Impartial achievement games for generating generalized dihedral groups

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Abstract

We study an impartial game introduced by Anderson and Harary. This game is played by two players who alternately choose previously-unselected elements of a finite group. The first player who builds a generating set from the jointly-selected elements wins. We determine the nimnumbers of this game for generalized dihedral groups, which are of the form $\text{Dih}(A) = \mathbb{Z}_2 \ltimes A$ for a finite abelian group A.

1 Introduction

Anderson and Harary [2] introduced a pair of two-player impartial games on a nontrivial finite group G. In both games, two players take turns choosing a previouslyunselected element of G. The player who first builds a generating set from the jointly-selected elements wins the achievement game GEN(G), while a player who cannot avoid building a generating set loses the avoidance game DNG(G). The theory of these games is further developed in [3, 4, 5, 7]. Similar algebraic games are studied in [6].

For a finite abelian group A, the generalized dihedral group Dih(A) is the semidirect product $\mathbb{Z}_2 \ltimes A$. Generalized dihedral groups share many of the same properties

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of dihedral groups. For example, every element of Dih(A) that is not in A has order 2 and acts on A by inversion. See [8] or [9] for an introduction to generalized dihedral groups.

A fundamental problem in game theory is to determine nim-numbers of impartial two-player games. The nim-number allows for the easy calculation of the outcome of the sum of games. A general theory on how to compute nim-numbers appears in [7]. The strategies for the avoidance game for many families of finite groups were presented in [3], and a complete theory for finding the nim-numbers for the avoidance games was developed in [5]. The strategies and nim-numbers for symmetric and alternating groups were determined in [4] for both the achievement and avoidance games.

The nim-numbers for $\mathsf{DNG}(\mathsf{Dih}(A))$ were classified in [5], which found that $\mathsf{DNG}(\mathsf{Dih}(A))$ is *3 if A is cyclic of odd order and *0 otherwise. The task in this paper is to determine the nim-numbers for $\mathsf{GEN}(\mathsf{Dih}(A))$ for all finite abelian A.

2 Preliminaries

In this section, we recall some general terminology and results from [4, 5, 7] related to impartial games, as well as the achievement game that is the focus of this paper.

2.1 Impartial Games

A comprehensive treatment of impartial games can be found in [1, 10]. An *impartial* game is a finite set X of positions together with a collection $\{Opt(P) \subseteq X \mid P \in X\}$, where Opt(P) is the set of possible options for a position P. Two players take turns replacing the current position P with one of the available options in Opt(P). The player who encounters a *terminal position* with an empty option set cannot move and therefore *loses*. All games must come to an end in finitely many turns, so we do not allow infinite lines of play.

The one-pile NIM game *n with n stones is the prototype of an impartial game. The set of options of *n is $Opt(*n) = \{*0, \ldots, *(n-1)\}$.

The minimum excludant mex(A) of a set A of ordinals is the smallest ordinal not contained in the set. The nim-number nim(P) of a position P is the minimum excludant of the set of nim-numbers of the options of P. That is,

$$\min(P) := \max\{\min(Q) \mid Q \in \operatorname{Opt}(P)\}.$$

The minimum excludant of the empty set is 0, so the terminal positions of a game have nim-number 0. The *nim-number of a game* is the nim-number of its starting position. A winning strategy exists at position P if and only if $\operatorname{nim}(P) \neq 0$. So, the nim-number of a game determines the outcome of the game. We write P = R if $\operatorname{nim}(P) = \operatorname{nim}(R)$, so $P = * \operatorname{nim}(P)$ for every impartial game P.

2.2 Achievement Games for Groups

We now give a more precise description of the achievement game $\operatorname{\mathsf{GEN}}(G)$ played on a finite group G. We also recall some definitions and results from [7]. The nonterminal positions of $\operatorname{\mathsf{GEN}}(G)$ are exactly the non-generating subsets of G, and the terminal positions are the generating sets S of G such that there is an $g \in S$ satisfying $\langle S \setminus \{g\} \rangle < G$. The starting position is the empty set since neither player has chosen an element yet. The first player chooses $x_1 \in G$, and the designated player selects $x_k \in G \setminus \{x_1, \ldots, x_{k-1}\}$ at the *k*th turn. A position Q is an option of P if $Q = P \cup \{g\}$ for some $g \in G \setminus P$. The player who builds a generating set wins the game.

The set \mathcal{M} of maximal subgroups play a significant role in the game. The last two authors [7] define the set

$$\mathcal{I} := \{ \cap \mathcal{N} \mid \emptyset \neq \mathcal{N} \subseteq \mathcal{M} \}$$

of intersection subgroups, which is the set of all possible intersections of maximal subgroups. The smallest intersection subgroup is the Frattini subgroup $\Phi(G)$ of G, which is the intersection of all maximal subgroups of G.

Example 2.1. Let $G = \langle r, s | r^4 = e = s^2, sr^3 = rs \rangle$ be the dihedral group of order 8 with identity e. The maximal subgroups $\langle r \rangle$, $\langle s, sr^2 \rangle$, and $\langle sr, sr^3 \rangle$ are the order 4 intersection subgroups. The only other intersection subgroup is the Frattini subgroup $\langle r^2 \rangle$, which has order 2.

The set \mathcal{I} of intersection subgroups is partially ordered by inclusion. We use interval notation to denote certain subsets of \mathcal{I} . For example, if $I \in \mathcal{I}$, then $(-\infty, I) := \{J \in \mathcal{I} \mid J \leq I\}.$

For each $I \in \mathcal{I}$ let

$$X_I := \mathcal{P}(I) \setminus \bigcup \{ \mathcal{P}(J) \mid J \in (-\infty, I) \}$$

be the collection of those subsets of I that are not contained in any other intersection subgroup properly contained in I. We let $\mathcal{X} := \{X_I \mid I \in \mathcal{I}\}$ and call an element of \mathcal{X} a structure class. We define an additional structure class X_G to be the set of terminal positions, and we let $\mathcal{Y} := \mathcal{X} \cup \{X_G\}$. For any position P of $\mathsf{GEN}(G)$, let $\lceil P \rceil$ be the unique element of $\mathcal{I} \cup \{G\}$ such that $P \in X_{\lceil P \rceil}$. Note that $\lceil P \rceil$ is the smallest intersection subgroup containing P if P is not a terminal position. We will write $\lceil A, g_1, \ldots, g_n \rceil$ to mean $\lceil A \cup \{g_1, \ldots, g_n\} \rceil$ for $A \subseteq G$ and $g_1, \ldots, g_n \in G$.

Example 2.2. We continue Example 2.1. The structure classes in \mathcal{X} are

$$X_{\langle r^2 \rangle} = \{ \emptyset, \{e\}, \{r^2\}, \{e, r^2\} \}$$
$$X_{\langle r \rangle} = \mathcal{P}(\langle r \rangle) \setminus X_{\langle r^2 \rangle}$$
$$X_{\langle s, sr^2 \rangle} = \mathcal{P}(\langle s, sr^2 \rangle) \setminus X_{\langle r^2 \rangle}$$
$$X_{\langle sr, sr^3 \rangle} = \mathcal{P}(\langle sr, sr^3 \rangle) \setminus X_{\langle r^2 \rangle}.$$

Thus, for example $\lceil \emptyset \rceil = \langle r^2 \rangle = \lceil e \rceil$, $\lceil r \rceil = \langle r \rangle$, and $\lceil s \rceil = \langle s, sr^2 \rangle = \lceil s, sr^2 \rceil$.

Parity plays a crucial role in the theory of impartial games. We define the parity of a natural number n via $pty(n) := (1 + (-1)^{n+1})/2$. The parity of a subset of a group is defined to be the parity of the size of the subset. The parity of a structure class is defined to be $pty(X_I) := pty(I)$. We will say an object is even if its parity is 0 and odd if its parity is 1.

The partition \mathcal{Y} of the set of game positions of $\mathsf{GEN}(G)$ is compatible with the option relationship between game positions [7, Corollary 3.11]: if $X_I, X_J \in \mathcal{X}$ and $P, Q \in X_I \neq X_J$, then $\mathsf{Opt}(P) \cap X_J \neq \emptyset$ if and only if $\mathsf{Opt}(Q) \cap X_J \neq \emptyset$. We say that X_J is an *option* of X_I and we write $X_J \in \mathsf{Opt}(X_I)$ if $\mathsf{Opt}(I) \cap X_J \neq \emptyset$.

If $P, Q \in X_I \in \mathcal{Y}$ and pty(P) = pty(Q), then nim(P) = nim(Q) by [7, Proposition 4.4]. The structure digraph of GEN(G) has vertex set \mathcal{Y} and edge set $\{(X_I, X_J) \mid X_J \in \text{Opt}(X_I)\}$. In a structure diagram, a structure class X_I is represented by a triangle pointing down if I is odd and by a triangle pointing up if I is even. The triangles are divided into a smaller triangle and a trapezoid, where the smaller triangle represents the odd positions of X_I and the trapezoid represents the even positions of X_I . The numbers in the smaller triangle and the trapezoid are the nim-numbers of these positions. There is a directed arrow from X_I to X_J provided $X_J \in \text{Opt}(X_I)$.

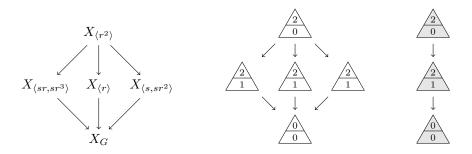


Figure 1: A depiction of the structure digraph, the structure diagram, and the simplified structure diagram, respectively, of GEN(G) for the dihedral group G of order 8.

Example 2.3. Figure 1 shows the structure digraph and the structure diagram of GEN(G) for the dihedral group G of Example 2.1. The numbers in the diagram were computed from the bottom up using nim-number arithmetic as depicted in Figure 2.

The type of the structure class X_I is the triple

$$type(X_I) := (pty(I), nim(P), nim(Q)),$$

where $P, Q \in X_I$ with pty(P) = 0 and pty(Q) = 1. Note that the type of a structure class X_I is determined by the parity of X_I and the types of the options of X_I as shown in Figure 2. We define $otype(X_I) := \{type(X_J) \mid X_J \in Opt(X_I)\}$.

The nim-number of the game is the nim-number of the initial position \emptyset , which is an even subset of $\Phi(G)$. Because of this, $\min(\mathsf{GEN}(G))$ is the second component of

$$\operatorname{type}(X_{\Phi(G)}) = (\operatorname{pty}(\Phi(G)), \operatorname{nim}(\emptyset), \operatorname{nim}(\{e\})),$$



Figure 2: Examples for the calculation of the type of a structure class using the types of the options.

which corresponds to the trapezoidal part of the triangle representing the source vertex $X_{\Phi(G)}$ of the structure diagram.

The simplified structure diagram of GEN(G) is built from the structure diagram by identifying structure classes X_I and X_J satisfying $\text{type}(X_I) = \text{type}(X_J)$ and $\text{otype}(X_I) \cup \{\text{type}(X_I)\} = \text{otype}(X_J) \cup \{\text{type}(X_J)\}$. This is followed by the removal of any resulting loops. See [7] for a more detailed description.

Example 2.4. The simplified structure diagram of GEN(G) is shown in Figure 1 for the dihedral group G of Example 2.1.

3 Deficiency

We will develop some general tools in this section. For a finite group G, the smallest possible size of a generating set is denoted by

$$d(G) := \min\{|S| : \langle S \rangle = G\}.$$

The following definition is closely related to d(G).

Definition 3.1. The *deficiency* of a subset P of a finite group G is the minimum size $\delta(P)$ of a subset Q of G such that $\langle P \cup Q \rangle = G$.

Note that the deficiency of a generating set is 0 and $\delta(\emptyset) = d(G)$. Also, $P \subseteq Q$ implies $\delta(P) \ge \delta(Q)$.

Proposition 3.2. If $P, Q \in X_I$, then $\delta(P) = \delta(Q)$.

Proof. Let \mathcal{L} be the set of maximal subgroups containing I. Then $I = \cap \mathcal{L}$ because I is an intersection subgroup. Let $g_1, \ldots, g_{\delta(P)} \in G$ such that $\langle P, g_1, \ldots, g_{\delta(P)} \rangle = G$. Then no element of \mathcal{L} contains $\{g_1, \ldots, g_{\delta(P)}\}$. The only maximal subgroups that contain Q are also in \mathcal{L} . Therefore, $\langle Q, g_1, \ldots, g_{\delta(P)} \rangle = G$ and $\delta(Q) \leq \delta(P)$. By a symmetric argument, $\delta(P) \leq \delta(Q)$, so $\delta(P) = \delta(Q)$.

Corollary 3.3. If G is a finite group, then $\delta(\Phi(G)) = d(G)$.

Proof. Since $\emptyset, \Phi(G) \in X_{\Phi(G)}$, we have $\delta(\Phi(G)) = \delta(\emptyset) = d(G)$.

Definition 3.4. We say that a structure class $X_I \in \mathcal{Y}$ is *m*-deficient in G and we write $\delta(X_I) = m$ if $\delta(I) = m$.

The terminal structure class is 0-deficient. That is, $\delta(X_G) = 0$.

Definition 3.5. Let G be a finite group, \mathcal{E} be the set of even structure classes and \mathcal{O} be the set of odd structure classes in \mathcal{Y} . We define the following sets:

$$\begin{aligned} \mathcal{D}_m &:= \{ X_I \in \mathcal{Y} \mid \delta(X_I) = m \}, \quad \mathcal{D}_{\geq m} := \cup_{k \geq m} \mathcal{D}_k \\ \mathcal{E}_m &:= \mathcal{E} \cap \mathcal{D}_m, \\ \mathcal{O}_m &:= \mathcal{O} \cap \mathcal{D}_m, \end{aligned}$$

$$\begin{aligned} \mathcal{D}_{\geq m} &:= \cup_{k \geq m} \mathcal{E}_k \\ \mathcal{O}_{\geq m} &:= \cup_{k \geq m} \mathcal{O}_k \end{aligned}$$

Note that $\mathcal{D}_0 = \{X_G\}.$

Proposition 3.6. If G is a finite group, then $\mathcal{D}_{\geq d(G)+1} = \emptyset$.

Proof. There must be a generating set Q for G with size d(G). If $X_I \in \mathcal{D}_m$, then $G = \langle Q \rangle \subseteq \langle I \cup Q \rangle$, and so $m \leq |Q| = d(G)$.

The next proposition follows immediately from Corollary 3.3.

Proposition 3.7. If G is a finite group, then $X_{\Phi(G)} \in \mathcal{D}_{d(G)}$.

Proposition 3.8. Let G be a finite group and m be a positive integer. If $X_I \in \mathcal{D}_m$, then X_I has an option X_K in \mathcal{D}_{m-1} , and every option of X_I is in either \mathcal{D}_m or \mathcal{D}_{m-1} .

Proof. Let X_J be an option of X_I ; that is, $Q := I \cup \{g\}$ is in X_J for some $g \in G$. Since $I \subseteq Q$,

$$\delta(X_I) = \delta(I) \ge \delta(Q) = \delta(J) = \delta(X_J).$$

Let $h_1, \ldots, h_{\delta(Q)} \in G$ such that $\langle Q, h_1, \ldots, h_{\delta(Q)} \rangle = G$. Then $\langle I, g, h_1, \ldots, h_{\delta(Q)} \rangle = G$, which implies

$$\delta(X_I) = \delta(I) \le \delta(Q) + 1 = \delta(X_J) + 1.$$

Hence $\delta(X_J) \in \{\delta(X_I), \delta(X_I) - 1\}.$

Now let $g_1, \ldots, g_{\delta(I)} \in G$ such that $\langle I, g_1, \ldots, g_{\delta(I)} \rangle = G$, and let $K = \lceil I \cup \{g_1\} \rceil$. Then $\delta(X_K) \in \{\delta(X_I), \delta(X_I) - 1\}$. Because $\langle I \cup \{g_1\}, g_2, \ldots, g_{\delta(I)} \rangle = G$, we conclude that $\delta(X_K) = \delta(X_I) - 1$.

Note that the deficiency of a structure class X_I is the directed distance from X_I to the terminal structure class X_G in the structure digraph.

Even structure classes only have even options by Lagrange's Theorem, so the next corollary follows by Proposition 3.8.

Corollary 3.9. If *m* is a positive integer and $X_I \in \mathcal{E}_m$, then X_I has an option in \mathcal{E}_{m-1} and every option of X_I is in either \mathcal{E}_m or \mathcal{E}_{m-1} .

The proof of the following result uses calculations similar to those shown in Figure 2.

Proposition 3.10. If G is a group of even order, then

$$\operatorname{type}(X_{I}) = \begin{cases} (0,0,0) & \text{if } X_{I} = X_{G} \in \mathcal{E}_{0} \\ (0,1,2) & \text{if } X_{I} \in \mathcal{E}_{1} \\ (0,0,2) & \text{if } X_{I} \in \mathcal{E}_{2} \\ (0,0,1) & \text{if } X_{I} \in \mathcal{E}_{\geq 3}. \end{cases}$$

Proof. We will use structural induction on the structure classes. By Corollary 3.9, structure classes in \mathcal{E}_m with $m \geq 1$ must have an option in \mathcal{E}_{m-1} and possibly more options in $\mathcal{E}_m \cup \mathcal{E}_{m-1}$. Hence $\mathcal{E}_m \neq \emptyset$ if and only if $m \leq d(G)$.

If $X_I \in \mathcal{E}_0$, then type (X_I) = type (X_G) = (0, 0, 0). If $X_I \in \mathcal{E}_1$, then type (X_I) = (0, 1, 2) since

otype
$$(X_I) = \begin{cases} \{(0,0,0)\} & \text{if } \operatorname{Opt}(X_I) \cap \mathcal{E}_1 = \emptyset \\ \{(0,0,0), (0,1,2)\} & \text{otherwise.} \end{cases}$$

If $X_I \in \mathcal{E}_2$, then type $(X_I) = (0, 0, 2)$ since

$$otype(X_I) = \begin{cases} \{(0,1,2)\} & \text{if } Opt(X_I) \cap \mathcal{E}_2 = \emptyset \\ \{(0,1,2), (0,0,2)\} & \text{otherwise.} \end{cases}$$

If $X_I \in \mathcal{E}_3$, then type $(X_I) = (0, 0, 1)$ since

otype
$$(X_I) = \begin{cases} \{(0,0,2)\} & \text{if } \operatorname{Opt}(X_I) \cap \mathcal{E}_3 = \emptyset \\ \{(0,0,2), (0,0,1)\} & \text{otherwise.} \end{cases}$$

If $X_I \in \mathcal{E}_{\geq 4}$, then type $(X_I) = (0, 0, 1)$, since every option of X_I has type (0, 0, 1). \Box

Corollary 3.11. If G is a group of even order with $d(G) \ge 4$, then $\mathsf{GEN}(G) = *0$.

Proof. If the first player initially selects an involution, the second player selects the identity; otherwise, the second player selects an involution t. This guarantees that the game position P after the second turn satisfies $P \in X_I \in \mathcal{E}_{\geq d(G)-2}$ for some structure class X_I . Because $d(G) - 2 \geq 2$, type (X_I) is (0, 0, 2) or (0, 0, 1) by Proposition 3.10. In both cases the even position P has nim-number 0, so the second player wins. \Box

4 The Generalized Dihedral Group Dih(A)

For a finite abelian group A, the generalized dihedral group Dih(A) is the semidirect product $C_2 \ltimes A$ where the generator x of the two-element cyclic group $C_2 = \{1, x\}$ acts on A on the right by inversion. That is, the operations on the elements of $C_2 \times A$ are defined by

$$(k,a)(1,b) := (k,ab), \quad (k,a)(x,b) := (kx,a^{-1}b)$$

for $a, b \in A$ and $k \in C_2$.

Note that we use multiplicative notation for group operations. We identify the elements of A and C_2 with their images in $C_2 \ltimes A$ through the natural inclusion maps $k \mapsto (k, 1) : C_2 \hookrightarrow \text{Dih}(A)$ and $a \mapsto (1, a) : A \hookrightarrow \text{Dih}(A)$ so that we can write ka for (k, a). This identification allows us to consider the action of C_2 on A as conjugation since $(ka)(xb) = kxx^{-1}axb = (kx)(a^xb)$.

Lemma 4.1. Let S be a subset of Dih(A) and $y \in Dih(A) \setminus A$. If $a, b \in A$, then $\langle S, ya, yb \rangle = \langle S, a^{-1}b, yb \rangle$.

Proof. Let $H = \langle S, ya, yb \rangle$ and $K = \langle S, a^{-1}b, yb \rangle$. Then $S \subseteq H$, $yb \in H$, and $a^{-1}b = (yay)b = (ya)(yb) \in H$, so $K \leq H$. Also, $S \subseteq K$, $yb \in K$, and $ya = (yb)(b^{-1}a) = (yb)(a^{-1}b)^{-1} \in K$, so $H \leq K$. Therefore, H = K.

Lemma 4.2. If $\langle B, y \rangle = \text{Dih}(A)$ for some $B \leq A$ and $y \in \text{Dih}(A) \setminus A$, then A = B.

Proof. It is easily verified that $y^{-1}by = b^{-1}$ for all $b \in B$. So $\langle y \rangle$ normalizes B, and hence B is normal in $\langle B, y \rangle = \text{Dih}(A)$. Thus, $\text{Dih}(A) = \langle B, y \rangle \subseteq B \langle y \rangle$, which implies $\text{Dih}(A) = B \langle y \rangle$. Since $B \cap \langle y \rangle$ is trivial, $B \langle y \rangle = \langle y \rangle \ltimes B$ after the natural identification. Since y acts on B by inversion, $\langle y \rangle \ltimes B = \text{Dih}(B)$. Therefore, Dih(A) = Dih(B), which shows that A = B.

Proposition 4.3. If A is a finite abelian group, then d(Dih(A)) = d(A) + 1.

Proof. Let $a_1, \ldots, a_{d(A)} \in A$ such that $\langle a_1, \ldots, a_{d(A)} \rangle = A$. Since $|\operatorname{Dih}(A) : A| = 2$, $\langle a_1, \ldots, a_{d(A)}, y \rangle = \operatorname{Dih}(A)$

for all $y \in \text{Dih}(A) \setminus A$, so $d(\text{Dih}(A)) \le d(A) + 1$.

Now let n = d(Dih(A)) and $g_1, \ldots, g_n \in \text{Dih}(A)$ such that $\langle g_1, \ldots, g_n \rangle = \text{Dih}(A)$. Without loss of generality, let $g_n \in \text{Dih}(A) \setminus A$. By possible repeated use of Lemma 4.1 we may find $B := \langle b_1, \ldots, b_{n-1} \rangle \leq A$ such that $\langle b_1, \ldots, b_{n-1}, g_n \rangle = \langle g_1, \ldots, g_n \rangle$. Then $\langle B, g_n \rangle = \text{Dih}(A)$, so B = A by Lemma 4.2. Therefore, $d(A) \leq n - 1 = d(\text{Dih}(A)) - 1$.

The following result is likely well-known, but we will include the proof for want of a reference.

Proposition 4.4. If G = Dih(A), then $\Phi(G) = \Phi(A)$.

Proof. Since A is a maximal subgroup of G, $\Phi(G)$ is contained in A.

We will first prove that $\Phi(A)$ is contained in $\Phi(G)$. Let z be in $\Phi(A)$ and M be a maximal subgroup of G. If M = A, then z is certainly in M, so we will assume that M is not equal to A. Then $M \cap A$ is maximal in A, so $z \in M \cap A \subseteq M$. Therefore, $\Phi(A) \subseteq \Phi(G)$.

We now prove that $\Phi(G)$ is contained in $\Phi(A)$. Let z be in $\Phi(G)$ and L be a maximal subgroup of A. Then $M = \langle x \rangle \ltimes L$ is a maximal subgroup of G, so $z \in M$. Since z is in $\Phi(G)$ and A is a maximal subgroup of G, z is in A. Therefore, $z \in M \cap A = L$. Therefore, $\Phi(G) \leq \Phi(A)$, and we conclude that $\Phi(G) = \Phi(A)$. \Box

5 The Achievement Game GEN(Dih(A))

We now determine the nim-number of GEN(Dih(A)). We consider three cases: d(A) = 1, $d(A) \ge 3$, and d(A) = 2. If d(A) = 1, then A is a cyclic group and Dih(A) is a usual dihedral group characterized by the following result.

Proposition 5.1. [7, Corollary 7.10] If A is a cyclic group of order n, then

$$\mathsf{GEN}(\mathrm{Dih}(A)) = \begin{cases} *0 & \text{if } n \equiv_4 0 \\ *1 & \text{if } n \equiv_4 2 \\ *3 & \text{if } n \equiv_2 1. \end{cases}$$

Proposition 5.2. If A is an abelian group such that $d(A) \ge 3$, then $\mathsf{GEN}(\mathsf{Dih}(A)) = *0$.

Proof. Since $d(A) \ge 3$, $d(\text{Dih}(A)) \ge 4$. The result follows by Corollary 3.11 because Dih(A) is even.

We now focus on the case when d(A) = 2. We start by considering the case when A is even. Figures 6(b) and 7(a) show our conjectured two possible simplified structure diagrams for this situation.

Proposition 5.3. If A is an abelian group of even order such that d(A) = 2, then GEN(Dih(A)) = *0.

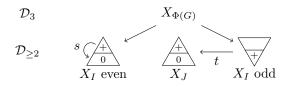


Figure 3: Visual representation for the proof of Proposition 5.3.

Proof. Because d(A) = 2, we know that d(Dih(A)) = 3 by Proposition 4.3, and so the initial position \emptyset has deficiency 3. We show that the second player can create a position with nim-number 0 in the second turn, no matter what the first player chooses in the first turn. This will show that the second player has a winning strategy and the nim-number of the game is 0.

Suppose the first player chooses $g \in \text{Dih}(A)$, and let $I = \lceil g \rceil$. Then X_I is in \mathcal{D}_2 or \mathcal{D}_3 by Proposition 3.8. We have provided a visualization in Figure 3.

If X_I is in \mathcal{E}_2 or \mathcal{E}_3 , then the second player picks any element $s \in I \setminus \{g\}$ to move to a position with nim-number 0 by Proposition 3.10 since type (X_I) is either (0, 0, 1)or (0, 0, 2). Now assume that X_I is in \mathcal{O}_2 or \mathcal{O}_3 . Because every element in $\text{Dih}(A) \setminus A$ is an involution, we conclude that I is an odd subgroup of A and g has odd order. The second player then selects an involution $t \in A$, and we consider a new structure class X_J for J = [I, t]. Because gcd(|g|, |t|) = 1 and A is abelian, $\langle g, t \rangle = \langle gt \rangle$ and hence J = [gt]. Therefore, X_J is in \mathcal{E}_2 or \mathcal{E}_3 . Since $\text{type}(X_J)$ is either (0, 0, 1) or (0, 0, 2) by Proposition 3.10, the second player moves to a position with nim-number 0 by choosing t.

We now consider the case when A is odd.

Lemma 5.4. Let G = Dih(A) where A is an abelian group of odd order and d(A) = 2. If $X_I \in \mathcal{O}_m$ with $m \in \{2, 3\}$, then X_I has an option in \mathcal{O}_{m-1} .

Proof. Recall that $C_2 = \langle x \rangle$. Since $X_I \in \mathcal{O}_m$, there are distinct g_1, \ldots, g_m such that $\langle I, g_1, \ldots, g_m \rangle = G$. Additionally, $I \leq A$ because I is odd. If g_i has odd order for some i, then $g_i \in A$ and $X_{\lceil I, g_i \rceil} \in \mathcal{O}_{m-1}$.

Now assume that all g_1, \ldots, g_m have even order. Since g_1 and g_2 are involutions, there are $a, b \in A$ such that $g_1 = xa$ and $g_2 = xb$. Then $\langle g_1, g_2 \rangle = \langle xa, xb \rangle = \langle a^{-1}b, xb \rangle$ if m = 2 and $\langle g_1, g_2, g_3 \rangle = \langle a^{-1}b, g_2, g_3 \rangle$ if m = 3 by Lemma 4.1, so $X_{[I,a^{-1}b]} \in \mathcal{O}_{m-1}$.

Lemma 5.5. Let G = Dih(A) where A is an abelian group of odd order such that d(A) = 2. If $X_I \in \mathcal{O}_m$ with $m \in \{1, 2, 3\}$, then X_I has an option in \mathcal{E}_{m-1} but not in \mathcal{E}_m .

Proof. Since X_I is odd, I is a subgroup of A. So $X_{\lceil I,h\rceil}$ is an option of X_I for any $h \in G \setminus A$, and this option is in \mathcal{E}_{m-1} or \mathcal{E}_m . It remains to show that X_I has no option in \mathcal{E}_m .

We argue by contradiction. Assume that $X_I \in \mathcal{O}_m$ has an option $X_J \in \mathcal{E}_m$ and m is as small as possible with this property. Let J = [I, y] for some $y \in G \setminus A$.

We first consider the case where m = 1 and $X_I \in \mathcal{O}_1$. Then $X_{\lceil I,g \rceil} \in \mathcal{D}_0$ for some $g \in G \setminus A$. Then $G = \langle I, g \rangle$, so I = A by Lemma 4.2. Thus $X_J = X_{\lceil A, y \rceil} = X_G \in \mathcal{E}_0$, which is a contradiction.

Now we consider the case when $m \geq 2$, as shown in Figure 4. By Lemma 5.4, there is a $g \in G$ such that $X_{[I,g]} \in \mathcal{O}_{m-1}$. Since m is minimal, $X_{[I,g]}$ has no option in \mathcal{E}_{m-1} , and so $X_{[I,g,y]} \in \mathcal{E}_{m-2}$. This is a contradiction, since $X_{[I,g,y]}$ is an option of X_J and X_J does not have an option in \mathcal{E}_{m-2} by Proposition 3.8.

Lemma 5.6. If G = Dih(A) where A is an abelian group of odd order and d(A) = 2, then

$$\text{type}(X_I) = \begin{cases} (1,2,1) & \text{if } X_I \in \mathcal{O}_1 \\ (1,3,0) & \text{if } X_I \in \mathcal{O}_2 \\ (1,3,1) & \text{if } X_I \in \mathcal{O}_3. \end{cases}$$

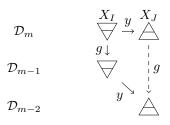


Figure 4: The $m \ge 2$ case of the proof of Lemma 5.5.

Proof. We will use structural induction on the structure classes. Proposition 3.8 and Lemmas 5.4 and 5.5 determine the possible options of our structure classes. Proposition 3.10 and induction yields the types of these options.

If $X_I \in \mathcal{O}_1$, then $\operatorname{Opt}(X_I) \cap \mathcal{E}_0 \neq \emptyset$ and $\operatorname{Opt}(X_I) \subseteq \mathcal{E}_0 \cup \mathcal{O}_1$. So type $(X_I) = (1, 2, 1)$ since

$$otype(X_I) = \begin{cases} \{(0,0,0)\} & \text{if } Opt(X_I) \cap \mathcal{O}_1 = \emptyset \\ \{(0,0,0), (1,2,1)\} & \text{otherwise.} \end{cases}$$

If $X_I \in \mathcal{O}_2$, then $\operatorname{Opt}(X_I) \cap \mathcal{E}_1 \neq \emptyset \neq \operatorname{Opt}(X_I) \cap \mathcal{O}_1$ and $\operatorname{Opt}(X_I) \subseteq \mathcal{E}_1 \cup \mathcal{O}_1 \cup \mathcal{O}_2$. So type $(X_I) = (1, 3, 0)$ since

$$otype(X_I) = \begin{cases} \{(0,1,2), (1,2,1)\} & \text{if } Opt(X_I) \cap \mathcal{O}_2 = \emptyset \\ \{(0,1,2), (1,2,1), (1,3,0)\} & \text{otherwise.} \end{cases}$$

If $X_I \in \mathcal{O}_3$, then $\operatorname{Opt}(X_I) \cap \mathcal{E}_2 \neq \emptyset \neq \operatorname{Opt}(X_I) \cap \mathcal{O}_2$ and $\operatorname{Opt}(X_I) \subseteq \mathcal{E}_2 \cup \mathcal{O}_2 \cup \mathcal{O}_3$. So $\operatorname{type}(X_I) = (1, 3, 1)$ since

otype(X_I) =

$$\begin{cases} \{(0,0,2), (1,3,0)\} & \text{if } \operatorname{Opt}(X_I) \cap \mathcal{O}_3 = \emptyset \\ \{(0,0,2), (1,3,0), (1,3,1)\} & \text{otherwise.} \end{cases}$$

Figure 5(b) shows our conjectured simplified structure diagram that corresponds to the next result.

Proposition 5.7. If A is an abelian group of odd order such that d(A) = 2, then GEN(Dih(A)) = *3.

Proof. By Proposition 3.7, we have $X_{\Phi(G)} \in \mathcal{D}_3$ since d(Dih(A)) = 3. By Lemma 4.4, we have $\Phi(G) = \Phi(A) \leq A$, so $\Phi(G)$ is odd. Therefore, $X_{\Phi(G)} \in \mathcal{O}_3$, so type $(X_{\Phi(G)}) = (1,3,1)$ by Lemma 5.6. The second component is 3, so we conclude that $\mathsf{GEN}(\mathrm{Dih}(A)) = *3$.

We conclude with our main result, which summarizes Propositions 5.1, 5.2, 5.3, and 5.7.

 \square

Theorem 5.8. If A is an abelian group, then

$$\mathsf{GEN}(\mathrm{Dih}(A)) = \begin{cases} *1 & \text{if } d(A) = 1, |A| \equiv_4 2 \\ *3 & \text{if } 1 \leq d(A) \leq 2, A \text{ odd} \\ *0 & \text{otherwise.} \end{cases}$$

6 Further Questions

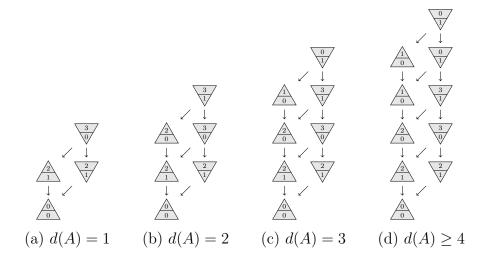


Figure 5: Conjectured simplified structure diagrams for GEN(Dih(A)) with odd A.

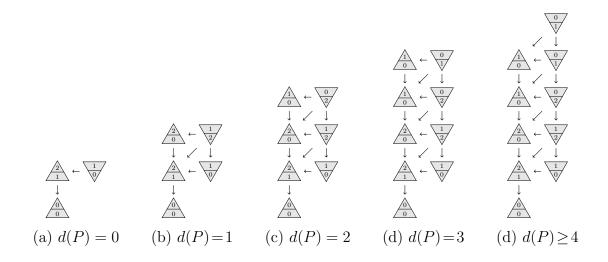


Figure 6: Conjectured simplified structure diagrams for GEN(Dih(A)) with $A = \mathbb{Z}_2 \times P$ and P odd.

We outline a few open problems.

- 1. The simplified structure diagrams in Figures 5–7 were found via computer calculations by considering several examples of GEN(Dih(A)). We conjecture that these figures cover all possible structure diagrams.
- 2. We believe that the key new insight of this paper is the use of deficiency. Can we use this tool to handle other semidirect products?
- 3. Can we extend the techniques in this paper to give a complete description of the nim-numbers of $\mathsf{GEN}(G)$ using covering properties as was done in [5] for $\mathsf{DNG}(G)$? It is likely that d(G) plays a role in this characterization.

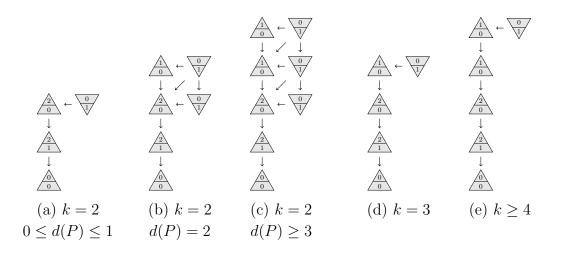


Figure 7: Conjectured simplified structure diagrams for $\mathsf{GEN}(\mathrm{Dih}(A))$ with $A = \mathbb{Z}_2^k \times P$ and P odd.

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