# A counterexample to a result on the tree graph of a graph<sup>\*</sup>

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### Abstract

Given a set of cycles C of a graph G, the tree graph of G, defined by C, is the graph T(G, C) whose vertices are the spanning trees of G and in which two trees R and S are adjacent if  $R \cup S$  contains exactly one cycle and this cycle lies in C. Li et al. [Discrete Math. 271 (2003), 303–310] proved that if the graph T(G, C) is connected, then C cyclically spans the cycle space of G. Later, Yumei Hu [Proc. 6th Int. Conf. Wireless Communications Networking and Mobile Comput. (2010), 1–3] proved that if C is an arboreal family of cycles of G which cyclically spans the cycle space of a 2-connected graph G, then T(G, C) is connected. In this note we present an infinite family of counterexamples to Hu's result.

## 1 Introduction

The tree graph of a connected graph G is the graph T(G) whose vertices are the spanning trees of G, in which two trees R and S are adjacent if  $R \cup S$  contains exactly one cycle. Li et al. [2] defined the tree graph of G with respect to a set of

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cycles C as the spanning subgraph T(G, C) of T(G) where two trees R and S are adjacent only if the unique cycle contained in  $R \cup S$  lies in C.

A set of cycles C of G cyclically spans the cycle space of G if for each cycle  $\sigma$  of G there are cycles  $\alpha_1, \alpha_2, \ldots, \alpha_m \in C$  such that:  $\sigma = \alpha_1 \Delta \alpha_2 \Delta \ldots \Delta \alpha_m$  and, for  $i = 2, 3, \ldots, m, \alpha_1 \Delta \alpha_2 \Delta \ldots \Delta \alpha_i$  is a cycle of G. Li et al. [2] proved the following theorem:

**Theorem 1.** If C is a set of cycles of a connected graph G such that the graph T(G,C) is connected, then C cyclically spans the cycle space of G.

A set of cycles C of a graph G is *arboreal* with respect to G if for every spanning tree T of G, there is a cycle  $\sigma \in C$  which is a fundamental cycle of T. Yumei Hu [1] claimed to have proved the converse theorem:

**Theorem 2.** Let G be a 2-connected graph. If C is an arboreal set of cycles of G that cyclically spans the cycle space of G, then T(G, C) is connected.

In this note we present a counterexample to Theorem 2 given by a triangulated plane graph G with 6 vertices and an arboreal family of cycles C of G such that Ccyclically spans the cycle space of G, while T(G, C) is disconnected. Our example generalises to a family of triangulated graphs  $G_n$  with 3(n + 2) vertices for each integer  $n \ge 0$ .

If  $\alpha$  is a face of a plane graph G, we denote, also by  $\alpha$ , the corresponding cycle of G as well as the set of edges of  $\alpha$ .

#### 2 Preliminary results

Let G be a plane graph. For each cycle  $\tau$ , let  $k(\tau)$  be the number of faces of G contained in the interior of  $\tau$ . A *diagonal edge* of  $\tau$  is an edge lying in the interior of  $\tau$  having both vertices in  $\tau$ . The following lemma will be used in the proof of Theorem 4.

**Lemma 3.** Let G be a triangulated plane graph and  $\sigma$  be a cycle of G. If  $k(\sigma) \geq 2$ , then there are two faces  $\phi$  and  $\psi$  of G, contained in the interior of  $\sigma$ , both with at least one edge in common with  $\sigma$ , and such that  $\sigma \Delta \phi$  and  $\sigma \Delta \psi$  are cycles of G.

*Proof.* If  $k(\sigma) = 2$ , let  $\phi$  and  $\psi$  be the two faces of G contained in the interior of  $\sigma$ . Clearly  $\sigma \Delta \phi = \psi$  and  $\sigma \Delta \psi = \phi$  which are cycles of G.

Assume  $k = k(\sigma) \geq 3$  and that the result holds for each cycle  $\tau$  of G with  $2 \leq k(\tau) < k$ . If  $\sigma$  has a diagonal edge uv, then  $\sigma$  together with the edge uv define two cycles  $\sigma_1$  and  $\sigma_2$  such that  $k(\sigma) = k(\sigma_1) + k(\sigma_2)$ . If  $\sigma_1$  is a face of G, then  $\sigma \Delta \sigma_1$  is a cycle of G and, if  $k(\sigma_1) \geq 2$ , then by induction there are two faces  $\phi_1$  and  $\psi_1$  of G, contained in the interior of  $\sigma_1$ , both with at least one edge in common with  $\sigma_1$ , and such that  $\sigma_1 \Delta \phi_1$  and  $\sigma_1 \Delta \psi_1$  are cycles of G. Without loss of generality, we assume uv is not an edge of  $\phi_1$  and therefore  $\phi = \phi_1$  has at least one edge in common

with  $\sigma$  and is such that  $\sigma\Delta\phi$  is a cycle of G. Analogously  $\sigma_2$  contains a face  $\psi$  with at least one edge in common with  $\sigma$  and such that  $\sigma\Delta\psi$  is a cycle of G.

For the remaining of the proof we may assume  $\sigma$  has no diagonal edges. Since  $k = k(\sigma) \geq 3$ , there is a vertex u of  $\sigma$  which is incident with one or more edges lying in the interior of  $\sigma$ . Let  $v_0, v_1, \ldots, v_m$  be the vertices in  $\sigma$  or in the interior of  $\sigma$  which are adjacent to u. Without loss of generality we assume  $v_0$  and  $v_m$  are vertices of  $\sigma$  and that  $v_0, v_1, \ldots, v_m$  is a path joining  $v_0$  and  $v_m$ , see Figure 1.



Figure 1: Cycle  $\sigma$  with no diagonal edges.

As  $\sigma$  has no diagonal edges, vertices  $v_1, v_2, \ldots, v_{m-1}$  are not vertices of  $\sigma$  and therefore faces  $\phi = uv_0v_1$  and  $\psi = uv_mv_{m-1}$  are such that  $\sigma\Delta\phi$  and  $\sigma\Delta\psi$  are cycles of G, each with one edge in common with  $\sigma$ .

**Theorem 4.** Let G be a triangulated plane graph and  $\alpha$  and  $\beta$  be two internal faces of G with one edge in common. If C is the set of internal faces of G with cycle  $\alpha$ replaced by the cycle  $\alpha \Delta \beta$ , then C cyclically spans the cycle space of G.

*Proof.* Let  $\sigma$  be a cycle of G. If  $k(\sigma) = 1$ , then  $\sigma \in C$  or  $\sigma = \alpha$  in which case  $\sigma = (\alpha \Delta \beta) \Delta \beta$ . In both cases  $\sigma$  is cyclically spanned by C.

We proceed by induction assuming  $k = k(\sigma) \ge 2$  and that if  $\tau$  is a cycle of G with  $k(\tau) < k$ , then  $\tau$  is cyclically spanned by C.

By Lemma 3, there are two faces  $\phi$  and  $\psi$  of G, contained in the interior of  $\sigma$ , such that both  $\sigma\Delta\phi$  and  $\sigma\Delta\psi$  are cycles of G; without loss of generality we assume  $\phi \neq \alpha$ . Clearly  $k(\sigma\Delta\phi) < k$ ; by induction, there are cycles  $\tau_1, \tau_2, \ldots, \tau_m \in C$  such that:  $\sigma\Delta\phi = \tau_1\Delta\tau_2\Delta\ldots\Delta\tau_m$  and, for  $i = 2, 3, \ldots, m, \tau_1\Delta\tau_2\Delta\ldots\Delta\tau_i$  is a cycle of G. As  $\sigma = (\sigma\Delta\phi)\Delta\phi = \tau_1\Delta\tau_2\Delta\ldots\Delta\tau_m\Delta\phi$ , cycle  $\sigma$  is also cyclically spanned by C.  $\Box$ 

#### 3 Main result

Let G be the skeleton graph of a octahedron (see Figure 2) and C be the set of cycles that correspond to the internal faces of G with cycle  $\alpha$  replaced by cycle  $\alpha \Delta \beta$ .



Figure 2: Graph G with internal faces  $\alpha, \beta, \gamma, \delta, \sigma, \tau$  and  $\rho$  and outer face  $\omega$ .

By Theorem 4, C cyclically spans the cycle space of G. Suppose C is not arboreal and let P be a spanning tree of G with none of its fundamental cycles in C. For this to happen, each of the cycles  $\beta, \gamma, \delta, \sigma, \tau$  and  $\rho$  of G, must have at least two edges which are not edges of P and since P has no cycles, at least one edge of cycle  $\alpha$  and at least one edge of cycle  $\omega$  are not edgs of P.

Therefore G has at least 7 edges which are not edges of P. These, together with the 5 edges of P sum up to 12 edges which is exactly the number of edges of G. This implies that each of the cycles  $\omega$  and  $\alpha$  has exactly two edges of P and that each of the cycles  $\beta$ ,  $\gamma$ ,  $\delta$ ,  $\sigma$ ,  $\tau$  and  $\rho$  has exactly one edge of P.

If edges xy and xz are edges of P, then vertex x cannot be incident to any other edge of P and therefore cycle  $\alpha$  can only have one edge of P, which is not possible.

If edges xy and yz are edges of P, then vertex y cannot be incident to any other edge of P. In this case, the edge in cycle  $\sigma$ , opposite to vertex y, must be an edge of P and cannot be incident to any other edge of P, which again is not possible. The case where edges xz and yz are edges of P is analogous. Therefore C is an arboreal set of cycles of G.

Let T, S and R be the spanning trees of G given in Figure 3. The graph T(G, C) has a connected component formed by the trees T, S and R since cycle  $\rho$  is the only cycle in C which is a fundamental cycle of either T, S or R. This implies that T(G, C) is disconnected.



Figure 3: Trees T (left), S (center) and R (right).

We proceed to generalise the counterexample to graphs with arbitrary large number of vertices. Let  $G_0 = G, x_0 = x, y_0 = y, z_0 = z$  and for  $t \ge 0$  define  $G_{t+1}$  as the graph obtained by placing a copy of  $G_t$  in the inner face of the skeleton graph of an octahedron as in Figure 4.



Figure 4:  $G_{t+1}$ .

Notice that each graph  $G_n$  contains a copy G' of G in the innermost layer. We also denote by  $\alpha, \beta, \gamma, \delta, \sigma, \tau$  and  $\rho$  the cycles of  $G_n$  that correspond to the cycles  $\alpha, \beta, \gamma, \delta, \sigma, \tau$  and  $\rho$  of G'. Let  $\omega_n$  denote the cycle given by the edges in the outer face of  $G_n$ .

For  $n \ge 0$  let  $C_n$  be the set of cycles that correspond to the internal faces of  $G_n$  with cycle  $\alpha$  replaced by cycle  $\alpha \Delta \beta$ . By Theorem 4,  $C_n$  cyclically spans the cycle space of  $G_n$ .

We claim that for  $n \ge 0$ , set  $C_n$  is an arboreal set of cycles of  $G_n$ . Suppose  $C_t$  is arboreal but  $C_{t+1}$  is not and let  $P_{t+1}$  be a spanning tree of  $G_{t+1}$  such that none of its fundamental cycles lies in  $C_{t+1}$ .

As in the case of graph G and tree P, above, each cycle in  $C_{t+1}$ , other than  $\alpha \Delta \beta$ , has exactly one edge in  $P_{t+1}$ , while cycles  $\alpha$  and  $\omega_{t+1}$  have exactly two edges in  $P_{t+1}$ . The reader can see that this implies that the edges of  $P_{t+1}$  which are not edges of  $G_t$ form a path with length 3. Then  $P_{t+1} - \{x_{t+1}, y_{t+1}, z_{t+1}\}$  is a spanning tree of  $G_t$  and by induction, one of its fundamental cycles lies in  $C_t \subset C_{t+1}$  which is a contradiction. Therefore  $C_{t+1}$  is arboreal.

Let  $T_0 = T$  and for  $t \ge 0$  define  $T_{t+1}$  as the spanning tree of  $G_{t+1}$  obtained by placing a copy of  $T_t$  in the inner face of the skeleton graph of an octahedron as in Figure 5.

Trees  $S_{t+1}$  and  $R_{t+1}$  are obtained from  $S_t$  and  $R_t$  in the same way with  $S_0 = S$ and  $R_0 = R$  respectively. We claim that, for each integer  $n \ge 0$ , cycle  $\rho$  is the only cycle in  $C_n$  which is a fundamental cycle of either  $T_n, S_n$  or  $R_n$ . Therefore  $T_n, S_n$  and  $R_n$  form a connected component of  $G_n$  which implies that  $T(G_n, C_n)$  is disconnected.



Figure 5: t = 3k (left), t = 3k + 1 (centre) and t = 3k + 2 (right).

## References

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