x)(x - 2x^2 - x^3 - x^4).$$

In order to obtain G(x), we found a way to express $H_1(x)$ as a function of $H_2(x)$ and G(x). Then we express $H_2(x)$ as a function of G(x) using the block-decomposition method [32]. Extensive case analysis and inclusion-exclusion arguments are additionally used. As a result, we obtain a system of two equations each of which is a polynomial of x, G(x) and $H_2(x)$. We eliminate $H_2(x)$ to obtain an equation P(x, G(x)) = 0, where P is a polynomial of G(x) with coefficients that are polynomials of x. The polynomial P is given below.

$$x^{8}G(x)^{6} - (x^{8} + 2x^{7} + 5x^{6} + 2x^{5} + 4x^{4} - 4x^{3} + 5x^{2})G(x)^{4} + (2x^{4} + 2x^{3} + 4x^{2} - 2x + 1)G(x)^{2} - 1.$$

One could use a generalization of the Lagrange inversion formula discussed in the work of Banderier and Drmota [4] to get a closed-form expression for the coefficients of G(x).

Recent work by Albert et al. [1] discusses a very general approach automating the discovery of similar decompositions of various sets of combinatorial objects. They claim that their computer method allowed them to find generating functions for the number of permutations avoiding other DPs of size 3 with two squares, namely $1\Box 2\Box 3$, $1\Box^2 32$ and $13\Box^2 2$.

5 Vincular distant patterns

In this section, we consider a particular kind of vincular distant patterns of size 3. The goal will be to establish a surprising fact related to the permutations avoiding the classical patterns 123 and 132.

5.1 Patterns of the form $ab\Box c$ and $a\Box bc$

There are 12 patterns of this kind, and the reverse and complement maps give at most three Wilf-equivalence classes listed below. As we will show, these turn out to be different.

Class 1	Class 2	Class 3
<u>12</u> □3	$1\square \underline{32}$	<u>13</u> □2
$32\square 1$	$\underline{21}\Box 3$	$31\square 2$
$1\square\underline{23}$	$\underline{23}\Box 1$	$2\square \underline{31}$
3□ <u>21</u>	$3\square\underline{12}$	$2\square\underline{13}$
ll .		

Table 1: The Wilf-equivalent classes of $\underline{ab}\Box c$ and $\underline{a}\Box\underline{bc}$ patterns.

We will first find a recurrence for the pattern $12\square 3$:

Theorem 5.1. If $a_n = av_n(\underline{12}\square 3)|$, then $a_n = n!$ for $0 \le n \le 3$, and for $n \ge 4$ we have

$$a_n = a_{n-1} + \sum_{i=2}^{n} \left(\binom{n}{i-1} - 1 \right) a_{n-i}.$$

Proof. Let $q = \underline{12}\square 3$ and let $\pi = \pi_1 \pi_2 \cdots \pi_n$ be a permutation of [n] that avoids q. Also let $n = \pi_i$ and let us consider the two cases when n is part of an occurrence of the classical pattern 123 or not, respectively:

Case 1 n is part of an occurrence of 123.

We must have $\pi_1 > \pi_2 > \cdots > \pi_{i-2}$, as $\pi \in Av_n(q)$. Then, the only way for $\pi_i = n$ to be part of an occurrence of 123 is to have $\pi_{i-2} < \pi_{i-1} < \pi_i$. For the rest of the permutation, we must have $\pi' = \pi_{i+1}\pi_{i+2}\cdots\pi_n \in Av_{n-i}(q)$ and in addition π_{i-1} must be larger then each entry of π' since otherwise, we will have an occurrence of q. Thus to determine π , it suffices to choose the numbers in the set $S_{pref} := \{\pi_1, \pi_2, \dots, \pi_{i-2}\}$ in $\binom{n-1}{i-2}$ ways. Then, we put these numbers in decreasing order at the beginning of the permutation and put π_{i-1} to be the largest among the numbers in $[n-1] \setminus S_{pref}$. The remaining n-i numbers will be the entries of π' . The only forbidden choice for S_{pref} is $S_{pref} = \{n-1, n-2, \dots, n-i+2\}$, because then π_{i-1} will be smaller than all the entries of $\pi_1 \dots \pi_{i-2}$. Hence we obtain $\binom{n-1}{i-2} - 1$ and such permutations π and the position i can be each number in [3, n].

Case 2 n is not part of an occurrence of 123.

Then, $\pi_1 > \pi_2 > \cdots > \pi_{i-1}$ and $\pi' = \pi_{i+1}\pi_{i+2}\cdots\pi_n \in Av_{n-i}(q)$. Therefore, we obtain $\binom{n-1}{i-1}a_{n-i}$ such permutations and the position i can be any number in [1, n].

In total, we have

$$\sum_{i=3}^{n} \left(\binom{n-1}{i-2} - 1 \right) a_{n-i} + \sum_{i=1}^{n} \binom{n-1}{i-1} a_{n-i} = a_{n-1} + \sum_{i=2}^{n} \left(\binom{n}{i-1} - 1 \right) a_{n-i}.$$

The first few elements of the sequence $av_n(\underline{12}\square 3)$ for $n \geq 4$ are

20, 75, 316, 1464, 7359, 39815, 230306.

This is not part of any sequence in the OEIS [34].

This enumerates avoidance for Class 1 patterns in Table 1. Similar recurrence can be found for the patterns in Class 2. We give it below using the pattern $1\square 32$. One can obtain a proof analogous to the proof of Theorem 5.1 by considering the cases when the number 1 is part of an occurrence of the classical pattern 132 or when it is not part of such occurrence.

Theorem 5.2. If $b_n = av_n(1\square \underline{32})$, then $b_n = n!$ for $0 \le n \le 3$ and for $n \ge 4$ we have

$$b_n = b_{n-1} + \sum_{i=2}^{n} \left((i-2) \binom{n-2}{i-2} + \binom{n-1}{i-1} \right) b_{n-i}.$$
 (3)

The first few elements of the sequence $av_n(1\square 32)$ for $n \ge 4$ are

This is not part of any sequence in the OEIS [34].

Applying the complement map after the reverse map, we see that $av_n(\underline{12}\square 3) = av_n(1\square \underline{23})$ for every positive n and we already placed those two patterns in the same of the three classes for the considered set of vincular DPs. Using this, we can easily prove the following.

Theorem 5.3. If n > 4, then $av_n(1 \square 23) < av_n(1 \square 32)$.

Proof. We just noted that $av_n(1\square 23)$ is given by the number a_n from Theorem 5.1, while $av_n(1\square 32)$ is given by the number b_n from Theorem 5.2. By substituting j = n - i + 1, we get that the sum in the right-hand side of Equation (3) can be written as

$$\sum_{j=4}^{n-1} ((n-j+1) \binom{n-2}{j-3} + \binom{n-1}{j-1}) b_{n-j}.$$

Then, in order to obtain this inequality, it suffices to prove that for every n > 4 and $4 \le i \le n - 1$:

$$\binom{n}{i-1} - 1 < (n-i+1)\binom{n-2}{i-3} + \binom{n-1}{i-1}.$$

This is equivalent to $\binom{n-1}{i-2} - 1 < (n-i+1)\binom{n-2}{i-3}$ or $\binom{n-1}{i-2} - 1 < \frac{(n-i+1)(i-2)}{n-1}\binom{n-1}{i-2}$. When n=5 and i=4, we check directly that the latter holds. When n>5, one can easily see that $\frac{(n-i+1)(i-2)}{n-1} > 1$, for $4 \le i \le n-1$.

It remains to investigate Class 3. A well-known proof technique in the area of permutation patterns helps to do that.

Theorem 5.4. For all $n \in \mathbb{Z}^+$, $\operatorname{Av}_n(\underline{13}\square 2) = \operatorname{Av}_n(13\square 2)$, which implies that $av_n(\underline{13}\square 2) = av_n(13\square 2)$.

Proof. We will prove that whenever an n-permutation contains the pattern $13\square 2$, then it must contain the pattern $\underline{13}\square 2$. Take $\sigma = \sigma_1\sigma_2\cdots\sigma_n \in S_n$ containing $q=13\square 2$ and let $\sigma_i\sigma_j\sigma_k$, $1\leq i< j< k-1< n$ be an occurrence of q with the smallest possible distance between the 1 and the 3, i.e., d=j-i is the smallest possible for such an occurrence. If d=1, then $\sigma_i\sigma_j\sigma_k$ would be an occurrence of $\underline{13}\square 2$ and we are done. Assume that d>1 and then consider the value of σ_{i+1} . If $\sigma_{i+1}<\sigma_k$, then $\sigma_{i+1}\sigma_j\sigma_k$ would be a q-occurrence with j-(i+1)=d-1< d. On the other hand, if $\sigma_{i+1}>\sigma_k$, then $\sigma_i\sigma_{i+1}\sigma_k$ would be a q-occurrence with (i+1)-i=1< d, which is again a contradiction.

The theorem that we just proved and the fact that $av_n(12\square 3) = av_n(13\square 2)$ (see [13] and subsection 4.1) imply that $av_n(\underline{13}\square 2)$ is given by the right-hand side of Equation (2) and sequence A049124 in OEIS. It turns out that the patterns in the corresponding Class 3 of Table 1 have the smallest avoiding sets out of the 3 classes.

Theorem 5.5. For all $n \geq 5$, $av_n(\underline{12}\Box 3) > av_n(\underline{13}\Box 2)$.

To establish this fact, we will first need a few additional definitions. For a given pattern q, let $Av_{i_1,i_2,...,i_k;n}(q)$ be the set of q-avoiders of size n beginning with $i_1,i_2,...,i_k$ and let $av_{i_1,i_2,...,i_k;n}(q)$ denote $|Av_{i_1,i_2,...,i_k;n}(q)|$. Let us first prove the following simple lemma.

Lemma 5.6. *If* $1 \le i \le n-1$ *and* $n \ge 4$, *then*

$$av_{i:n}(\underline{13}\Box 2) \le av_{n:n}(\underline{13}\Box 2) = av_{n-1}(\underline{13}\Box 2).$$

Moreover, if $1 \le i \le n-2$, then the inequality is strict, i.e.,

$$av_{i:n}(\underline{13}\Box 2) < av_{n:n}(\underline{13}\Box 2).$$

Proof. For every $\pi \in \operatorname{Av}_{n-1}(\underline{13}\square 2)$, we have that $n\pi \in \operatorname{Av}_{n;n}(\underline{13}\square 2)$, since n cannot participate in any occurrences of $\underline{13}\square 2$, being at first position. Conversely, for every $n\sigma \in \operatorname{Av}_{n;n}(\underline{13}\square 2)$, one has that $\sigma \in \operatorname{Av}_{n-1}(\underline{13}\square 2)$. Thus, $av_{n;n}(\underline{13}\square 2) = av_{n-1}(\underline{13}\square 2)$. In addition, for every $1 \le i \le n-1$ and $i\sigma \in \operatorname{Av}_{i;n}(\underline{13}\square 2)$, we have $fl(\sigma) \in \operatorname{Av}_{n-1}(\underline{13}\square 2)$, which implies that $av_{i;n}(\underline{13}\square 2) \le av_{n-1}(\underline{13}\square 2)$.

Since $n \geq 4$, when $1 \leq i \leq n-2$, then we have at least one *n*-permutation $\pi = in\sigma'$, beginning with i, where $n\sigma' = \sigma$ is such that $fl(\sigma) \in \operatorname{Av}_{n-1}(\underline{13}\square 2)$ and i is obviously part of an $\underline{13}\square 2$ -occurrence. An example is $\pi = ina \cdots (n-1)$ for any $a \in [n]$, where $a \neq i, n, (n-1)$. Therefore, $fl(\sigma) \in \operatorname{Av}_{n-1}(\underline{13}\square 2)$, but $\pi = i\sigma \notin \operatorname{Av}_{i,n}(\underline{13}\square 2)$.

We will need this lemma together with a few other definitions. Given a permutation (pattern) σ , let $C_n(\sigma) = S_n \setminus \operatorname{Av}_n(\sigma)$ be the permutations of S_n containing σ . Then, let

$$U_n := C_n(\underline{12}\square 3) \cap \operatorname{Av}_n(\underline{13}\square 2)$$

and

$$V_n := \operatorname{Av}_n(\underline{12}\square 3) \cap C_n(\underline{13}\square 2),$$

with $u_n := |U_n|$ and $v_n := |V_n|$. In addition, let us denote with $U_{i_1,i_2,\dots,i_k;n}$ (respectively, $V_{i_1,i_2,\dots,i_k;n}$) the set of permutations in U_n (respectively, in V_n) beginning with $i_1i_2 \dots i_k$. Furthermore, let $u_{i_1,i_2,\dots,i_k;n} := |U_{i_1,i_2,\dots,i_k;n}|$ and $v_{i_1,i_2,\dots,i_k;n} := |V_{i_1,i_2,\dots,i_k;n}|$. Now, we will prove the following.

Lemma 5.7. For each $n \ge 4$ and $1 \le i \le n$,

$$u_{i;n} \leq v_{i;n}$$
.

Proof. Note that the statement implies $u_n \leq v_n$ and $|C_n(\underline{12}\square 3)| \leq |C_n(\underline{13}\square 2)|$ (respectively, $av_n(\underline{12}\square 3) \ge av_n(\underline{13}\square 2)$, for each $n \ge 4$. Indeed, if $T_n = C_n(\underline{12}\square 3) \cap$ $C_n(\underline{13}\square 2)$, then $C_n(\underline{12}\square 3) = U_n \cup T_n$ and $C_n(\underline{13}\square 2) = V_n \cup T_n$. Thus, $u_n \leq v_n$ implies $|C_n(\underline{12}\square 3)| \leq |C_n(\underline{13}\square 2)|$. We will proceed by induction on n. One can directly check that $u_{i;4} \leq v_{i;4}$ for each $1 \leq i \leq 4$. Now assume that $u_{i;n'} \leq v_{i;n'}$, for each $4 \le n' \le n-1$ and $1 \le i \le n'$, for a given $n \ge 5$. Consider $u_{i,n}$ and $v_{i,n}$ for $1 \le i \le n$. If i = n, then using the induction hypothesis, we have $u_{n,n} = u_{n-1} \leq v_{n-1} = v_{n,n}$. Similarly, if i = n - 1, then we have $u_{n-1,n} = u_{n-1} \leq v_{n-1} = v_{n-1,n}$. Now, let $1 \le i \le n-2$. By the induction hypothesis, $u_{i,i-k,n} = u_{i-k,n-1} \le v_{i-k,n-1} = v_{i,i-k,n}$, for each $1 \leq k \leq i-1$. It remains to compare the numbers $u_{i,i+k,n}$ and $v_{i,i+k,n}$ for $1 \leq k \leq n-i$. Note that when $k \geq 3$, then $u_{i,i+k,n} = 0$, since for these values of k, any n-permutation beginning with i(i+k) will contain an occurrence of <u>13</u> \square 2. Similarly, $v_{i,i+k;n} = 0$, when i + k < n - 1, since for these values of k, any *n*-permutation beginning with i(i+k) will contain an occurrence of 12 \square 3. We will show that $u_{i,i+1,n} \leq v_{i,n,n}$ and that $u_{i,i+2,n} \leq v_{i,n-1,n}$, which will complete the proof. Consider the sets $U_{i,i+1,n}$ and $V_{i,n,n}$. First, let us look at those $\pi \in U_{i,i+1,n}$ (respectively, $\pi \in V_{i,n,n}$ that do not begin with an $\underline{12}\square 3$ occurrence (respectively, not with an <u>13</u> \square 2 occurrence). Then, note that π must begin with (n-2)(n-1)n (respectively with (n-2)n(n-1)). However, we have

$$u_{n-2,n-1,n;n} = u_{n-3} \le v_{n-3} = v_{n-2,n;n-1} \tag{4}$$

using the induction hypothesis, again.

Now, let us look at those $\pi \in U_{i,i+1,n}$ beginning with a $\underline{12}\square 3$ occurrence. Their number is given by

$$av_{i:n-1}(\underline{13}\square 2) - av_{n-2,n-1:n-1}(\underline{13}\square 2).$$
 (5)

Indeed, after we remove from π its first element i and standardize π , we obtain an (n-1)-avoider of $\underline{13}\square 2$. Conversely, for any permutation $\pi=i\pi_2\dots\pi_{n-1}\in \mathrm{Av}_{i;n-1}(\underline{13}\square 2)$, one can increase by 1 all the elements of π greater than or equal to i and then add i at the beginning, to obtain a permutation in $U_{i,i+1;n}$. This permutation will begin with a $\underline{12}\square 3$ occurrence, unless it begins with (n-2)(n-1)n, i.e., when i=n-2 and when $\pi\in \mathrm{Av}_{n-2,n-1;n-1}(\underline{13}\square 2)$. Therefore, we should subtract $av_{n-2,n-1;n-1}(\underline{13}\square 2)$. Respectively, for the number of permutations $\pi\in V_{i,n;n}$, beginning with an $\underline{13}\square 2$ occurrence, one would have

$$av_{n-1;n-1}(\underline{12}\square 3) - av_{n-1,n-2;n-1}(\underline{12}\square 3).$$
 (6)

It is not difficult to see that $av_{n-2,n-1;n-1}(\underline{13}\square 2) = av_{n-3}(\underline{13}\square 2)$ and that $av_{n-1,n-2;n-1}(\underline{12}\square 3) = av_{n-3}(\underline{12}\square 3)$. Hence, by using expressions (5), (6) and Equation (4), we see that in order to establish that $u_{i,i+1;n} \leq v_{i,n;n}$, it remains to prove the inequality below for each $1 \leq i \leq n-2$:

$$av_{i:n-1}(\underline{13}\square 2) - av_{n-3}(\underline{13}\square 2) \le av_{n-1:n-1}(\underline{12}\square 3) - av_{n-3}(\underline{12}\square 3).$$
 (7)

By Lemma 5.6, we have that $av_{i;n-1}(\underline{13}\Box 2) \leq av_{n-1;n-1}(\underline{13}\Box 2) = av_{n-2}(\underline{13}\Box 2)$. We also have that $av_{n-1;n-1}(\underline{12}\Box 3) = av_{n-2}(\underline{12}\Box 3)$. Thus, it suffices to prove that

$$av_{n-2}(\underline{13}\square 2) - av_{n-3}(\underline{13}\square 2) \le av_{n-2}(\underline{12}\square 3) - av_{n-3}(\underline{12}\square 3).$$
 (8)

Using that $av_{n-3}(q) = av_{n-2;n-2}(q)$ for any of the patterns $q = \underline{12}\square 3$ or $q = \underline{13}\square 2$, as well as the relation

$$av_{i:n-2}(\underline{13}\square 2) \le av_{i:n-2}(\underline{12}\square 3) \iff v_{i:n-2} \ge u_{i:n-2},$$

we see that Equation (8) is equivalent to

$$\sum_{i=1}^{n-3} v_{i;n-2} \ge \sum_{i=1}^{n-3} u_{i;n-2},\tag{9}$$

which follows directly, because by the induction hypothesis $u_{i;n-2} \leq v_{i;n-2}$, for all $i \in [1, n-3]$. From Equations (4) and (9), we conclude that $u_{i,i+1;n} \leq v_{i,n;n}$.

One could establish that $u_{i,i+2;n} \leq v_{i,n-1;n}$ in almost the same way, by first noticing that $U_{i,i+2;n} = U_{i,i+2,i+1;n}$ and that $V_{i,n-1;n} = V_{i,n-1,n;n}$ since the permutations in $U_{i,i+2;n}$ (respectively, in $V_{i,n-1;n}$) do not have a $\underline{13}\square 2$ (respectively, a $\underline{12}\square 3$) occurrence. Then, the only thing that remains is to consider the cases i = n-2 and $i \neq n-2$ and to use the induction hypothesis and Lemma 5.6. In particular, if i = n-2, then

$$u_{n-2,n,n-1;n} = u_{n-3} \le v_{n-3} = v_{n-2,n-1,n;n}.$$

If $i \neq n-2$, then $\pi \in U_{i,i+2,i+1,n}$ (respectively in $V_{i,n-1,n,n}$) begins with an $\underline{12}\square 3$ (respectively an $\underline{13}\square 2$) occurrence and

$$u_{i,i+2,i+1;n} = av_{i,n-2}(\underline{13}\square 2) \le av_{n-2;n-2}(\underline{13}\square 2)$$

by Lemma 5.6. In addition,

$$av_{n-2;n-2}(\underline{13}\square 2) \le av_{n-2;n-2}(\underline{12}\square 3) = v_{i,n-1,n;n}$$

by the induction hypothesis.

As we have pointed out, Lemma 5.7 implies that $av_n(\underline{12}\square 3) \geq av_n(\underline{13}\square 2)$, for each $n \geq 4$. In order to obtain a proof of Theorem 5.5, we should just use the second part of Lemma 5.6 to see that Inequality (7) is strict when $n-1 \geq 4$, i.e., when $n \geq 5$.

We are now ready to formulate an interesting conclusion. The last Theorem 5.5 together with Theorem 5.3, the result on consecutive patterns of Elizalde [10] and the corollary of the result of Hopkins (Theorem 4.1) imply the following:

Corollary 5.8. Consider the set of distant patterns

$$X = \{1 \square 23, 12 \square 3, 1 \square 2 \square 3, 123\}.$$

Take any pattern $p \in X$ and switch the places of the letters 2 and 3 to get a pattern p' in $Y = \{1 \square 32, 13 \square 2, 1 \square 3 \square 2, 132\}$. We have that $Av_n(p) > Av_n(p')$ for all n > 5, $p \in X$ and the corresponding $p' \in Y$, except for $1 \square 23$ that is avoided by fewer permutations of size n, compared to its counterpart $1 \square 32$.

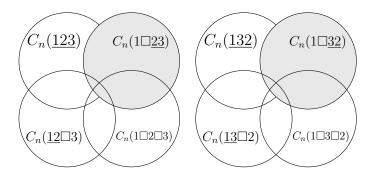


Figure 1: Diagrams for the *n*-permutations containing 123 and 132

Figure 1 depicts the sets of permutations containing each of the patterns in X and Y. Corollary 5.8 is somewhat surprising since each occurrence of the classical pattern 123 (respectively, 132) is an occurrence of a pattern in X (respectively, Y) and as it was shown in [36], $av(123) = av_n(132)$. Thus the total "area" of the union of the four sets on the left is the same as the total "area" of the union of the four sets on the right. However, each of the three unmarked sets on the left contains fewer elements than its counterpart on the right.

5.2 Consecutive distant patterns

Recall that when all the constraints for the gap sizes in a distant pattern are tight, then we call these patterns consecutive distant patterns and we underline the whole pattern to denote that. Considering POGP, Kitaev mentioned in the introduction of [23] that $\operatorname{Av}_n(\underline{1\square 2}) = \binom{n}{\lfloor \frac{n}{2} \rfloor}$. Indeed, we may use that the letters in the odd and the even positions of a permutation avoiding this pattern do not affect each other. Thus we can choose the letters in odd positions in $\binom{n}{\lfloor \frac{n}{2} \rfloor}$ ways, and we must arrange them in decreasing order. We then must arrange the letters in even positions in decreasing order, too. Using the same reasoning one can easily find, for example $\operatorname{Av}_n(\underline{1\square 2}\underline{2})$ or $\operatorname{Av}_n(\underline{1\square 2}\underline{\square 3})$. This can be further generalized by the fact given below. Recall that if $q = q_1q_2\cdots q_k$ is a classical pattern of size k, then $\underline{q} = \underline{q_1q_2\cdots q_k}$ is the corresponding consecutive pattern. We also use $\underline{\operatorname{dist}_r(q)}$ to denote the corresponding consecutive distant pattern $q_1\square^r q_2\square^r \cdots \square^r q_k$.

Theorem 5.9. [21, Theorem 11] For a given classical pattern q of size k, given distance $r \ge 0$ and a natural n, denote $l = \lfloor \frac{n}{r+1} \rfloor$. Set $u := n \mod (r+1) \in [0,r]$. Then

$$av_{n}(\underline{\text{dist}_{r}(q)}) = \frac{n!}{(l!)^{r+1-u}((l+1)!)^{u}} (av_{l}(\underline{q}))^{r+1-u} (av_{l+1}(q))^{u}.$$
 (10)

This gives us a formula for the size of the set of permutations avoiding any consecutive distant pattern, knowing the size of the avoidance set for the corresponding classical consecutive pattern. Corollaries of this simple fact were previously stated in [13, 14]. We state another simple corollary here, which shows a surprising relationship between the former fact and avoidance of arithmetic progressions in permutations.

Theorem 5.10. The number of permutations of size n avoiding arithmetic progressions of length k > 1 and difference r > 0, i.e., the number of permutations π without subsequences $\pi_{i_1}\pi_{i_2}\cdots\pi_{i_k}$, such that $\pi_{i_1} = x$, $\pi_{i_2} = x + r$, \cdots , and $\pi_{i_k} = x + (k-1)r$ for some $x, r \in [n]$, is $av_n(\operatorname{dist}_r(12\cdots k))$, which can be obtained using Equation (10).

Proof. Consider $\pi \in S_n$, containing an arithmetic progression $\pi_{i_1}\pi_{i_2}\cdots\pi_{i_k}$ of size k and difference r>0. Then in the inverse permutation π^{-1} , $i_1i_2\cdots i_k$ will be an occurrence of the distant pattern $\underline{\operatorname{dist}_r(12\cdots k)}$ since $\pi^{-1}(x)=i_1, \pi^{-1}(x+r)=i_2, \cdots, \pi^{-1}(x+(k-1)r)=i_k$. Conversely, if $\pi \in S_n$ contains $\underline{\operatorname{dist}_r(12\cdots k)}$, then π^{-1} contains an arithmetic progression of length k and difference r>0. Therefore, the number of permutations of [n] containing arithmetic progressions of length k and difference r>0 equals the number of permutations of [n] containing $\underline{\operatorname{dist}_r(12\cdots k)}$. This implies the same for the set of avoiders, i.e., what we aim to prove.

6 Interpretations of other results

Here, we will demonstrate that DPs can be very useful when interpreting already known results (including ones obtained with a computer). One previous work that gives several conjectures about the enumeration of pattern-avoiding classes containing many size four patterns is the work of Kuszmaul [28]. He listed ten conjectures about simultaneous pattern avoidance of many size four patterns. One can find brief solutions, using both computer programs and manual work, to many of these conjectures in the two articles of Mansour and Schork [30, 31].

Below, we give direct solutions to two of the conjectures without using a computer. To do that, we interpret the respective big set of size four patterns as a smaller set of both classical and distant patterns. Our approach is similar to the technique introduced in [29].

Theorem 6.1. (conjectured in [28, p.20, sequence 6]) The generating function of $av_n(2431, 2143, 3142, 4132, 1432, 1342, 1324, 1423, 1243)$

is given by $C + x^3C$, where C is the generating function for the Catalan numbers.

Proof. Note that the set of patterns above can be written as

$$\Pi = \{ \Box 132, 132 \Box, 1342 \}$$

When n=1,2,3, $av_n(\Pi)=1,2,6$, respectively, and these are indeed the first three coefficients of $C+x^3C$. Consider values $n\geq 4$. If $\sigma\in \operatorname{Av}_n(132)$, then $\sigma\in \operatorname{Av}_n(\Pi)$. As we know, $av_n(132)$ has generating function C [27]. It remains to find the generating function for those σ containing 132, but avoiding Π . Take one such $\sigma=\sigma_1\sigma_2\cdots\sigma_n$ and notice that any occurrence of 132 in σ must have σ_1 as its first letter and σ_n as its last letter. Otherwise, given that $n\geq 4$, an occurrence of at least one of the patterns 132 \square or \square 132 will be present. Now, let $\sigma_k=n$ be the biggest element of σ . Clearly, $\sigma_1\sigma_k\sigma_n$ must be an occurrence of 132. If not, then this biggest element must be either at the first or the last position in σ and thus σ would not contain any 132-occurrences that either begin at σ_1 or end at σ_n . In Figure 2 are shown three black points representing σ_1 , σ_k and σ_n , as well as three segments of the diagram of σ denoted with A, B and C and defined below. We further show that σ will not contain any elements in these three segments. Here is why:

- A is empty There is no element x among $\sigma_2, \sigma_3, \ldots, \sigma_{k-1}$, such that $x < \sigma_n$. If there is such x, then $x\sigma_k\sigma_n$ would be a forbidden occurrence of 132.
- B is empty There is no element x among $\sigma_{k+1}, \sigma_{k+2}, \dots, \sigma_{n-1}$, such that $\sigma_1 < x < \sigma_k$. If there is such x, then $\sigma_1 \sigma_k x$ would be a forbidden occurrence of 132.
- C is empty There is no element x among $\sigma_1, \sigma_2, \ldots, \sigma_{k-1}$, such that $\sigma_n < x < \sigma_k$. If there is such x, then $\sigma_1 x \sigma_k \sigma_n$ would be an occurrence of 1342.

Therefore, the biggest element σ_k in σ must be at position 2, i.e., k=2 and the only non-empty segment could be the one denoted by α in figure 2. In other words, we must have $\sigma = (n-2)n\alpha(n-1)$, for some permutation $\alpha \in \text{Av}(132)$. Otherwise, an occurrence of 132, such that $\sigma_1 = n-2$ is not part of it, would be formed. Conversely, for any $\alpha \in \text{Av}(132)$, $\sigma = (n-2)n\alpha(n-1)$ belongs to $\text{Av}(\Pi)$.

Then, we get x^3C for the generating function of the permutations in $Av(\Pi)$ containing 132 and therefore we will have $C + x^3C$ for the generating function of $Av(\Pi)$, since C is the generating function for Av(132).

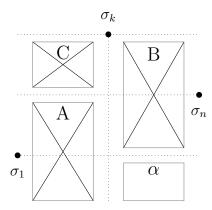


Figure 2: Decomposition for $\sigma \in \text{Av}(\Pi)$ from the proof of Theorem 6.1

As we know, $C = 1 + xC^2$, so we can write

$$C + x^{3}C = C + x^{3}(1 + xC^{2}) = x^{3} + C \cdot (1 + x^{4}C)$$

and this indeed corresponds to sequence [34, A071742] given by $C \cdot (1 + x^4C)$, as reported in [28], with the subtle difference that for n = 3, we have one extra permutation in $Av_3(\Pi)$, namely 132. The same structure for the decomposition of the permutations in $Av(\Pi)$ was also found recently with a computer by Bean et al., who used a particular algorithm called the *Struct algorithm* [5]. As we saw, rewriting the problem in terms of distant patterns helped us to prove the result directly and to give an interpretation of the already discovered decomposition.

Below, we will give a direct proof of another former conjecture listed in [28].

Theorem 6.2. (conjectured in [28, p.19, sequence 5]) The generating function of

$$av_n(2431, 2413, 3142, 4132, 1432, 1342, 1324, 1423)$$

is given by $C \cdot (1+x^3C)$, where C is the generating function for the Catalan numbers.

Proof. Note that the set of permutations above can be written as

$$\Pi = \{13\square 2, 1324, 2431, 3142, 4132\}.$$

If σ has no occurrences of 132 at all, then obviously $\sigma \in \operatorname{Av}(\Pi)$ and the generating function for these permutations is C. Let us consider those σ that have some occurrences of 132 and are in $\operatorname{Av}(\Pi)$. The set Π contains 13 \square 2 thus all the occurrences of 132 in σ , are occurrences of 132. Take the occurrence $\sigma_i \sigma_j \sigma_{j+1}$ of 132 that ends at the largest possible position, i.e., with j maximal. Denote by α the segment in σ of largest possible size that ends at σ_i and such that $\alpha < \sigma_{j+1} < \sigma_j$. Let us first consider $\sigma' = \sigma_{j+2}\sigma_{j+3}\cdots\sigma_n$. We will show that σ' is the empty permutation, i.e., n = j + 1 and the segments A,B and C, defined below and shown at Figure 3 are empty.

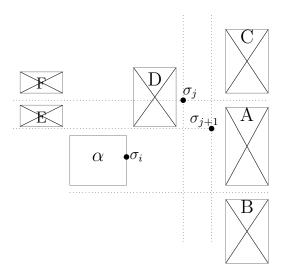


Figure 3: Decomposition for $\sigma \in Av(\Pi)$ from the proof of Theorem 6.2.

\bullet A is empty

There is no element x in σ' , such that $x < \sigma_j$ and $x \not< \alpha$. If there is such x, then $\sigma_i \sigma_j x$ would be an occurrence of 132 that is not an 132-occurrence.

\bullet B is empty

There is no element x in σ' , such that $x < \alpha$. If there is such x, then $\sigma_i \sigma_j \sigma_{j+1} x$ would be an occurrence of 2431, which is not allowed.

\bullet C is empty

There is no element x in σ' , such that $x > \sigma_j$. If there is such x, then $\sigma_i \sigma_j \sigma_{j+1} x$ would be an occurrence of 1324, which is not allowed.

Next, let us consider the segment $\sigma'' = \sigma_{i+1} \cdots \sigma_{j-1}$. We will show that σ'' is the empty permutation, i.e., the segment D, shown at Figure 3, is empty and i = j - 1.

• D is empty

There is no element x in σ'' , such that $x > \sigma_{j+1}$. If there is such x, then $\sigma_i x \sigma_{j+1}$ would be an occurrence of 132 that is not an 132-occurrence.

Finally, consider the segment σ''' that is the part of σ in front of α . We will show that σ''' is the empty permutation, i.e., the segments E and F, shown at Figure 3 are empty.

• E is empty

There is no element x in σ''' , such that $x > \sigma_{j+1}$ and $x < \sigma_j$. If there is such x, then $x\sigma_i\sigma_j\sigma_{j+1}$ would be an occurrence of 3142, which is forbidden.

• F is empty

There is no element x in σ''' , such that $x > \sigma_j$. If there is such x, then $x\sigma_i\sigma_j\sigma_{j+1}$ would be an occurrence of 4132, which is forbidden.

Therefore, we must have $\sigma = \alpha n(n-1)$, where α is non-empty and $\alpha \in \text{Av}(132)$. The latter holds since if α contains 132 then after appending σ_j at the end, we will get an occurrence of 1324. Conversely, one may readily check that for each non-empty $\alpha \in \text{Av}(132)$, $\sigma = \alpha n(n-1)$ would be a permutation in $\text{Av}(\Pi)$. Thus, the generating function of the number of permutations in $\text{Av}(\Pi)$ is $C + x^2(C-1)$, since the generating function of such non-empty $\alpha \in \text{Av}(132)$ is C-1. Furthermore, we have

$$C + x^{2}(C - 1) = C + x^{2}xC^{2} = C \cdot (1 + x^{3}C).$$

The sequence of the coefficients for this generating function is given by A071726 in OEIS.

7 Stanley-Wilf type conjectures for DPs

A popular former conjecture on the classical permutation patterns formulated independently by Stanley and Wilf states that for any given classical pattern q, there exists a constant c_q , such that $\sqrt[n]{av_n(q)} < c_q$, when $n \to \infty$. In 1999, Arratia ([2]) observed that this is equivalent to the existence of the limit $\lim_{n\to\infty} \sqrt[n]{av_n(q)}$. The conjecture was resolved in 2004 by Marcus and Tardos ([33]) who actually proved a conjecture of Füredi and Hajnal, which had been shown earlier to imply the Stanley-Wilf conjecture.

In Theorem 2.1, we saw that the avoidance of every distant pattern is equivalent to simultaneous avoidance of several classical patterns. The Stanley-Wilf conjecture is true for any of these classical patterns. Thus we will have that $\sqrt[n]{av_n(q)} < const$, for any distant pattern q, when $n \to \infty$. Arratia's observation that $av_{m+n}(q) \ge av_m(q).av_n(q)$, also holds for distant patterns, if the considered distant pattern does not start with a square. Thus for those kind of distant patterns, we can rely on Fekete's lemma on subadditive sequences, exactly as Arratia did, to obtain that $\sqrt[n]{av_n(dp)}$ exists. As for the DPs beginning with r > 0 number of squares, we can use Theorem 2.2 to write

$$\sqrt[n]{av_n(\Box^r q)} = \sqrt[n]{n^{(r)}av_{n-r}(q)} = (n^{(r)})^{\frac{1}{n}}av_{n-r}(q)^{\frac{1}{n-r}\frac{n-r}{n}} \underset{n \to \infty}{\longrightarrow} c_q,$$

where q is a distant pattern that does not start with a square and $\sqrt[n]{av_n(q)} \longrightarrow c_q$. This yields the following Stanley-Wilf type result for DPs.

Theorem 7.1. For any distant pattern q, there exists a constant $c_q > 0$, such that

$$\sqrt[n]{av_n(q)} \underset{n \to \infty}{\longrightarrow} c_q.$$

An interesting continuation might be to consider avoidance of $\operatorname{dist}_r(q)$, for a classical pattern q and size of r that increases with n. Obviously, if $r \geq n - 1$, then

 $av_n(\operatorname{dist}_r(q)) = n!$ for any pattern q. However, one may ask what will happen if r is a fixed positive fraction of n. Is it still true in this case that the corresponding number of n-permutations is bounded above by c^n for some constant c? That is the question that is answered in the negative with Theorem 3.2 below.

Theorem 7.2. It is not true that for any given classical pattern q, there exist constants c > 0 and 0 < t < 1, such that

$$\sqrt[n]{av_n(\operatorname{dist}_r(q))} \overset{n \to \infty}{\underset{r = \lfloor tn \rfloor}{\longrightarrow}} c.$$

Proof. Consider the classical pattern q=12. By Theorem 3.2, the number of n-permutations avoiding $1\square^r 2$, for any $r\geq 1$, will be the same as the number of n-permutations for which in each of their cycles, any two elements differ by at most r. Denote this set of permutations by S_n^r . Furthermore, let C_n^r be the set of permutations in S_n for which each cycle is of length exactly r, except possibly 1 cycle of smaller length, if r does not divide n, and where each cycle consists of consecutive elements. Therefore, since $C_n^r \subset S_n^r$, we can use that $|C_n^r| < |S_n^r| = av_n(1\square^r 2)$. In addition, we have the obvious bound $|C_n^r| \geq ((r-1)!)^{\lfloor \frac{n}{r} \rfloor}$. Thus for any given 0 < t < 1, if r = |tn|, then for big enough values of n, we have:

$$av_n(1\square^r 2) > ((r-1)!)^{\lfloor \frac{n}{r} \rfloor} > ((\lfloor tn-2 \rfloor)!)^{\frac{1}{t}-1} \ge \left(\left(\frac{\lfloor tn-2 \rfloor}{e} \right)^{\lfloor tn-2 \rfloor} \right)^{\frac{1}{t}-1} > \left(\frac{tn-3}{e} \right)^{(1-t)n-\frac{3}{t}+3},$$

where the last inequality follows from the Stirling approximation. It is easy to check that

$$\left(\frac{tn-3}{e}\right)^{(1-t)n-\frac{3}{t}+3} = \Omega(D^n)$$

for any constant D > 0. One can do that by taking logarithms of both sides. \square

The latter fact motivates us to consider the following conjecture.

Conjecture 7.3. For any given classical pattern q, such that $|q| \ge 2$, and for every 0 < t < 1, there exists 0 < w < 1, such that:

$$\lim_{n \to \infty} \left(\frac{av_n(\operatorname{dist}_{\lfloor tn \rfloor}(q))}{n!} \right)^{\frac{1}{n}} = w.$$

The approach of Elizalde ([11, Section 4]) for consecutive patterns might be useful when one tries to prove the latter conjecture, even though this approach cannot be applied directly. Here, we prove one lemma that might help confirming the conjecture.

Lemma 7.4. For any given classical pattern $q \in S_k$ and for every 0 < t < 1, there exists d < 1, such that if $r = \lfloor tn \rfloor$ and $n \ge k(r+1)$, then

$$av_n(\operatorname{dist}_r(q)) < d^n n!.$$

Proof. Assume that $n \geq k(r+1)$ for some t and n and let us take an arbitrary permutation $\pi = \pi_1 \pi_2 \cdots \pi_n$. Consider the following set of subsequences of π :

$$\mathbb{F} := \{ \pi_1 \pi_{r+2} \cdots \pi_{(k-1)r+k}, \pi_2 \pi_{r+3} \pi_{2r+4} \cdots \pi_{(k-1)r+k+1}, \dots, \pi_{r+1} \pi_{2r+2} \cdots \pi_{(r+1)k} \},$$

This is a set of r+1 subsequences of π ; Each of them is of size k and each entry of the prefix of π , $\pi_1\pi_2\cdots\pi_{(r+1)k}$ is part of exactly one subsequence in \mathbb{F} . At the same time, if $\pi \in Av_n(dist_r(q))$, then none of the subsequences in \mathbb{F} is order-isomorphic to q. Denote the event that when sampling a permutation $\pi \in S_n$ uniformly at random, no subsequence in \mathbb{F} is order isomorphic to q by $E_{\mathbb{F}}$. Since $|\mathbb{F}| = r + 1$, we will have

$$\mathbb{P}(E_{\mathbb{F}}) = \left(1 - \frac{1}{k!}\right)^{r+1}$$

Therefore, if we write $E_{\text{dist}_r(q)}$ for the event that π contains the pattern $\text{dist}_r(q)$, then

$$\mathbb{P}(E_{\text{dist}_r(q)}) > 1 - \left(1 - \frac{1}{k!}\right)^{r+1}$$

and thus the number of permutations π containing $\operatorname{dist}_r(q)$ is at least $n!(1-(1-\frac{1}{k!})^{r+1})$. Hence,

$$av_n(\operatorname{dist}_r(q)) \leq n! \left(\left(1 - \frac{1}{k!} \right)^{r+1} \right) = \left(\left(1 - \frac{1}{k!} \right)^{\frac{r+1}{n}} \right)^n n!$$

$$< \left(\left(\left(1 - \frac{1}{k!} \right)^t \right)^n n! = d^n n!.$$

The latter inequality holds since $\frac{r+1}{n} = \frac{\lfloor tn \rfloor + 1}{n} > t$.

An analogous fact could be conjectured about the bound from below.

Conjecture 7.5. For any given classical pattern $q \in S_k$ and for every 0 < t < 1, there exists c > 0, such that if r = |tn| and $n \ge k(r+1)$, then

$$av_n(\operatorname{dist}_r(q)) > c^n n!.$$

We saw that when r is a positive fraction of n, the number of n-permutations avoiding the corresponding distant pattern may become huge. Thus it would be reasonable to consider a Stanley-Wilf type conjecture, where r is asymptotically smaller than $\mathcal{O}(n)$, e.g., a function of the kind n^t , for 0 < t < 1.

Conjecture 7.6. For any given classical pattern q, there exist $c_q > 0$ and 0 < t < 1, such that

$$\sqrt[n]{av_n(\operatorname{dist}_r(q))} \xrightarrow[r=|n^t|]{n\to\infty} c_q.$$

If the latter conjecture is true, then one might ask which are the allowable growth rates c_q when c_2 is a fixed positive constant. Furthermore, an interesting additional question would be to find a function g(n), such that

$$\sqrt[n]{av_n(\operatorname{dist}_r(q))} \xrightarrow[r=\lfloor\Theta(g(n))\rfloor]{n\to\infty} c,$$

for some constant c > 0, but

$$\sqrt[n]{av_n(\operatorname{dist}_r(q))} \xrightarrow[r=|\Omega(g(n))|]{n\to\infty} c,$$

for any c > 0.

8 Open problems and future work

In this section, we list some ideas for possible further investigations, related to the current work on distant patterns:

- The approach using the map g in Section 3 can be applied to the set $A'_{n-1} = S_{n-1} \setminus \operatorname{Av}_n(1 \square 23) \times [n]$ with the hope to find an equation that will give us a recurrence formula for $av_n(1 \square 2 \square 3)$. We have checked that a statement analogous to Theorem 3.4 holds and that no permutation is the image of g for more than two different elements of A'_{n-1} .
- One can try to prove or refute the following surprising conjecture (see Section 4.1 for a discussion):

Conjecture 8.1. Choose one of the 3 places between consecutive letters in the patterns {1234, 1243, 2143} and put a square at that place for each of the three given classical patterns. You will obtain three Wilf-equivalent distant patterns. For example

$$av_n(1\square 234) = av_n(1\square 243) = av_n(2\square 143).$$

We should note that a similar statement does not hold for any two Wilf-equivalent classical patterns, because $av_n(4132) = av_n(3142)$ [37], for all n > 1, but $av_7(4\square 132) = 3592 \neq 3587 = av_7(3\square 142)$.

• Three more conjectures related to the least and most avoided uniform distant patterns can be investigated (see Section 4 for a discussion):

Conjecture 8.2. For every $m \geq 3$ and $r \geq 1$, there exists $n_0 \in \mathbb{N}$ such that for every natural $n > n_0$, we have

$$av_n(\operatorname{dist}_r(12\cdots m)) > av_n(\operatorname{dist}_r(q)),$$

for any given classical pattern q of size m.

Conjecture 8.3. For every $m \geq 3$ and $r \geq 1$, there exists $n_0 \in \mathbb{N}$ such that for every natural $n > n_0$, we have

$$av_n(\operatorname{dist}_r(12\cdots(m-2)m(m-1))) \le av_n(\operatorname{dist}_r(q)),$$

for any given classical pattern q of size m.

A weaker version of these two conjectures would be the one below and a suitable injection establishing the fact is desired.

Conjecture 8.4. For every $m \geq 3$ and $r \geq 1$, there exists $n_0 \in \mathbb{N}$ such that for every natural $n > n_0$:

$$av_n(\operatorname{dist}_r(12\cdots m)) \ge av_n(\operatorname{dist}_r(12\cdots (m-2)m(m-1))).$$

- The PermPAL database [35] currently contains about 17000 permutation pattern avoidance results and one may try to find equivalent interpretations in terms of distant patterns, as well as new combinatorial proofs using those interpretations (similar to the two proofs in Section 6). One may also use the database [9].
- Section 7 mentions a few unresolved conjectures, namely Conjecture 7.3, 7.5 and 7.6, all related to Stanley–Wilf type results.
- One may investigate the case of bivincular distant patterns, i.e., when we have constraints for the values of the letters in an occurrence of a distant pattern in addition to gap size constraints.
- Several statistics over permutations avoiding certain distant patterns might be considered and perhaps some facts related to their distribution could be established. We were not able to find much previous work related to the topic.

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