# On the number of tilings of the rectangular board with T-tetrominoes 

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#### Abstract

The classical combinatorial problem of counting domino tilings of a $2 n \times 2 m$ rectangle was solved by P.W. Kasteleyn and also by H.N.V. Temperley and M.E. Fisher in 1961. We shall consider the similar problem for T-tetrominoes, that is, pieces formed by 4 unit squares in the shape of a T . We give explicit formulae for the number of tilings with T-tetrominoes for the $4 n \times 4 m$ rectangle when $n=1,2,3$ and 4 , and a computational method for values of $n$ up to 8 .


## 1 Introduction

One of the classical results in enumerative combinatorics is the formula for the number of domino tilings of a $2 n \times 2 m$ chess-board. The formula is the following:

$$
4^{n m} \prod_{s=1}^{n} \prod_{t=1}^{m}\left(\cos ^{2} \frac{s \pi}{2 n+1}+\cos ^{2} \frac{t \pi}{2 m+1}\right) .
$$

This was obtained originally by Kasteleyn [7], and also by Temperley and Fisher [14]. Later Lieb [10] and more recently V. Strehl [13] gave different methods to obtained the same formula.

In order to state our problem we give some definitions first. A tetromino is a twodimensional shape made by connecting 4 unit squares along their edges. There are 5 possible tetrominoes (up to rotations and reflections) of which one has a T-shape, and we called it $T$-tetromino. A tiling of a plane region $R$ is a covering of $R$ using a given set of tiles, completely and without any overlap. We have the following:

[^0]Theorem 1 (D.W. Walkup 1965). An $r \times s$ rectangular board has a tiling with $T$-tetrominoes as tiles if and only if $r$ and $s$ are multiples of 4.

Our problem is to count the number of tilings of a $4 n \times 4 m$ rectangular board with T-tetrominoes or $T$-tilings, which we denote by $f(n, m)$. We use $f(n)$ for the number of T-tilings of the square board.

The rectangular lattice $L_{n, m}$ is a well-known graph (see Biggs [2]), but for reference we include the definition here. The graph $L_{n, m}$ has as vertices the set $\{0, \ldots, n-1\} \times\{0, \ldots, m-1\}$ where two vertices $(i, j)$ and $\left(i^{\prime}, j^{\prime}\right)$ are adjacent if $\left|i-i^{\prime}\right|+\left|j-j^{\prime}\right|=1$. This definition also gives a natural planar embedding of $L_{n, m}$, in which each bounded face is a unit square. We consider the region defined by this planar embedding of $L_{n+1, m+1}$ as the $n \times m$ rectangular board. For the basic graph theory required in this paper the reader is referred to the book of Bondy and Murty [3].

## 2 T-tilings and The Tutte polynomial

For a graph $G=(V, E)$, we identify a subset of edges $A \subseteq E$ of $G$ with the subgraph of $G$ that $A$ spans, that is, the spanning subgraph $(V, A)$. Then the number of connected components of $A$ is denoted by $k(A)$ and the rank of $A$ is defined by $r(A)=|V|-k(A)$. When we have a planar embedding of $G$, any subgraph of $G$ has an induced planar embedding. In this case, we defined the number $f(A)$, as the number of bounded faces of $A$.

The Tutte polynomial $T(G ; x, y)$ of a graph $G=(V, E)$ is a two variable polynomial and has the following expansion

$$
\begin{equation*}
T(G ; x, y)=\sum_{A \subseteq E}(x-1)^{r(E)-r(A)}(y-1)^{|A|-r(A)} . \tag{2.1}
\end{equation*}
$$

If $G$ is a connected plane graph, $r(E)-r(A)=k(A)-1$ and $|A|-r(A)=$ $f(A)$. The last equality is obtained by applying Euler's formula to each connected component of $A$.

When evaluating the Tutte polynomial along different curves and points we get several interesting invariants of graphs. Among them we have the chromatic and flow polynomials of a graph, the all terminal reliability probability of a network and the partition function of the $Q$-state Potts model. Further details of many of the invariants can be found in the survey article of Brylawski and Oxley [4] and in the book of Welsh [17].

Theorem 2 (Korn and Pak, 2003). The number of T-tilings of a $4 n \times 4 m$ board satisfy

$$
\begin{equation*}
f(n, m)=2 T\left(L_{n, m} ; 3,3\right) . \tag{2.2}
\end{equation*}
$$

Sketch. We defined the top of a T-tetromino as the line segment of length 3 on its boundary and its bottom as the unit line segment on the boundary opposite to the top.

Consider the lattice induced by the points $(a, b)$ such that $a, b \equiv 2(\bmod 4)$ in a $4 n \times 4 m$ board. We obtain a graph isomorphic to $L_{n, m}$, where each edge has length 4. Now, let us take any subset of edges $A$ in $L_{n, m}$ and let $C$ be the set of boundary points of a bounded (unbounded) face $F$ in $A$. Then, let $R$ be the region inside (outside) $F$ made of by all the points at Chebyshev distance at most 2 from $C$. Here the Chebyshev distance is the one induced by the $l^{\infty}$-norm. The left-hand side of Figure 1 shows a set $A$ of edges of $L_{6,5}$ with 4 bounded faces and 3 unbounded faces.


Figure 1: The figure on the left shows a partial T-tiling of a 24 by 20 rectangular board. The figure on the right shows a complete T-tiling of the board.

In general, the region $R$ can be tiled with T-tetrominoes in many different ways. But there are exactly two T-tilings of $R$ where the top and bottom of the tiles alternatively intersect the set $C$. An example is shown on the left-hand side of Figure 1.

For a fixed $A$, the union of these regions is the $4 n \times 4 m$ rectangular board. Thus, to each subset $A$ of edges of the lattice $L_{n, m}$ we have associated $2^{f(A)+k(A)}$ T-tilings.

Clearly, from different edge-sets we get different tilings. Also, each tiling can be obtained b this procedure from an edge-set $A$ and a choice of T-tilings o each region. Summing up these quantities for all the subsets of edges in $L_{n, m}$ we get that the number of T-tilings is

$$
\begin{aligned}
f(n, m) & =\sum_{A \subseteq E\left(L_{n, m}\right)} 2^{k(A)+f(A)} \\
& =2 \sum_{A \subseteq E\left(L_{n, m}\right)} 2^{k(A)-1} 2^{f(A)} \\
& =2 \sum_{A \subseteq E\left(L_{n, m}\right)} 2^{r(E)-r(A)} 2^{|A|-r(A)}=2 T\left(L_{n, m} ; 3,3\right) .
\end{aligned}
$$

For example, the Tutte polynomial of the square lattice $L_{3,3}$ is

$$
\begin{aligned}
T\left(L_{3,3} ; x, y\right)= & x^{8}+4 x^{7}+10 x^{6}+16 x^{5}+19 x^{4}+16 x^{3}+10 x^{2}+3 x \\
& +4 x^{5} y+12 x^{4} y+20 x^{3} y+20 x^{2} y+13 x y+3 y \\
& +4 x^{3} y^{2}+10 x^{2} y^{2}+12 x y^{2}+6 y^{2}+4 x y^{3}+4 y^{3}+y^{4}
\end{aligned}
$$

and so, the number of T-tilings of a 12-by-12 board is 78696 .

## 3 The number of T-tilings

Using Equation 2.2, we can easily compute the number of T-tilings for the $4 \times 4 m$ board, as the corresponding lattice $L_{1, m}$ is a path of length $m-1$, whose Tutte polynomial is just $x^{m-1}$. Thus, for the 4 x 4 board we get $f(1,1)=2$ and, in general, we have that

$$
f(1, m)=2(3)^{m-1}
$$

This sequence is already known as A008776 [11] and could have been obtained more easily by direct counting. However, direct counting is not so straightforward for the $8 \times 4 m$ board. But the corresponding lattice in Theorem 2 is a ladder, and it is easy to find the recurrence relation $f(2, m)=16 f(2, m-1)-27 f(2, m-2)$, for $m \geq 3$, with initial conditions $f(2,1)=6$ and $f(2,2)=84$. After solving we get the following formula:

$$
\begin{equation*}
f(2, m)=\left(3-\frac{18}{\sqrt{37}}\right)(8-\sqrt{37})^{m-1}+\left(3+\frac{18}{\sqrt{37}}\right)(8+\sqrt{37})^{m-1} \tag{3.1}
\end{equation*}
$$

For the board of width 12 we have to evaluate $T\left(L_{3, m} ; 3,3\right)$ that could still be done by finding a recurrence relation and solving it. This method is going to give a formula for the boards of fixed width $n$ for small $n$; but it is bound to fail for the general case. There is another way to compute $T\left(L_{n, m} ; x, y\right)$ for a fixed width $n$ at point ( $x, y$ ) which is described in Calkin et al. [6] and it is an application of the transfer-matrix method. In this case we have the same restriction as before, a fixed width $n$ for small values of $n$, but it has the advantage of being easily automatized.

Theorem 3 (Calkin et al. 2001). For real values $x$ and $y$ with $x \neq 0$ and integers $n, m \geq 2$, we have

$$
T\left(L_{n, m} ; x+1, y+1\right)=x^{n m-1} X_{n}^{t} \cdot\left(\Lambda_{n}\right)^{m-1} \cdot \overrightarrow{1}
$$

where $X_{n}$, a vector of length $c_{n}$, and $\Lambda_{n}, a c_{n} \times c_{n}$ matrix, depend on $x, y$ and $n$ but not $m$. And $\overrightarrow{1}$ is the vector of length $c_{n}$ with all entries equal to 1 .

The quantity $c_{n}$ is the $n$-Catalan number, so the method is just practical for small values of $n$. Computing the vectors $X_{n}$ and the matrix $\Lambda_{n}$ can be easily done in a computer. Following the example in Calkin et al. [6] we get

$$
T\left(L_{2, m} ; 3,3\right)=2^{2 m-1}(1, \quad 1 / 2)\left(\begin{array}{cc}
9 / 4 & 9 / 8 \\
2 & 7 / 2
\end{array}\right)^{m-1}\binom{1}{1}
$$

In the above equation we can get integer entries by multiplying by an appropriate power of 2. Thus, by using Equation 2.2 we get the formula

$$
f(2, m)=\frac{1}{2^{m-2}}(2, \quad 1)\left(\begin{array}{cc}
18 & 9 \\
16 & 14
\end{array}\right)^{m-1}\binom{1}{1} .
$$

The method applied to the rectangular lattice $L_{3, m}$ gives

$$
f(3, m)=\frac{1}{2^{2 m-4}}(4, \quad 2, \quad 0, \quad 2, \quad 1)\left(\begin{array}{ccccc}
108 & 54 & 0 & 54 & 27 \\
96 & 84 & 0 & 48 & 42 \\
96 & 48 & 12 & 48 & 48 \\
96 & 48 & 0 & 84 & 42 \\
80 & 64 & 8 & 64 & 96
\end{array}\right)^{m-1}\left(\begin{array}{l}
1 \\
1 \\
1 \\
1 \\
1
\end{array}\right)
$$

The first 5 terms of the sequence are 18, 1182, 78696, 5253822, 350950482.
In general, we have the formula

$$
f(n, m)=\frac{1}{2^{(n-1)(m-1)}} \vec{v}_{n} A_{n}^{m-1} \overrightarrow{1},
$$

where $A_{m}=2^{2 m-1} \Lambda_{n}$ and $\vec{v}_{n}=2^{m-1} X_{n}$. For $n=4$ the matrix $A_{4}$ is the 14 -by- 14 matrix

$$
\left(\begin{array}{cccccccccccccc}
648 & 324 & 0 & 0 & 324 & 162 & 0 & 0 & 0 & 324 & 162 & 0 & 162 & 81 \\
576 & 504 & 0 & 0 & 288 & 252 & 0 & 0 & 0 & 288 & 252 & 0 & 144 & 126 \\
576 & 288 & 72 & 0 & 288 & 288 & 0 & 0 & 0 & 288 & 144 & 36 & 144 & 144 \\
576 & 288 & 0 & 72 & 288 & 144 & 36 & 0 & 36 & 288 & 144 & 36 & 144 & 162 \\
576 & 288 & 0 & 0 & 504 & 252 & 0 & 0 & 0 & 288 & 144 & 0 & 252 & 126 \\
480 & 384 & 48 & 0 & 384 & 576 & 0 & 0 & 0 & 240 & 192 & 24 & 192 & 288 \\
512 & 256 & 0 & 64 & 448 & 224 & 56 & 0 & 32 & 256 & 128 & 32 & 224 & 240 \\
576 & 288 & 0 & 0 & 288 & 144 & 0 & 72 & 36 & 288 & 144 & 0 & 288 & 144 \\
480 & 384 & 0 & 48 & 240 & 192 & 24 & 48 & 120 & 240 & 192 & 24 & 216 & 348 \\
576 & 288 & 0 & 0 & 288 & 144 & 0 & 0 & 0 & 504 & 252 & 0 & 252 & 126 \\
512 & 448 & 0 & 0 & 256 & 224 & 0 & 0 & 0 & 448 & 392 & 0 & 224 & 196 \\
480 & 240 & 48 & 48 & 240 & 216 & 24 & 0 & 24 & 384 & 192 & 120 & 192 & 348 \\
480 & 240 & 0 & 0 & 384 & 192 & 0 & 48 & 24 & 384 & 192 & 0 & 576 & 288 \\
384 & 288 & 32 & 32 & 288 & 400 & 16 & 32 & 80 & 288 & 192 & 80 & 400 & 776
\end{array}\right)
$$

and the vector $\vec{v}_{4}$ is $(8,4,0,0,4,2,0,0,0,4,2,0,2,1)$. The first 5 terms of the sequence are $54,16644,5253822,1668091536,530454033510$.

The vectors and matrices for $n$ up to 8 can be found in http://www.matem. unam.mx/~merino/e_publications.html. All the matrices are nonnegative real matrices, in fact all are nonnegative primitive matrices. So in principle, we can get a formula similar to Equation 3.1 that involves the eigenvalues of the matrix. Also, the asymptotic behaviour can be obtained by computing the spectral radius of the aforementioned matrix $\Lambda_{n}$ and then taking the $n$-root. Thus, for $n=2,3$ and 4 we have

$$
\begin{aligned}
\lim _{n \rightarrow \infty} f(2, m)^{\frac{1}{2 m}} & =\sqrt{8+\sqrt{37}}=3.75270 \ldots \\
\lim _{n \rightarrow \infty} f(3, m)^{\frac{1}{3 m}} & =4.05769 \ldots \\
\lim _{n \rightarrow \infty} f(4, m)^{\frac{1}{4 m}} & =4.22351 \ldots
\end{aligned}
$$

The values of the limits for $n$ from 5 to 8 are respectively $4.32788 \ldots, 4.39966 \ldots$, 4.45208... and 4.49199....

## 4 The square board

To compute $f(n)$ we used the above method and obtained the values for $n$ up to 8 , that is, for the square board of side 32 . There are also computer programs which compute the Tutte polynomial of moderated size graphs. One is TuLiC by Rodolfo Conde which is freely available at http://ada.fciencias.unam.mx/~rconde/tulic/. Using this program we computed the Tutte polynomial for all the square lattices from $L_{2,2}$ to $L_{12,12}$. The values for the number of T-tilings are given in the following list:

$$
\begin{aligned}
f(1) & =2 \\
f(2) & =84 \\
f(3) & =78696 \\
f(4) & =1668091536 \\
f(5) & =804175873700640 \\
f(6) & =8840889502844537044800 \\
f(7) & =2219885416449546846322852561536 \\
f(8) & =12743498392347171159734108119436194009344 \\
f(9) & =1673655934365810075982323780364346176451059139240448 \\
f(10) & =5031230898942160933982250013114536314591579141675092922832 \\
& 491520 \\
f(11)= & 3463131467175532420109700259258995693571631028572354110874 \\
& 72918271870994016256 \\
f(12)= & 5459704314008539845990741829270093470430155282377357168707 \\
& 98525852987071373589503117526536192
\end{aligned}
$$

The values up to $n=8$ were checked with both methods. All these numbers have the form $2^{n} q$, for an odd integer $q$. This is true in general.

Proposition 4. The number $f(n, m)=2^{\operatorname{gcd}(n, m)} q$, where $q$ is an odd integer.
Proof. Las Vergnas [9] proved that if $G$ is a plane graph with medial graph $H$, then $T(G ; 3,3)=2^{c(H)-1} q$, where $q$ is an odd integer. Given a connected plane graph $G$, its medial graph $H$ is constructed by putting a vertex on each edge of $G$ and two vertices of $H$ are joined by an edge if the corresponding edges in $G$ are neighbours in the cyclic order of edges around a vertex. Thus, $H$ is an Eulerian 4-regular graph. The medial graph of $L_{9,3}$ is shown in Figure 2. The graph invariant $c(H)$ is the number of crossing circuits of $H$, see Las Vergnas [9], that is, the number of circuits in the Eulerian partition of the edges of $H$ defined by choosing at each vertex opposite edges to be on the same circuit. An example of one crossing circuit for the medial graph of $L_{9,3}$ is shown in Figure 2.

In the case of $L_{n, m}$, a crossing circuit of its medial graph corresponds to the trajectory of a ball thrown at $45^{\circ}$ in a billiard table of $2 n$ by $2 m$ from a point with integer coordinates $(0, j)$, for some odd integer $j$ with $1 \leq j \leq 2 m-1$. The $2 n+2 m$
points in the boundary with integer coordinates $(i, j)$ such that $i+j$ is an odd integer are partitioned by the trajectories. The points in the same equivalent class will be those that correspond to rebounds of the ball of a particular trajectory. The same method as for the classical billiard problem, see Steinhaus [12], gives that the number of rebounds in a trajectory is $2 r+2 s$, where $r$ and $s$ are the integers satisfying the equation $\operatorname{lcm}(n, m)=r m=s n$. Thus, the number of different trajectories is $\operatorname{gcd}(n, m)$. The proposition follows from Theorem 2 .


Figure 2: The medial graph of $L_{9,3}$ where the big dots are the vertices of the graph. Also one of its 3 crossing circuits is in thick line.

Corollary 5. The number $f(n)=2^{n} q$, where $q$ is an odd integer.
The asymptotic behaviour for the square lattice is known; the value corresponds to the free energy of the 4-Potts model and the result is due to Baxter [1]:

$$
\lim _{n \rightarrow \infty} f(n)^{\frac{1}{n^{2}}}=\left(\frac{\Gamma(1 / 4)}{2 \Gamma(3 / 4)}\right)^{4}=4.78926 \ldots
$$

## 5 Conclusion

We used Korn and Pak's result [8] to obtain general formulae for the number of T-tilings of a rectangular board of $4 n \times 4 m$, when $n$ is at most 4 and $m$ is arbitrary. The same method gives formulae when $n$ is at most 8 . We also give some numerical values for $f(n)$, the number of T-tilings of a square board of side $4 n$. A generalization of Theorem 2 is given by Jacobsen in [5]. There each tile is assigned a weight that depends on its orientation and position on the board and, for a particular choice of the weights, the generating function of weighted tilings is shown to be the evaluation of the multivariate Tutte polynomial ZG $(Q, v)$.

For this approach, it is necessary to evaluate the Tutte polynomial of $L_{n, m}$ at the point $(3,3)$. Evaluating $T(G ; 3,3)$ is \#P-hard, even for planar bipartite graphs, see Vertigan and Welsh [15], so better understanding about the Tutte polynomial of the lattice is needed to extend our results. However, it is possible that using a different method a similar formula to Kasteleyn's could be found.

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