

The crossing number of the strong product of two paths

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Abstract

Let $P_m \boxtimes P_n$ be the strong product of two paths P_m and P_n . In 2013, Klešč et al. conjectured that the crossing number of $P_m \boxtimes P_n$ is equal to $(m - 1)(n - 1) - 4$ for $m \geq 4$ and $n \geq 4$. In this paper we show that the above conjecture is true except when $m = 4$ and $n = 4$, and that the crossing number of $P_4 \boxtimes P_4$ is four.

1 Introduction

Let G and H be two disjoint graphs. The strong product $G \boxtimes H$ of G and H is the graph with vertex set $V(G) \times V(H)$ and edge set $\{(u, v)(x, y) : u = x \text{ and } vy \in E(H), \text{ or } v = y \text{ and } ux \in E(G), \text{ or } ux \in E(G) \text{ and } vy \in E(H)\}$.

Suppose that j is a positive integer. Let P_j be a path with j vertices. In 2013, Klešč et al. [1] firstly studied the crossing number of the strong product of two graphs. They showed that the crossing number of $P_3 \boxtimes P_n$ is equal to $n - 3$ if $n \geq 3$, and they established

Lemma 1.1 [1] *The crossing number of $P_m \boxtimes P_n$ is at most $(m - 1)(n - 1) - 4$ for $m \geq 4$ and $n \geq 4$.*

Subsequently, Klešč et al. conjectured that the crossing number of $P_m \boxtimes P_n$ is equal to $(m - 1)(n - 1) - 4$ for $m \geq 4$ and $n \geq 4$. In this paper we shall show that the above conjecture is true except for $m = 4$ and $n = 4$, and that the crossing number of $P_4 \boxtimes P_4$ is four.

The arrangement of the paper is as follows. In Section 2, we give some lemmas and show that the crossing number of a subgraph of $P_4 \boxtimes P_5$ is at least eight. In

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Section 3, we first show using induction on n that the crossing number of a subgraph $W_{4,n}$ of $P_4 \boxtimes P_n$ is equal to $3(n-1)-4$ for $n \geq 5$. Then we prove using induction on m that the crossing number of a subgraph $W_{m,n}$ of $P_m \boxtimes P_n$ is equal to $(m-1)(n-1)-4$ for $m \geq 4$ and $n \geq 4$ and $(m, n) \neq (4, 4)$. Subsequently, we determine the crossing number of $P_m \boxtimes P_n$ when $m \geq 4$ and $n \geq 4$ and $(m, n) \neq (4, 4)$. In Section 4 we show that the crossing number of $P_4 \boxtimes P_4$ is four.

The rest of this section is contributed to some terminology for crossing numbers and graph theory.

Let G be a graph. By a drawing of G , we mean a drawing of G in the plane in which: no edge has a vertex as an interior point; no two adjacent edges cross each other; no two edges cross each other more than once; and no three edges cross in a common point.

Suppose that Φ is a drawing of a graph. The number of edge crossings in Φ is denoted by $\text{cr}(\Phi)$. The *crossing number* of a graph G , denoted by $\text{cr}(G)$, is the minimum number of edges crossings over all drawings of G . A drawing Ψ of G is *optimal* if $\text{cr}(\Psi) = \text{cr}(G)$.

A graph G' is a *subdivision* of G if G' is isomorphic to G or G' can be obtained from G by inserting vertices of degree two in some edges. Obviously, $\text{cr}(G') = \text{cr}(G)$ if G' is a subdivision of G . A graph H is a *minor* of G if H is isomorphic to a graph obtained from a subgraph F of G by contracting some edges in $E(F)$. A vertex of G is called a *branch vertex* if its degree is at least three in G . The complete graph with n vertices is denoted by K_n .

By Kuratowski's theorem [2], a graph is planar if and only if it contains no subdivision of either the complete graph K_5 or the complete bipartite graph $K_{3,3}$. Hence, if a graph G has a subgraph isomorphic to a subdivision of K_5 or $K_{3,3}$, then $\text{cr}(G) \geq 1$.

2 Basic lemmas and the crossing number of the graph in Figure 10

Let $P_m = u_1u_2 \dots u_m$ and $P_n = z_1z_2 \dots z_n$ be two paths, where $m \geq 4$ and $n \geq 4$. For brevity, the vertex (u_i, z_j) in $P_m \boxtimes P_n$ is labeled by $w_{i,j}$. It is easy to find that there are many induced subgraphs in $P_m \boxtimes P_n$ such that each is isomorphic to K_4 . For example, the graph $P_4 \boxtimes P_5$ has twelve induced subgraphs in which each is isomorphic to K_4 . But K_4 is a planar graph. In order to give a lower bound for the crossing number of $P_m \boxtimes P_n$, we need a nonplanar graph which contains K_4 as subgraph. We shall define this graph in next paragraph. The drawing of $P_4 \boxtimes P_5$ shown in Figure 1 can be generalized to obtain a drawing of $P_m \boxtimes P_n$ with $(m-1)(n-1)-4$ crossings.

Let T_6 be the graph shown in Figure 2. Obviously, T_6 has a subgraph isomorphic to $K_{3,3}$. So $\text{cr}(T_6) \geq 1$. It is easy to find that there are many subgraphs in $P_m \boxtimes P_n$ such that each is isomorphic to a subdivision of T_6 .

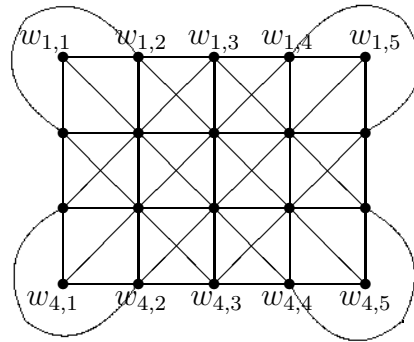


Figure 1 The graph $P_4 \boxtimes P_5$

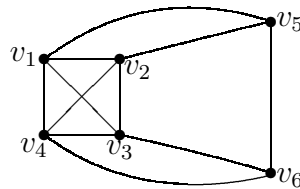


Figure 2 The graph T_6

Lemma 2.1 *Let F be a graph isomorphic to a subdivision of T_6 . Let Q be a subgraph of F which is isomorphic to a subdivision of K_4 . Then at least one edge of Q is crossed in any drawing of F .*

Proof Suppose that Φ is a drawing of F . Since F is isomorphic to a subdivision of T_6 , Φ has at least one crossing.

Suppose that v_1, v_2, v_3 and v_4 are the four branch vertices of Q . If some edge of Q is crossed by some other edge of Q , then we have the desired result. Otherwise, the drawing of Φ restricted in Q divides the plane into four regions in which one is unbounded and its closure contains three branch vertices of Q . Also, the boundary of each of the other three regions contains three branch vertices of Q . Let F' be the graph obtained from F by deleting all edges in $E(Q)$. Then F' is a connected graph. Suppose that v_5 and v_6 are the other two branch vertices of F . By the Jordan Curve Theorem, no matter which regions v_5 and v_6 are in, at least one edge of Q is crossed by some edge in F' . □

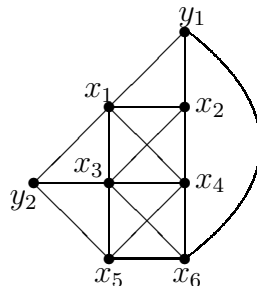


Figure 3 The graph H_1

Lemma 2.2 *Let H_1 be the graph shown in Figure 3. Then $cr(H_1) = 2$.*

Proof Let Q be the subgraph of H_1 induced by the four vertices x_1, x_2, x_3 and x_4 . Then Q is isomorphic to K_4 . It is easy to find that H_1 has a subgraph isomorphic to T_6 which contains Q . By Lemma 2.1, some edge in Q has at least one crossing in any drawing of H_1 . For any edge e in Q , it can be checked that $H_1 - e$ has a minor isomorphic to K_5 . So $cr(H_1) \geq cr(K_5) + 1 \geq 2$. Also, Figure 3 exhibits a drawing of H_1 with two crossings. Hence $cr(H_1) = 2$. \square

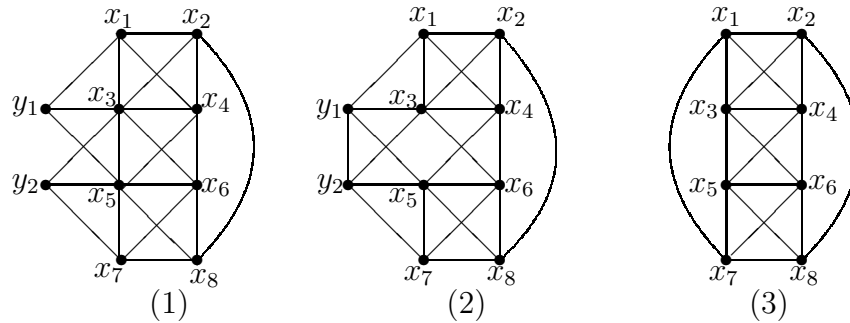


Figure 4 Three graphs H_2, H_3 and H_4

Lemma 2.3 *Let H_2 be the graph shown in Figure 4(1). Then $cr(H_2) \geq 3$.*

Proof Suppose that Φ is an optimal drawing of H_2 . Let Q be the subgraph of H_2 induced by the four vertices $x_1, x_2, x_3,$ and x_4 . Clearly, Q is isomorphic to K_4 , and H_2 has a subgraph isomorphic to a subdivision of T_6 which contains Q . By Lemma 2.1, some edge e in Q has at least one crossing in Φ . If e is some edge in $E(Q) \setminus \{x_2x_3, x_2x_4, x_3x_4\}$, it is easy to find that $H_2 - e$ contains a subgraph isomorphic to H_1 . If e is x_2x_3 or x_2x_4 , then $H_2 - e$ contains the path $x_1x_2x_8$ and the path $x_3x_1x_4$. So $H_2 - e$ has a subgraph isomorphic to a subdivision of H_1 . If e is the edge x_3x_4 , then $H_2 - e$ contains the path $x_3x_1x_4$ and the path $x_3x_2x_4$. Hence $H_2 - e$ has a subgraph isomorphic to a subdivision of H_1 . So $cr(H_2) \geq cr(H_1) + 1 \geq 3$. \square

Lemma 2.4 *Let H_3 be the graph shown in Figure 4(2). Then $cr(H_3) \geq 3$.*

Proof Let Q be the subgraph of H_3 induced by the four vertices $x_1, x_2, x_3,$ and x_4 . Proceeding the similar argument to that in the proof of Lemma 2.3, one can show that $cr(H_3) \geq 3$. \square

Lemma 2.5 *Let H_4 be the graph shown in Figure 4(3). Then $cr(H_4) \geq 3$.*

Proof Suppose that $cr(H_4) = k$, and that Φ is an optimal drawing of H_4 .

For $i = 1, 2, 3$, let Q_i be the subgraph of H_4 induced by the four vertices $x_{2i-1}, x_{2i}, x_{2i+1}$ and x_{2i+2} . Then each Q_i is isomorphic to K_4 . It is easy to find that H_4 contains a subgraph isomorphic to a subdivision of T_6 which contains Q_1 . By Lemma 2.1, some edge e_1 in Q_1 has at least one crossing in Φ . Let Φ'_1 be the drawing obtained from Φ by deleting e_1 . Then $cr(\Phi'_1) \leq k - 1$. It is easy to find that $H_4 - e_1$ contains a subgraph which is isomorphic to a subdivision of T_6 which contains Q_3 . By Lemma 2.1, some edge e_2 in Q_3 has at least one crossing in Φ'_1 . Let Φ'_2 be the drawing obtained from Φ'_1 by deleting e_2 . Then $cr(\Phi'_2) \leq k - 2$.

We observe that $H_4 - e_1 - e_2$ has the following properties.

- (1) It contains the path $x_3x_1x_4$ or $x_3x_2x_4$.
- (2) It contains the path $x_5x_7x_6$ or $x_5x_8x_6$.

Without loss of generality, suppose that $H_4 - e_1 - e_2$ contains the path $x_3x_1x_4$. If $H_4 - e_1 - e_2$ contains $x_5x_7x_6$, then $H_4 - e_1 - e_2$ contains a subgraph isomorphic to $K_{3,3}$ if the cycle $x_3x_5x_4x_6x_3$ is considered. Otherwise, $H_4 - e_1 - e_2$ contains $x_5x_8x_6$. Moreover, $H_4 - e_1 - e_2$ contains x_7x_8 . In this case $H_4 - e_1 - e_2$ contains a subgraph isomorphic to a subdivision of $K_{3,3}$ if the cycle $x_3x_5x_4x_6x_3$ is considered. So $H_4 - e_1 - e_2$ is nonplanar. Thus, $k \geq 3$. □

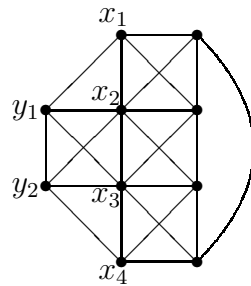


Figure 5 The graph H_5

Lemma 2.6 *Let H_5 be the graph shown in Figure 5. Then $cr(H_5) = 4$.*

Proof Suppose that $cr(H_5) = k$, and suppose that Φ is an optimal drawing of H_5 .

Let Q be the induced subgraph of H_5 by the four vertices y_1, y_2, x_2 and x_3 . Obviously, Q is isomorphic to K_4 , and H_5 has a subgraph isomorphic to a subdivision of T_6 which contains Q . By Lemma 2.1, some edge e in Q has at least one crossing in Φ . Let Φ' be the drawing obtained from Φ by deleting e . If e is some edge in $\{y_1x_2, y_1x_3, y_2x_2, y_2x_3\}$, then $H_5 - e$ has a subgraph isomorphic to a subdivision of H_4 defined in Lemma 2.5. Thus, we have that $k \geq 4$ in this case. Otherwise, we consider the edge y_1y_2 . If it has at least one crossing in Φ , then we take it as e . So $H_5 - e$ is isomorphic to H_2 defined in Lemma 2.3. Then $k \geq 4$. If y_1y_2 has not any crossing in Φ , then e is exactly x_2x_3 in Φ . So $H_5 - e$ is isomorphic to H_3 defined in Lemma 2.4. So $k \geq 4$. Also, Figure 5 exhibits a drawing of H_5 with four crossings. Hence $cr(H_5) = 4$. □

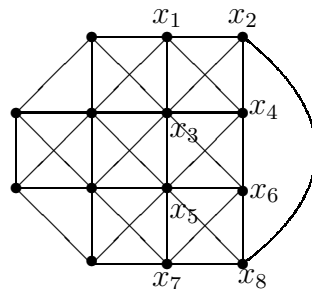


Figure 6 The graph G_1

Lemma 2.7 *Let G_1 be the graph shown in Figure 6. Then $cr(G_1) = 7$.*

Proof Suppose that $cr(G_1) = k$, and that Ψ is an optimal drawing of G_1 .

For $i = 1, 2, 3$, let Q_i be the induced subgraph of G_1 by the four vertices x_{2i-1} , x_{2i} , x_{2i+1} and x_{2i+2} . Then each Q_i is isomorphic to K_4 . Moreover, G_1 has a subgraph isomorphic to a subdivision of T_6 which contains Q_1 . By Lemma 2.1, some edge in Q_1 has at least one crossing in Ψ . We now apply the following operations.

(1) If some edge in $\{x_1x_2, x_1x_4, x_2x_3, x_2x_4\}$ has at least one crossing in Ψ , then it is deleted and turn to (4). Otherwise, turn to (2).

(2) Suppose that the edge x_1x_3 has at least one crossing in Ψ . Notice that x_2x_1 and x_2x_3 are not successive if the edges incident with x_2 are oriented in clockwise or anticlockwise in Ψ , otherwise, we redraw x_1x_3 near to the path $x_1x_2x_3$, obtaining a drawing of G_1 with at most $k - 1$ crossings, a contradiction. Now x_1x_3 is redrawn near to $x_1x_2x_3$ such that it crosses exactly x_2x_4 . Delete x_2x_4 and turn to (4). If x_1x_3 has not any crossing in Ψ , turn to (3).

(3) The edge x_3x_4 has at least one crossing in Ψ . A similar argument as the one used in (2) shows that we may redraw x_3x_4 near to the path $x_3x_2x_4$ such that it crosses exactly x_1x_2 . Delete x_1x_2 and turn to (4).

(4) Let Ψ'_1 be the obtained drawing, and let G'_1 the obtained graph. Then $cr(\Psi'_1) \leq k - 1$, and G'_1 contains the edge x_1x_4 or the path $x_1x_2x_4$.

Since G'_1 contains the edge x_2x_8 , it is easy to find that G'_1 has a subgraph isomorphic to a subdivision of T_6 which contains Q_3 . By Lemma 2.1, some edge in Q_3 has at least one crossing in Ψ'_1 . An argument similar to the one used for Q_1 shows that there is a drawing Ψ'_2 which is obtained from Ψ'_1 by deleting some edge e in $E(Q_3) \setminus \{x_5x_6, x_5x_7\}$ and $cr(\Psi'_2) \leq k - 2$. Let G''_1 be the graph obtained from G'_1 by deleting e . Then G''_1 contains the edge x_7x_6 or the path $x_7x_8x_6$.

Since G''_1 contains one of x_1x_4 and $x_1x_2x_4$ and one of x_7x_6 and $x_7x_8x_6$, G''_1 has a subgraph isomorphic to a subdivision of T_6 which contains Q_2 . By Lemma 2.1, some edge in Q_2 has at least one crossing in Ψ'_2 . We consider two cases.

Case 1: G''_1 contains one of x_1x_2 and x_7x_8 . Without loss of generality, suppose that G''_1 contains x_1x_2 . If some edge in $E(Q_2) \setminus \{x_3x_5\}$ has at least one crossing in Ψ'_2 , then it is deleted. Otherwise, x_3x_5 must has at least one crossing in Ψ'_2 . We now redraw x_3x_5 near to $x_3x_4x_5$. Notice that each of x_4x_3 and x_4x_5 has not any crossing in Ψ'_2 . But x_3x_5 may crosses some edges in $\{x_4x_1, x_4x_2, x_4x_6\}$. If this case occur, then we delete those edges. Thus, the new drawing of x_3x_5 has not any crossing. Let Ψ'_3 be the obtained drawing in the above procedure, and let \bar{G}_1 be the obtained graph. Then $cr(\Psi'_3) \leq k - 3$, and \bar{G}_1 contains one of $x_1x_2x_8x_7$ and $x_1x_2x_8x_6x_7$. Thus \bar{G}_1 has a subgraph isomorphic to a subdivision of the graph H_5 defined in Lemma 2.6. So $k - 3 \geq 4$. Thus $k \geq 7$.

Case 2: G''_1 contains none of x_1x_2 and x_7x_8 . In this case, G''_1 is isomorphic to the graph shown in Figure 7(1).

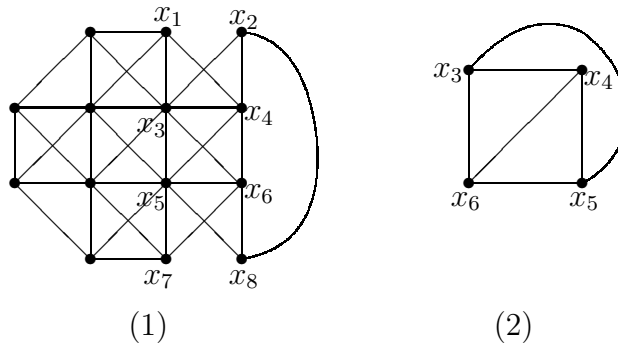


Figure 7 Two graphs defined in Case 2

If some edge in $E(Q_2) \setminus \{x_3x_5\}$ has at least one crossing in Ψ'_2 , then it is deleted. The graph obtained has a subgraph isomorphic to a subdivision of H_5 defined in Lemma 2.6. Following a similar argument to that in Case 1, we have $k \geq 7$. If not, then x_3x_5 has at least one crossing in Ψ'_2 . Since each edge in $E(Q_2) \setminus \{x_3x_5\}$ is not crossed, the drawing of Ψ'_2 restricted in $E(Q_2)$ is shown in Figure 7(2).

We observe that none of x_2 and x_8 is in the interior of the region whose boundary is the cycle $x_3x_4x_6x_3$. Otherwise, the existence of the path $x_2x_8x_5$ in G'_1 , and Jordan Curve Theorem show that one edge of the cycle $x_3x_4x_6x_3$ must be crossed. Similarly, none of x_2 and x_8 is in the interior of the region whose boundary is the cycle $x_4x_5x_6x_4$. Let F_1 be the region whose boundary is the cycle $x_3x_4x_5x_3$, let F_2 the unbounded region in Figure 7(2). If x_2 and x_8 are in F_1 and F_2 , respectively, then x_2x_8 must intersect x_3x_5 . In this case, x_2x_8 is deleted. Clearly, the obtained drawing has at most $k - 3$ crossings, and the obtained graph has a subgraph isomorphic to a subdivision of H_5 defined in Lemma 2.6. So $k \geq 7$. If x_2 and x_8 are in the same region, there are two cases to consider. If they are in F_1 , then x_6x_8 must intersect x_3x_5 . In this case, x_6x_8 is deleted. Then the obtained graph has a subgraph isomorphic to a subdivision of H_5 . Thus, $k \geq 7$. If they are in F_2 , then x_2x_4 must intersect x_3x_5 . Similarly, we have that $k \geq 7$. Notice that Figure 6 exhibits a drawing of G_1 with four crossings. Hence $\text{cr}(G_1) = 7$. □

Lemma 2.8 *Let G_2 be the graph shown in Figure 8. Then $\text{cr}(G_2) \geq 7$.*

Proof Suppose that $\text{cr}(G_2) = k$, and that Ψ is an optimal drawing of G_2 .

Let Q_1 be the subgraph of G_2 induced by the four vertices x_1, x_2, x_3 and x_4 . Clearly, Q_1 is isomorphic to K_4 , and G_2 has a subgraph isomorphic to a subdivision of T_6 which contains Q_1 . By Lemma 2.1, some edge in $E(Q_1)$ has at least one crossing in Ψ . We now apply the operations to Q_1 which are similar to those in the proof of Lemma 2.7. If some edge in $E(Q_1) \setminus \{x_1x_3, x_3x_4\}$ has at least one crossing in Ψ , then it is deleted. Otherwise, one of x_1x_3 and x_3x_4 has at least one crossing in Ψ . In this case the edge is redrawn and some edge in $\{x_1x_2, x_2x_4\}$ is deleted if necessary such that at least one crossing is eliminated.

Let Ψ'_1 be the drawing so obtained, and let G'_2 the graph. Then $\text{cr}(\Psi'_1) \leq k - 1$, and G'_2 contains one of the following three subgraphs.

- (a) The path $x_1x_2x_4$.
- (b) The path $x_1x_4x_2$.
- (c) $\{x_3x_2\} \cup \{x_1x_4\}$.

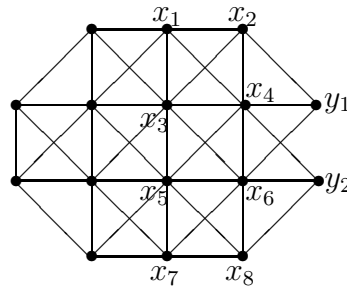


Figure 8 The graph G_2

Let Q_2 be the induced subgraph of G'_2 by the four vertices x_5, x_6, x_7 and x_8 . Obviously, Q_2 is isomorphic to K_4 , and G'_2 has a subgraph isomorphic to a subdivision of T_6 which contains Q_2 . By Lemma 2.1, some edge in $E(Q_2)$ has at least one crossing in Ψ'_1 . An argument similar to one used for Q_1 shows that there is a drawing Ψ'_2 obtained from Ψ'_1 by deleting some edge e in $E(Q_2) \setminus \{x_5x_6, x_5x_7\}$ and $\text{cr}(\Psi'_2) \leq k - 2$. Let G''_2 be the graph obtained. We observe that G''_2 contains one of the following subgraphs.

- (a) The path $x_7x_8x_6$.
- (b) The path $x_7x_6x_8$.
- (c) $\{x_5x_8\} \cup \{x_7x_6\}$.

Let Q_3 be the induced subgraph of G''_2 by the four vertices x_3, x_4, x_5 and x_6 . Then Q_3 is isomorphic to K_4 , and G''_2 has a subgraph isomorphic to a subdivision of T_6 which contains Q_3 . By Lemma 2.1, some edge in $E(Q_3)$ has at least one crossing in Ψ'_2 . We consider two cases.

Case 1: G''_2 contains one of x_1x_2 and x_7x_8 . Without loss of generality, suppose that G''_2 contains the edge x_1x_2 . If some edge e_1 in $E(Q_3) \setminus \{x_3x_5\}$ has at least one crossing in Ψ'_2 , then the drawing obtained from Ψ'_2 by deleting e_1 has at most $k - 3$ crossings. Moreover, $G''_2 - e_1$ contains a subdivision of the graph H_5 defined in Lemma 2.6, since there is a path $x_1x_2y_1x_6x_7$ or $x_1x_2y_1x_6x_8x_7$ in $G''_2 - e_1$. So $k \geq 7$. If not, then x_3x_5 has at least one crossing in Ψ'_2 . We now delete edges x_4y_1 and x_4y_2 . If x_2x_4 was not removed, then it is deleted. Next, x_3x_5 can be drawn near to the path $x_3x_4x_5$ such that it has at most one crossing. If x_3x_5 has one crossing, then it can be drawn such that it crosses exactly x_4x_6 . Thus, the drawing obtained from Ψ'_2 by deleting x_4x_6 has at most $k - 3$ crossings, and the obtained graph has a subgraph isomorphic to a subdivision of the graph H_5 defined in Lemma 2.6. So $k \geq 7$. If x_3x_5 has not any crossing, we also let Ψ'_2 be the obtained drawing. Clearly, $\text{cr}(\Psi'_2) \leq k - 3$. So $k \geq 7$.

Case 2: G''_2 contains none of x_1x_2 and x_7x_8 . In this case, G''_2 is the graph shown in Figure 9.

If some edge in $E(Q_3) \setminus \{x_3x_5\}$ has at least one crossing in Ψ'_2 , then it is deleted. The obtained graph has a subgraph isomorphic to a subdivision of the graph H_5 defined in Lemma 2.6. Hence, $k \geq 7$. Otherwise, x_3x_5 has at least one crossing in Ψ'_2 . The drawing of Ψ'_2 restricted in Q_3 is as in Figure 7(2). Let F_1 and F_2 be the regions whose boundaries are $x_3x_4x_6x_3$ and $x_4x_5x_6x_4$, respectively. Proceeding the similar argument as x_2 and x_8 in Case (2) in the proof of Lemma 2.7, none of x_1 and

x_7 is in F_1 or F_2 .

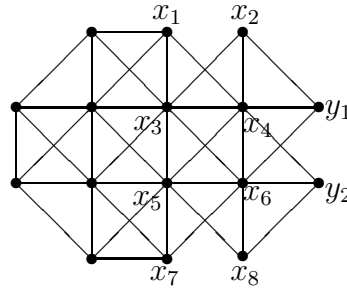


Figure 9 The graph defined in Case 2

Let J be the graph obtained from G_2'' by deleting $x_2, x_3, x_4, x_5, x_6, x_8, y_1$ and y_2 . Then J is connected graph. It is easy to find that there are two internally disjoint paths P_1 and P_2 from x_1 to x_7 in J . Let F_3 be the region whose boundary is the cycle $x_3x_4x_5x_3$, and let F_4 the unbounded region in Figure 7(2).

If x_1 and x_7 are in F_3 and F_4 , respectively, it can be found that x_3x_5 has at least two crossings if P_1 and P_2 are considered. We now delete x_3x_5 . Then the obtained graph has a subgraph isomorphic to a subdivision of the graph H_3 defined in Lemma 2.4. So $k \geq 7$.

If x_1 and x_7 are in the same region, we consider two cases.

(a) Both x_1 and x_7 are in F_3 . Then x_6x_7 must intersect x_3x_5 by Jordan Curve Theorem. If x_3x_5 has at least two crossings, then it is deleted. Considering that the obtained graph has a subgraph isomorphic to a subdivision of H_3 , we have that $k \geq 7$. If x_3x_5 has exactly one crossing, then it is produced by x_3x_5 and x_6x_7 . We now consider the vertex x_8 . We claim that x_8 must be in F_2 . For, if x_8 is in F_1 , then x_8x_5 must cross some edge in the cycle $x_3x_4x_6x_3$ by Jordan Curve Theorem. If x_8 is in F_4 , then the path $x_8y_2x_4$ must cross some edge in the cycle $x_3x_5x_6x_3$. If x_8 is in F_3 , then the edge x_8x_6 must cross some edge in the cycle $x_3x_4x_5x_3$. So x_8 is in F_2 .

We now discuss how many crossings are eliminated after x_7x_8 has been deleted to obtain G_2'' . If there are at least two crossings being eliminated, then Ψ_2' has at most $k - 3$ crossings. In this case, we delete x_3x_5 . Considering that the obtained graph has a subgraph isomorphic to a subdivision of H_3 defined in Lemma 2.4, we have that $k \geq 7$. If there is exactly one crossing being eliminated, then the crossing must be produced by x_7x_8 and x_4x_5 . Now x_7x_8 is added back in the primitive way. Next, x_7x_6 is newly drawn such that it is near to the path $x_6x_8x_7$. We now delete all edges incident with x_8 other than x_8x_6 and x_8x_7 in the interior of the region F_2 . Then x_6x_7 has exactly one crossing which is produced by x_6x_7 and x_4x_5 . Notice that x_4x_5 has at least two crossings in this case. Next, x_4x_5 is deleted. The obtained drawing has at most $k - 3$ crossings, and the obtained graph has a subgraph isomorphic to a subdivision of H_5 defined in Lemma 2.6. Hence, we have $k \geq 7$.

(b) Both x_1 and x_7 are in F_4 . Then x_1x_4 must intersect x_3x_5 . Next, we proceed the similar argument as x_6x_7 . The difference are that x_8 is replaced by x_2 , that x_7x_8 is replaced by x_1x_2 , and that F_2 is replaced by F_1 . So $k \geq 7$. □

Lemma 2.9 *Let G_3 be the graph shown in Figure 10. Then $cr(G) \geq 8$.*

Proof Suppose that $cr(G_3) = k$, and that Ψ is an optimal drawing of G_3 .

Let Q be the subgraph of G_3 induced by the four vertices x_1, x_2, x_3 and x_4 . Obviously, Q is isomorphic to K_4 , and G_3 has a subgraph isomorphic to a subdivision of T_6 which contains Q . By Lemma 2.1, some edge in $E(Q)$ has at least one crossing in Ψ .

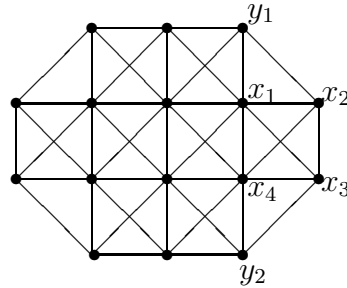


Figure 10 The graph G_3

If some edge in $\{x_1x_2, x_1x_3, x_2x_4, x_3x_4\}$ has at least one crossing in Ψ , then the edge is deleted. The obtained graph has a subgraph isomorphic to a subdivision of G_1 defined in Lemma 2.7. So $k \geq 8$. Otherwise, we consider x_2x_3 . If x_2x_3 has at least one crossing in Ψ , then it is deleted. The obtained graph is isomorphic to G_2 defined in Lemma 2.8. Thus, we have that $k \geq 8$. If not, then x_1x_4 has at least one crossing in Ψ . If we can redraw x_1x_4 so that it is not crossed, then we obtain a drawing of G_3 with less crossings than in Ψ , a contradiction. Otherwise, x_1x_4 can be redrawn such that it crosses exactly x_2x_3 . Since the graph obtained from G_3 by deleting x_2x_3 is isomorphic to G_2 , we have $k \geq 8$. □

3 The crossing number of $P_m \boxtimes P_n$ for $m \geq 4, n \geq 4$ and $(m, n) \neq (4, 4)$

Let $W_{m,n}$ be the graph obtained from $P_m \boxtimes P_n$ by deleting the four vertices $w_{1,1}, w_{1,n}, w_{m,1}$ and $w_{m,n}$.

Lemma 3.1 $cr(W_{4,n}) \geq 3(n - 1) - 4$ for $n \geq 5$.

Proof We use the induction on n . If $n = 5$, then $cr(W_{4,5}) \geq 8$, since $W_{4,5}$ is isomorphic to the graph defined in Lemma 2.9.

Assume that $cr(W_{4,t}) \geq 3(t - 1) - 4$, where $t \geq 5$. Suppose that $cr(W_{4,t+1}) = k$, and that Π is an optimal drawing of $W_{4,t+1}$. Let Q_1 be the subgraph of $W_{4,t+1}$ induced by the four vertices $w_{1,t-1}, w_{2,t-1}, w_{1,t}$ and $w_{2,t}$. Obviously, Q_1 is isomorphic to K_4 , and $W_{4,t+1}$ has a subgraph isomorphic to a subdivision of T_6 which contains Q_1 . By Lemma 2.1, some edge in Q_1 has at least one crossing in Π . We now apply the following operations which are similar to that in the proof of Lemma 2.7.

- (1) If some edge in $E(Q_1) \setminus \{w_{1,t-1}w_{2,t-1}, w_{2,t-1}w_{2,t}\}$ has at least one crossing in Π ,

then it is deleted and turn to (4). Otherwise, turn to (2).

(2) Suppose that the edge $w_{1,t-1}w_{2,t-1}$ has at least one crossing in Π . We can redraw $w_{1,t-1}w_{2,t-1}$ near to the path $w_{1,t-1}w_{1,t}w_{2,t-1}$ such that it has at most one crossing. If the new drawing of $w_{1,t-1}w_{2,t-1}$ has no crossing, then a drawing of $W_{4,t+1}$ with at most $k - 1$ crossings is obtained, a contradiction. Otherwise, $w_{1,t-1}w_{2,t-1}$ can be redrawn near to $w_{1,t-1}w_{2,t-1}$ such that it crosses exactly $w_{1,t}w_{2,t}$. Delete $w_{1,t}w_{2,t}$ and turn to (4). If $w_{1,t-1}w_{2,t-1}$ has not any crossing in Π , turn to (3).

(3) The edge $w_{2,t-1}w_{2,t}$ has at least one crossing in Ψ . A similar argument as the one used in (2) shows that we may redraw $w_{2,t-1}w_{2,t}$ near to the path $w_{2,t-1}w_{1,t}w_{2,t}$ such that it crosses exactly $w_{1,t-1}w_{1,t}$. Delete $w_{1,t-1}w_{1,t}$ and turn to (4).

(4) Let Π'_1 be the obtained drawing, and let $W_{4,t+1}^{(1)}$ the obtained graph.

Then $\text{cr}(\Pi'_1) \leq k - 1$, and $W_{4,t+1}^{(1)}$ contains the edge $w_{1,t-1}w_{2,t}$ or the path $w_{1,t-1}w_{1,t}w_{2,t}$.

Let Q_2 be the subgraph of $W_{4,t+1}^{(1)}$ induced by the four vertices $w_{3,t-1}$, $w_{3,t}$, $w_{4,t-1}$ and $w_{4,t}$. Next, we proceed the similar argument as Q_1 . Let Π'_2 be the drawing so obtained, and let $W_{4,t+1}^{(2)}$ be the graph corresponding to Π'_2 . Then $\text{cr}(\Pi'_2) \leq \text{cr}(\Pi'_1) - 1 \leq k - 2$ and $W_{4,t+1}^{(2)}$ contains the edge $w_{4,t-1}w_{3,t}$ or the path $w_{4,t-1}w_{4,t}w_{3,t}$.

Let Q_3 be the induced subgraph of $W_{4,t+1}^{(2)}$ by the four vertices $w_{2,t}$, $w_{2,t+1}$, $w_{3,t}$ and $w_{3,t+1}$. Obviously, Q_3 is isomorphic to K_4 , and $W_{4,t+1}^{(2)}$ has a subgraph isomorphic to a subdivision of T_6 which contains Q_3 . By Lemma 2.1, some edge in $E(Q_3)$ has at least one crossing in Π'_2 . We now apply the following operations.

(a) If some edge in $E(Q_3) \setminus \{w_{2,t}w_{3,t}\}$ has at least one crossing in Π'_2 , then it is deleted and turn to (c). Otherwise, turn to (b).

(b) The edge $w_{2,t}w_{3,t}$ is crossed in Π'_2 . We redraw this edge near to the path $w_{2,t}w_{2,t+1}w_{3,t}$. If in such new drawing $w_{2,t}w_{3,t}$ is not crossed, then turn to (c). Otherwise, $w_{2,t}w_{3,t}$ can be drawn such that it crosses exactly the edge $w_{2,t+1}w_{3,t+1}$. Now the edge $w_{2,t+1}w_{3,t+1}$ is deleted, and turn to (c).

(c) Let Π'_3 be the obtained drawing, and let $W_{4,t+1}^{(3)}$ be the obtained graph.

Then $\text{cr}(W_{4,t+1}^{(3)}) \leq \text{cr}(W_{4,t+1}^{(2)}) - 1 \leq k - 3$, and $W_{4,t+1}^{(3)}$ has a subgraph isomorphic to a subdivision of $W_{4,t}$. By the inductive assumption, $\text{cr}(W_{4,t}) \geq 3(t - 1) - 4$. This implies that $k \geq 3t - 4$. So $\text{cr}(W_{4,t+1}) \geq 3t - 4$. Therefore, $\text{cr}(W_{4,n}) \geq 3(n - 1) - 4$. \square

Lemma 3.2 $\text{cr}(W_{m,n}) \geq (m - 1)(n - 1) - 4$ for $m \geq 4$ and $n \geq 5$.

Proof We use the induction on m . By Lemma 3.1, $\text{cr}(W_{m,n}) \geq (m - 1)(n - 1) - 4$ if $m = 4$. Assume that $\text{cr}(W_{m,n}) \geq (m - 1)(n - 1) - 4$ if $m = q$. Suppose that $\text{cr}(W_{q+1,n}) = k$, and that Π is an optimal drawing of $W_{q+1,n}$.

Let Q'_1 (Q'_{n-1} , respectively) be the induced subgraph of $W_{q+1,n}$ by the four vertices $w_{q-1,1}$, $w_{q-1,2}$, $w_{q,1}$ and $w_{q,2}$ ($w_{q-1,n-1}$, $w_{q-1,n}$, $w_{q,n-1}$ and $w_{q,n}$, respectively). For $i = 1, 2, \dots, n - 3$, let Q'_{i+1} be the induced subgraph of $W_{q+1,n}$ by the four vertices

$w_{q,i+1}, w_{q,i+2}, w_{q+1,i+1}$ and $w_{q+1,i+2}$. It is obvious that Q'_j is isomorphic to K_4 for $j = 1, 2, \dots, n - 1$.

We start with Q'_1 and deal with it as Q_1 in the proof of Lemma 3.1. If some edge in $E(Q'_1) \setminus \{w_{q-1,1}w_{q-1,2}, w_{q-1,2}w_{q,2}\}$ has at least one crossing in Π , then it is deleted. Otherwise, one of $w_{q-1,1}w_{q-1,2}$ and $w_{q-1,2}w_{q,2}$ has at least one crossing. In this case the edge is redrawn and some edge in $\{w_{q-1,1}w_{q,1}, w_{q,1}w_{q,2}\}$ is deleted if necessary such that at least one crossing is eliminated.

After Q'_1 has been dealt with, the obtained graph has a subgraph isomorphic to a subdivision of T_6 which contains Q'_2 . By Lemma 2.1, some edge in Q'_2 has at least one crossing in the present drawing. We now deal with Q'_2 in the similar way to that of Q_1 in the proof of Lemma 3.1. If some edge in $E(Q'_2) \setminus \{w_{q,2}w_{q,3}, w_{q,3}w_{q+1,3}\}$ has at least one crossing in the present drawing, then it is deleted. Otherwise, one of $w_{q,2}w_{q,3}$ and $w_{q,3}w_{q+1,3}$ has at least one crossing in the present drawing. In this case the edge is redrawn and some edge in $\{w_{q,2}w_{q+1,2}, w_{q+1,2}w_{q+1,3}\}$ is deleted if necessary such that at least one crossing is eliminated.

For $i = 3, \dots, n - 2$, Q'_i is dealt with as Q'_2 . At last, Q'_{n-1} is dealt with in the similar way to Q_3 in the proof of Lemma 3.1. Let G be the obtained graph after removing at least one crossing for each of $Q_i, i \in \{1, 2, \dots, n - 1\}$. Then G has a subgraph isomorphic to a subdivision of the graph $W_{q,n}$. Thus, $\text{cr}(G) \leq k - (n - 1)$. By the inductive assumption, $\text{cr}(G) \geq (q - 1)(n - 1) - 4$. This implies that $k \geq q(n - 1) - 4$. So $\text{cr}(W_{q+1,n}) \geq q(n - 1) - 4$. Therefore, $\text{cr}(W_{m,n}) \geq (m - 1)(n - 1) - 4$. \square

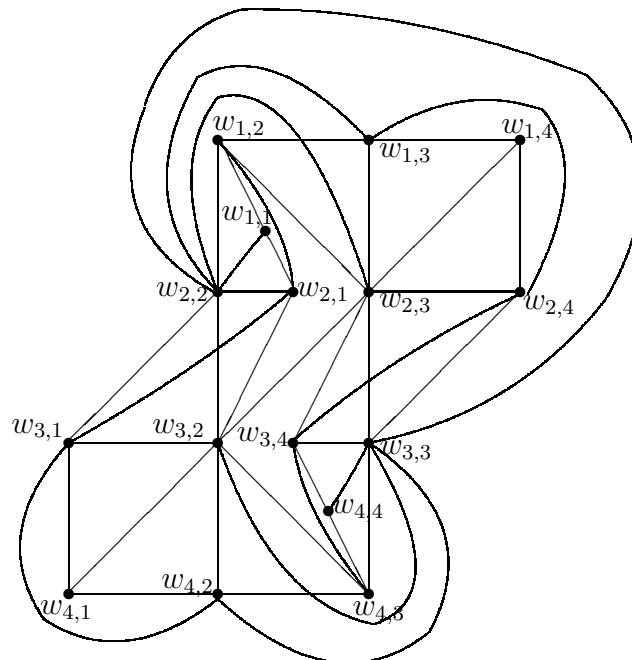


Figure 11 A drawing of $P_4 \boxtimes P_4$

Since $P_m \boxtimes P_n$ is isomorphic to $P_n \boxtimes P_m$, we have that $\text{cr}(P_m \boxtimes P_n) = \text{cr}(P_n \boxtimes P_m)$.

Theorem 3.3 $cr(P_m \boxtimes P_n) = (m - 1)(n - 1) - 4$ for $m \geq 4$, $n \geq 4$ and $(m, n) \neq (4, 4)$.

Proof The theorem follows from Lemmas 1.1 and 3.2. \square

4 The crossing number of $P_4 \boxtimes P_4$

Theorem 4.1 $cr(P_4 \boxtimes P_4) = 4$.

Proof The drawing of $P_4 \boxtimes P_4$ shown in Figure 11 implies that $cr(P_4 \boxtimes P_4) \leq 4$. Since $P_4 \boxtimes P_4$ has a subgraph isomorphic to a subdivision of H_5 defined in Lemma 2.6, $cr(P_4 \boxtimes P_4) \geq 4$. Hence $cr(P_4 \boxtimes P_4) = 4$. \square

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