

Indices of convergence on four digraph operators*

WEIGEN YAN

Department of Mathematics
Jimei University
Xiamen 361021
P.R. China
weigenyan@263.net

FUJU ZHANG

Department of Mathematics
Xiamen University
Xiamen 361005
P.R. China

Abstract

Let D be a digraph. We denote by $L(D)$, $M(D)$, $M'(D)$ and $T(D)$ the line digraph, the middle digraph, the special middle digraph and the total digraph of D , by $k(D)$, $p(D)$ and $Diam(D)$ the index of convergence, the period and the diameter of D , respectively. Zuo (in *Acta Mathematica Applicatae Sinica*, 21(1) (1998), 144–147) proved that $k(D) - 1 \leq k(L(D)) \leq k(D) + 1$. In this paper, we prove that:

1. $\max\{k(M(D)), k(M'(D)), k(T(D)) + 1\} \leq \max\{2p(D), 2k(D) + 2\}$.
2. $k(T(D)) \leq k(M(D)) \leq k(T(D)) + 1$;
 $k(T(D)) \leq k(M'(D)) \leq k(T(D)) + 1$.
3. If there do not exist both sources and sinks in D , then
 $k(M(D)) \leq k(M'(D)) \leq k(M(D)) + 1$.
4. If D is a strongly connected digraph, then
 $\min\{k(M(D)), k(M'(D)) - 1, k(T(D))\} \geq Diam(D) + 1$.
5. If D is a primitive digraph, then
 $\max\{k(M(D)), k(M'(D)) - 1, k(T(D))\} \leq k(D) + 1$.

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1 Introduction

Throughout this paper, let D be a digraph with the vertex-set $V(D) = \{v_1, v_2, \dots, v_n\}$ and the arc-set $A(D) = \{x_1, x_2, \dots, x_m\}$. Digraphs in this paper will be allowed to have loops but not multiple arcs (arcs of the form (v_i, v_j) and (v_j, v_i) are allowed). Let C_n denote a directed cycle of length n . The directed cycle of length 1 is a loop. A **source** is a vertex with in-valency 0, and a **sink** is a vertex with out-valency 0. We suppose that the operations of matrices in this paper are Boolean operations. We use Hemminger and Beineke [1] for terminologies and notations not defined here.

In 1960, Harary and Norman [2] introduced the concept of the line digraph. For a digraph D , the **line digraph**, denoted by $L(D)$, has as its vertex-set the arc-set of D , (a, b) is an arc of $L(D)$ if and only if there are vertices u, v and w in D with $a = (u, v)$ and $b = (v, w)$. Line digraphs have been discussed in [1–8].

In 1966, Chartrand [9] introduced the concept of the total digraph. for a digraph D , the **total digraph**, denoted by $T(D)$, has its vertex-set $V(T(D)) = V(D) \cup A(D)$, there is an arc $(a, b) \in A(T(D))$ from vertex a to vertex b in $V(T(D))$ if and only if one of the following four cases holds: 1. If $a \in V(D)$ and $b \in V(D)$, then $(a, b) \in A(D)$. 2. If $a \in V(D)$ and $b \in A(D)$, then a is the tail of arc b in D . 3. If $a \in A(D)$ and $b \in V(D)$, then b is the head of arc a in D . 4. If $a \in A(D)$ and $b \in A(D)$, then the head of arc a in D is the tail of arc b in D . The total digraph has been discussed in [9–11].

In 1977, Zamfirescu [10] introduced the concept of the middle digraph. For a digraph D , the **middle digraph**, denoted by $M(D)$, has its vertex-set $V(M(D)) = V(D) \cup A(D)$, there is an arc $(a, b) \in A(M(D))$ from vertex a to vertex b in $V(M(D))$ if and only if one of the following three cases holds: 1. If $a \in V(D)$ and $b \in A(D)$, then a is the tail of arc b in D . 2. If $a \in A(D)$ and $b \in V(D)$, then b is the head of arc a in D . 3. If $a \in A(D)$ and $b \in A(D)$, then the head of arc a in D is the tail of arc b in D . The middle digraph has been discussed in [10,11].

Similarly, we can define the special middle digraph of D as follows.

Definition 1.1 For a digraph D , the **special middle digraph**, denoted by $M'(D)$, has its vertex $V(M'(D)) = V(D) \cup A(D)$. there is an arc $(a, b) \in A(M'(D))$ from vertex a to vertex b in $V(M'(D))$ if and only if one of the following three cases holds: 1. If $a \in V(D)$ and $b \in V(D)$, then $(a, b) \in A(D)$. 2. If $a \in V(D)$ and $b \in A(D)$, then a is the tail of arc b in D . 3. If $a \in A(D)$ and $b \in V(D)$, then b is the head of arc a in D .

Suppose that A is the adjacency matrix of D , whose entries are 0 or 1. Then the index of convergence and the period of D equal the index of convergence and the period of Boolean matrix A , respectively, defined as follows (see [12–14]):

Suppose that A is a Boolean matrix. Its Boolean sequence of powers is denoted by $(A^j) = I, A, A^2, \dots, A^k, \dots$. The index of convergence (say $k(A)$) and the period (say $p(A)$) of A are the least non-negative integer k and the least positive integer p such that $A^k = A^{k+p}$, respectively. By this definition, the Boolean sequence of

powers of A is as follows:

$$(A^j) = I, A, A^2, \dots, A^{k-1}, A^k, \dots, A^{k+p-1}, A^k, \dots, A^{k+p-1}, \dots$$

where $A^i \neq A^j$ if $\max(i, j) < k + p$ and $i \neq j$.

It is well known [12–14] that the digraph D is primitive if and only if D is strongly connected and the greatest common divisor of the lengths of all (elementary) directed cycles of D is 1.

In paper [7], Zuo considered relations between $k(D)$ and $k(L(D))$ and between $p(D)$ and $p(L(D))$, and obtained the following proposition.

Proposition 1.2 [7] Let D be a digraph; then:

1. $p(L(D)) = p(D)$.
2. $k(D) - 1 \leq k(L(D)) \leq k(D) + 1$.
3. If D is a primitive digraph, then $k(L(D)) = k(D) + 1$.
4. $k(L(D)) = k(D) - 1 = l$ if there exist no directed cycles in D , where l is the length of the longest directed path in D .

The above results have been proved by Zhou [15] by using a simpler method—algebraic method. In particular, by using the algebraic method, Yan and Zhang [8] obtained a stronger result than that of Proposition 1.1 as follows: If a digraph D has no sources or sinks, then $k(D) \leq k(L(D)) \leq k(D) + 1$. Moreover, if D has no sources and no sinks, then $k(L(D)) = k(D) + 1$ if there is at least one connected component of D which is not a directed cycle, and $k(L(D)) = k(D) = 0$ if every connected component of D is a directed cycle.

The following results are useful.

Proposition 1.3 [12] Let D be a digraph.

1. If D is strongly connected, then: (1). $p(D)$ equals the greatest common divisor of lengths of all directed cycles of D . (2). $\text{Diam}(L(D)) = \text{Diam}(D) + 1$ unless D is a directed cycle, where $\text{Diam}(D)$ denotes the diameter of D .
2. If D is weakly connected, then $p(D)$ equals the least common multiple of the periods of strongly connected components of D .
3. If $D_i (1 \leq i \leq c)$ are all of weakly connected components of D , then $k(D) = \max_{i=1}^c(k(D_i))$, and $p(D)$ equals the least common multiple of $p(D_1), p(D_2), \dots$, and $p(D_c)$.

The following propositions are due to Hemminger and Beineke [1] and Lin and Zhang [4–6], respectively.

Proposition 1.4 [1] Suppose that D is a digraph with no isolated vertices. Then

- (1). $L(D)$ is strongly connected if and only if D is strongly connected.
- (2). $L(D)$ is a directed cycle if and only if D is a directed cycle.

Proposition 1.5 [4–6] Let D be a digraph, with vertex-set $V(D) = \{v_1, v_2, \dots, v_n\}$, and arc-set $A(D) = \{x_1, x_2, \dots, x_m\}$, B_0 and B_1 be the following two $n \times m$ matrices: $B_0 = (b_{ij}^0), B_1 = (b_{ij}^1)$, respectively, where

$$b_{ij}^0 = \begin{cases} 1 & \text{if } v_i \text{ is the tail of arc } x_j \text{ in } D; \\ 0 & \text{otherwise.} \end{cases}$$

$$b_{ij}^1 = \begin{cases} 1 & \text{if } v_i \text{ is the head of arc } x_j \text{ in } D; \\ 0 & \text{otherwise.} \end{cases}$$

Then $A = B_0 B_1^T, A_L = B_1^T B_0, A_T = \begin{pmatrix} A & B_0 \\ B_1^T & A_L \end{pmatrix}$, where A, A_L and A_T are the adjacency matrices of $D, L(D)$ and $T(D)$, respectively, and B_1^T denotes the transpose of B_1 .

Definition 1.6 [4,12] We say that \mathbf{B}_0 and \mathbf{B}_1 in Proposition 1.5 are the **out-incidence matrix** and the **in-incidence matrix** of D , respectively.

2 Some lemmas

Similarly to Proposition 1.5, we can prove the following lemma.

Lemma 2.1 Let D be a digraph with n vertices and m arcs and let A_M and $A_{M'}$ be the adjacency matrices of $M(D)$ and $M'(D)$, respectively. Then

$$A_M = \begin{pmatrix} 0_n & B_0 \\ B_1^T & A_L \end{pmatrix}, \quad A_{M'} = \begin{pmatrix} A & B_0 \\ B_1^T & 0_m \end{pmatrix}$$

where 0_n is the $n \times n$ matrix with all entries equal zero, and B_0 and B_1 are the out-incidence matrix and the in-incidence matrix of D , respectively.

Lemma 2.2 Let D be a digraph. Then

1. $A_L^k = B_1^T A^{k-1} B_0, A^k = B_0 A_L^{k-1} B_1^T$, for $k \geq 1$.
2. If k is odd ($k > 1$), then

$$A_M^k = \begin{pmatrix} A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix};$$

If k is even ($k \geq 2$), then

$$A_M^k = \begin{pmatrix} A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k}{2}} \end{pmatrix}.$$

Proof By Proposition 1.5, the first assertion can be easily proved.

We prove the second assertion by induction on k .

When $k = 2$ or 3 ,

$$A_M^2 = \begin{pmatrix} A & AB_0 \\ B_1^T A & A_L^2 + A_L \end{pmatrix},$$

$$A_M^3 = A_M A_M^2 = \begin{pmatrix} 0_n & B_0 \\ B_1^T & A_L \end{pmatrix} \begin{pmatrix} A & AB_0 \\ B_1^T A & A_L^2 + A_L \end{pmatrix} = \begin{pmatrix} A^2 & (A^2 + A)B_0 \\ B_1^T(A^2 + A) & A_L^3 + A_L^2 \end{pmatrix}.$$

Hence, when $k = 2$ or 3 , the second assertion holds. We assume the second assertion holds for k . First, we suppose that k is odd. Then

$$A_M^k = \begin{pmatrix} A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix}.$$

By the first assertion, then

$$B_0(A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}}) = (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0.$$

Hence, we have

$$\begin{aligned} A_M^{k+1} &= A_M A_M^k = \\ &\begin{pmatrix} 0_n & B_0 \\ B_1^T & A_L \end{pmatrix} \begin{pmatrix} A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix} = \\ &\begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}) & A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix}. \end{aligned}$$

Similarly, we can prove that when k is even the second assertion holds. Hence our proof follows.

Similarly, we can prove the following two lemmas.

Lemma 2.3 Let D be a digraph. Then

1. If k is odd ($k > 1$), then

$$A_{M'}^k = \begin{pmatrix} A^k + A^{k-2} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^{k-1} + A_L^{k-2} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix};$$

2. If k is even ($k \geq 2$), then

$$A_{M'}^k = \begin{pmatrix} A^k + A^{k-2} + \dots + A^{\frac{k}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}}) & A_L^{k-1} + A_L^{k-2} + \dots + A_L^{\frac{k}{2}} \end{pmatrix}.$$

Lemma 2.4 Let D be a digraph. Then

1. If k is odd ($k > 1$), then

$$A_T^k = \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix};$$

2. If k is even ($k \geq 2$), then

$$A_T^k = \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k}{2}}) & A_L^k + A_L^k + \dots + A_L^{\frac{k}{2}} \end{pmatrix}.$$

Lemma 2.4 was proved by Yan and Zhang [16] and You and Liu et al. [17]. The following lemma is obvious.

Lemma 2.5 Suppose that there do not exist directed cycles in D . Then

1. $p(M(D)) = p(M'(D)) = p(T(D)) = 1$.

2. If we denote by $l(D)$ the length of the longest directed path of D , then $l(M(D)) = l(M'(D)) = l(T(D)) = 2l(D)$.

Partial results on $T(D)$ in Lemma 2.5 were proved by Yan and Zhang [16] and You and Liu et al. [17].

Lemma 2.6 Let D be a strongly connected digraph. Then all of $M(D)$, $M'(D)$ and $T(D)$ are primitive digraphs.

Proof By the definition of $M(D)$, it is obvious that $M(D)$ is strongly connected. Note that there exists at least a directed cycle in D . If we denote by l the length of this directed cycle. Then there at least exist two directed cycles in $M(D)$ whose lengths are l and $l + 1$, respectively. By Proposition 1.3, then $p(M(D)) = 1$. Hence $M(D)$ is a primitive digraph. Similarly, we can prove that both of $M'(D)$ and $T(D)$ are primitive digraphs. Thus our proof follows.

Corollary 2.7 Let D be a digraph. Then $p(M(D)) = p(M'(D)) = p(T(D)) = 1$.

Proof If there are not directed cycles in D , then, by Lemma 2.5, $p(M(D)) = p(M'(D)) = p(T(D)) = 1$. If there is at least a directed cycle in D , denoted by l its length, then there at least exist two directed cycles in $M(D)$ ($M'(D)$ and $T(D)$) whose lengths are l and $l + 1$, respectively. we distinguish the following three cases:

Case 1 Let D be a strongly connected digraph.

By Lemma 2.6, all of $M(D)$, $M'(D)$, and $T(D)$ are primitive digraphs, Hence $p(M(D)) = p(M'(D)) = p(T(D)) = 1$.

Case 2 Let D be a weakly connected digraph.

It is easy to see that $M(D)$ ($M'(D)$ and $T(D)$) is weakly connected. By Case 1 and Proposition 1.3, then $p(M(D)) = p(M'(D)) = p(T(D)) = 1$.

Case 3 Let $D_i(1 \leq i \leq c)$ be all of weakly connected components of D .

By Case 2 and Lemma 2.5, then $p(M(D_i)) = p(M'(D_i)) = p(T(D_i)) = 1$. Hence $p(M(D)) = p(M'(D)) = p(T(D)) = 1$.

Thus we have completed the proof of Corollary 2.7.

Partial results on $T(D)$ in Lemma 2.6 and Corollary 2.7 were proved by Yan and Zhang [16] and You and Liu et al. [17].

Lemma 2.8 Let D be a digraph with no sources, and let C_1 and C_2 be two $n \times n$ matrices whose entries are 0 or 1. If $B_1^T C_1 = B_1^T C_2$, then $C_1 = C_2$, where B_1 is the in-incidence matrix of D .

Proof Let $C_1 = (c_{ij}^1)$ and $C_2 = (c_{ij}^2)$. We need to prove that $c_{ij}^1 = c_{ij}^2$ for $1 \leq i, j \leq n$. By the definition of B_1 , for vertex v_i of D , then there is an arc x_m in D such that v_i is the head of arc x_m . Hence $b_{im}^1 = 1$. So the mj -entry $(B_1^T C_1)_{mj}$ of matrix $B_1^T C_1$ equals $\sum_{k=1}^n b_{km}^1 c_{kj}^1 = b_{im}^1 c_{ij}^1 = c_{ij}^1$, since there only exists one entry b_{im}^1 which is not zero in the m -th column of B_1 . Similarly, the mj -entry $(B_1^T C_2)_{mj}$ of matrix $B_1^T C_2$ equals c_{ij}^2 . Noting that $B_1^T C_1 = B_1^T C_2$, hence $c_{ij}^1 = c_{ij}^2$. This shows that $C_1 = C_2$. Our proof thus follows.

Lemma 2.9 Let D be a digraph, $k = k(D)$ and $p = p(D)$. Then $\sum_{i=s}^t A^i = A^k + A^{k+1} + \dots + A^{k+p-1}$, for $s \geq k$ and $t \geq s + p - 1$.

By the definitions of $k(D)$ and $p(D)$, this is clear.

3 Main results

Theorem 3.1 Let D be a digraph. Then

1. $k(M(D)) \leq \max\{2p(D), 2k(D) + 2\}$.
2. $k(M'(D)) \leq \max\{2p(D), 2k(D) + 2\}$.
3. $k(T(D)) \leq \max\{2p(D) - 1, 2k(D) + 1\}$.

Proof First, we prove that $k(M(D)) \leq \max\{2p(D), 2k(D) + 2\}$.

Let $k = k(D), p = p(D)$. We distinguish the following two cases:

Case 1 If $2p(D) \geq 2k(D) + 2$, then $p \geq k + 1$. By Corollary 2.7, we only need to prove that $A_M^{2p} = A_M^{2p+1}$. By Lemma 2.2,

$$A_M^{2p} = \begin{pmatrix} A^{2p-1} + A^{2p-2} + \dots + A^p & (A^{2p-1} + A^{2p-2} + \dots + A^p)B_0 \\ B_1^T(A^{2p-1} + A^{2p-2} + \dots + A^p) & A_L^{2p} + A_L^{2p-1} + \dots + A_L^p \end{pmatrix};$$

$$A_M^{2p+1} = \begin{pmatrix} A^{2p} + A^{2p-1} + \dots + A^{p+1} & (A^{2p} + A^{2p-1} + \dots + A^p)B_0 \\ B_1^T(A^{2p} + A^{2p-1} + \dots + A^p) & A_L^{2p+1} + A_L^{2p} + \dots + A_L^{p+1} \end{pmatrix}.$$

Noting that $p \geq k + 1$, by Lemma 2.9, then

$$A^{2p-1} + A^{2p-2} + \dots + A^p = A^k + A^{k+1} + \dots + A^{k+p-1}.$$

Similarly, we have

$$A^{2p} + A^{2p-1} + \dots + A^{p+1} = A^k + A^{k+1} + \dots + A^{k+p-1}.$$

Hence

$$A^{2p-1} + A^{2p-2} + \dots + A^p = A^{2p} + A^{2p-1} + \dots + A^{p+1}.$$

By Proposition 1.2, $k(L(D)) \leq k(D) + 1$. Similarly, we can prove that

$$(A^{2p-1} + A^{2p-2} + \dots + A^p)B_0 = (A^{2p} + A^{2p-1} + \dots + A^p)B_0;$$

$$B_1^T(A^{2p-1} + A^{2p-2} + \dots + A^p) = B_1^T(A^{2p} + A^{2p-1} + \dots + A^p);$$

$$A_L^{2p} + A_L^{2p-1} + \dots + A_L^p = A_L^{2p+1} + A_L^{2p} + \dots + A_L^{p+1}.$$

Hence $A_M^{2p} = A_M^{2p+1}$. This shows that $k(M(D)) \leq 2p = \max\{2p, 2k + 2\}$.

Case 2 If $2p(D) \leq 2k(D) + 2$, then $p \leq k + 1$. By Lemma 2.9, we only need to prove that $A_M^{2k+2} = A_M^{2k+3}$. By Lemma 2.2, we have

$$A_M^{2k+2} = \begin{pmatrix} A^{2k+1} + A^{2k} + \dots + A^{k+1} & (A^{2k+1} + A^{2k} + \dots + A^{k+1})B_0 \\ B_1^T(A^{2k+1} + A^{2k} + \dots + A^{k+1}) & A_L^{2k+2} + A_L^{2k+1} + \dots + A_L^{k+1} \end{pmatrix};$$

$$A_M^{2k+3} = \begin{pmatrix} A^{2k+2} + A^{2k+1} + \dots + A^{k+2} & (A^{2k+2} + A^{2k+1} + \dots + A^{k+1})B_0 \\ B_1^T(A^{2k+2} + A^{2k+1} + \dots + A^{k+1}) & A_L^{2k+3} + A_L^{2k+2} + \dots + A_L^{k+2} \end{pmatrix}.$$

Similarly to the proof of Case 1, we can see that

$$A^{2k+1} + A^{2k} + \dots + A^{k+1} = A^{2k+2} + A^{2k+1} + \dots + A^{k+2};$$

$$\begin{aligned} (A^{2k+1} + A^{2k} + \dots + A^{k+1})B_0 &= (A^{2k+2} + A^{2k+1} + \dots + A^{k+1})B_0; \\ B_1^T(A^{2k+1} + A^{2k} + \dots + A^{k+1}) &= B_1^T(A^{2k+2} + A^{2k+1} + \dots + A^{k+1}); \\ A_L^{2k+2} + A_L^{2k+1} + \dots + A_L^{k+1} &= A_L^{2k+3} + A_L^{2k+2} + \dots + A_L^{k+2}. \end{aligned}$$

Hence $A_M^{2k+2} = A_M^{2k+3}$. This shows that $k(M(D)) \leq 2k + 2 = \max\{2p, 2k + 2\}$.

Combining Case 1 and Case 2, we have

$$k(M(D)) \leq \max\{2p(D), 2k(D) + 2\}.$$

Similarly, we can prove assertions 2 and 3.

The third assertion in Theorem 3.1 was proved by Yan and Zhang [15].

Remark 1 Let $D_1 = (V(D_1), A(D_1))$, $V(D_1) = \{1, 2, 3\}$, and $A(D_1) = \{(1, 2), (2, 1), (2, 3), (3, 2)\}$. Then $p(D_1) = 2$, $k(D_1) = 1$, $k(M(D_1)) = k(M'(D_1)) = 4$, $k(T(D_1)) = 3$. This example shows that the upper bounds in Theorem 3.1 are obtained.

Theorem 3.2 Let D be a primitive digraph; then

$$k(M(D)) \leq k(D) + 1, k(M'(D)) \leq k(D) + 2, k(T(D)) \leq k(D) + 1.$$

Proof Let $k(D) = k$, J_n be the $n \times n$ matrix with all entries equal one. Then $A^k = J_n$. By Proposition 1.2, $A_L^{k+1} = J_m$. If k is odd, then

$$A_M^{k+1} = \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}) & A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix}.$$

Noting that $J_n B_0 = J_{nm}$, $B_1^T J_n = J_{mn}$, where J_{nm} and J_{mn} denote the $n \times m$ matrix with all entries equal one and the $m \times n$ matrix with all entries equal one, respectively. Hence $A_M^{k+1} = J_{n+m}$. Similarly, if k is even then $A_M^{k+1} = J_{n+m}$. This shows that $k(M(D)) \leq k(D) + 1$. Similarly, we can prove that $k(M'(D)) \leq k(D) + 2$, $k(T(D)) \leq k(D) + 1$. Our theorem is thus proved.

The third assertion in Theorem 3.2 was obtained by Yan and Zhang [16] and You and Liu et al. [17].

Remark 2 Let $D_2 = (V(D_2), A(D_2))$, $V(D_2) = \{1, 2\}$, and $A(D_2) = \{(1, 1), (1, 2), (2, 1)\}$. It is obvious that D_2 is a primitive digraph and $k(D_2) = 2$, $k(M(D_2)) = k(T(D_2)) = 3$, $k(M'(D_2)) = 4$. This example shows that the upper bounds in Theorem 3.2 are obtained.

Theorem 3.3 Let D be a digraph. Then

1. If D is a directed cycle with n vertices, then $k(M(D)) = k(M'(D)) = k(T(D)) + 1 = 2n$.
2. If there exist no directed cycles in D , then $k(M(D)) = k(M'(D)) = k(T(D)) = 2k(D) - 1 = 2l(D) + 1$, where $l(D)$ denotes the length of the longest directed path in D .

Proof It is easy to prove that the first assertion holds.

Assume that there exist no directed cycles in D . By Lemma 2.5, then $p(M(D)) = p(M'(D)) = p(T(D)) = 1$, $l(M(D)) = l(M'(D)) = l(T(D)) = 2l(D)$. By Proposition 1.2, then $k(M(D)) = k(M'(D)) = k(T(D)) = 2l(D) + 1$, and $k(D) = l(D) + 1$. Hence $k(M(D)) = k(M'(D)) = k(T(D)) = 2k(D) - 1 = 2l(D) + 1$. Thus the second assertion holds. Our proof thus follows.

Partial results on $T(D)$ in Theorem 3.3 were obtained by Yan and Zhang [16] and You and Liu et al. [17].

Theorem 3.4 Let D be a strongly connected digraph. Then $k(M(D)) \geq \text{Diam}(D) + 1, k(M'(D)) \geq \text{Diam}(D) + 2, k(T(D)) \geq \text{Diam}(D) + 1$.

Proof Let $\text{Diam}(D) = d$. If D is a directed cycle with n vertices, then by Theorem 3.3 we have $k(M(D)) = k(M'(D)) = k(T(D)) + 1 = 2n \geq d + 1$. If D is not a cycle, then by Proposition 1.3 we have $\text{Diam}(L(D)) = d + 1$. Let $x_i x_{i+1} \cdots x_{i+d+1}$ be the directed path of length $d + 1$ in $L(D)$. Then the $(i, i+d+1)$ -entry of matrix $I + A_L + A_L^2 + \cdots + A_L^d$ equals zero. Hence $I + A_L + A_L^2 + \cdots + A_L^d \neq J_m$. By Lemma 2.2, then $A_M^d \neq J_{n+m}$. By Lemma 2.6, $M(D)$ is a primitive digraph. Hence $k(M(D)) \geq d+1$. Similarly, we can prove that $k(M'(D)) \geq d+2$ and $k(T(D)) \geq d+1$. Our proof thus follows.

The third result in Theorem 3.4 was obtained by Yan and Zhang [16].

Remark 3 Let $D_3 = (V(D_3), A(D_3)), V(D_3) = \{1, 2\}, A(D_3) = \{(1, 1), (1, 2), (2, 1), (2, 2)\}$. It is easy to prove that $k(M(D_3)) = k(T(D_3)) = 2, k(M'(D_3)) = 3$, and $\text{Diam}(D_3) = 1$. Hence $k(M(D_3)) = k(T(D_3)) = \text{Diam}(D_3) + 1$ and $k(M'(D_3)) = \text{Diam}(D_3) + 2$. This example shows that the lower bounds in Theorem 3.4 are obtained.

Theorem 3.5 Let D be a digraph. Then $k(T(D)) \leq k(M(D)) \leq k(T(D)) + 1, k(T(D)) \leq k(M'(D)) \leq k(T(D)) + 1$.

Proof First, we prove that $k(T(D)) \leq k(M(D)) \leq k(T(D)) + 1$. Let $k(M(D)) = k$. By Corollary 2.7, then $A_M^k = A_M^{k+1}$. We distinguish the following two cases:

Case 1 Let k be odd. Since $A_M^k = A_M^{k+1}$, by Lemma 2.2, we have

$$\begin{aligned} A_M^k &= \begin{pmatrix} A^{k-1} + A^{k-2} + \cdots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \cdots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \cdots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \cdots + A_L^{\frac{k+1}{2}} \end{pmatrix} \\ &= \begin{pmatrix} A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}}) & A_L^{k+1} + A_L^k + \cdots + A_L^{\frac{k+1}{2}} \end{pmatrix} = A_M^{k+1}. \end{aligned}$$

Hence

- (1) $A^{k-1} + A^{k-2} + \cdots + A^{\frac{k+1}{2}} = A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}}$;
- (2) $(A^{k-1} + A^{k-2} + \cdots + A^{\frac{k-1}{2}})B_0 = (A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}})B_0$;
- (3) $B_1^T(A^{k-1} + A^{k-2} + \cdots + A^{\frac{k-1}{2}}) = B_1^T(A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}})$;
- (4) $A_L^k + A_L^{k-1} + \cdots + A_L^{\frac{k+1}{2}} = A_L^{k+1} + A_L^k + \cdots + A_L^{\frac{k+1}{2}}$.

By (1), we have

$$A(A^{k-1} + A^{k-2} + \cdots + A^{\frac{k+1}{2}}) + A^{\frac{k+1}{2}} = A(A^k + A^{k-1} + \cdots + A^{\frac{k+1}{2}}) + A^{\frac{k+1}{2}}.$$

Thus

$$(1') \quad A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} = A^{k+1} + A^k + \dots + A^{\frac{k+1}{2}}.$$

By Lemma 2.4, then

$$A_T^k = \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix};$$

$$A_T^{k+1} = \begin{pmatrix} A^{k+1} + A^k + \dots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}) & A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix}.$$

By (1'), (2), (3), and (4), then $A_T^k = A_T^{k+1}$. Thus $k(T(D)) \leq k = k(M(D))$.

Case 2 Let k be even. Similarly, we can prove that $k(T(D)) \leq k = k(M(D))$.

Hence we proved that $k(T(D)) \leq k(M(D))$.

Now we prove that $k(M(D)) \leq k(T(D)) + 1$. Let $k(T(D)) = k'$. By Corollary 2.7, then $A_T^{k'} = A_T^{k'+1}$. We distinguish the following two cases:

Case a Let k' be odd. Since $A_T^{k'} = A_T^{k'+1}$, by Lemma 2.4 we have

$$A_T^{k'} = \begin{pmatrix} A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}} & (A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}})B_0 \\ B_1^T(A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}}) & A_L^{k'} + A_L^{k'-1} + \dots + A_L^{\frac{k'+1}{2}} \end{pmatrix}$$

$$= \begin{pmatrix} A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}} & (A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0 \\ B_1^T(A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}}) & A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+1}{2}} \end{pmatrix} = A_T^{k'+1}.$$

Hence

$$(5) \quad A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}} = A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}};$$

$$(6) \quad (A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}})B_0 = (A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0;$$

$$(7) \quad B_1^T(A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}}) = B_1^T(A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}});$$

$$(8) \quad A_L^{k'} + A_L^{k'-1} + \dots + A_L^{\frac{k'+1}{2}} = A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+1}{2}}.$$

By (6), then

$$(A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}})B_0 B_1^T = (A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0 B_1^T.$$

Hence

$$(5') \quad A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}} = A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+3}{2}}.$$

By (5), then

$$(6') \quad (A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0 = (A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}})B_0;$$

$$(7') \quad B_1^T(A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}}) = B_1^T(A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}}).$$

By (7), then

$$B_1^T(A^{k'-1} + A^{k'-2} + \dots + A^{\frac{k'-1}{2}})B_0 = B_1^T(A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0.$$

Hence

$$(9) \quad A_L^{k'} + A_L^{k'-1} + \dots + A_L^{\frac{k'+1}{2}} = A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+3}{2}}.$$

By (8), then

$$(10) \quad A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+3}{2}} = A_L^{k'+2} + A_L^{k'+1} + \dots + A_L^{\frac{k'+3}{2}}.$$

We plus both sides of the equation (9) by $A_L^{k'+1}$, then we have

$$(11) \quad A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+1}{2}} = A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+3}{2}}.$$

By (10) and (11), thus

$$(8') \quad A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+1}{2}} = A_L^{k'+2} + A_L^{k'+1} + \dots + A_L^{\frac{k'+3}{2}}.$$

By Lemma 2.2, then

$$A_M^{k'+1} = \begin{pmatrix} A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}} & (A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}})B_0 \\ B_1^T(A^{k'} + A^{k'-1} + \dots + A^{\frac{k'+1}{2}}) & A_L^{k'+1} + A_L^{k'} + \dots + A_L^{\frac{k'+1}{2}} \end{pmatrix};$$

$$A_M^{k'+2} = \begin{pmatrix} A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+3}{2}} & (A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}})B_0 \\ B_1^T(A^{k'+1} + A^{k'} + \dots + A^{\frac{k'+1}{2}}) & A_L^{k'+2} + A_L^{k'+1} + \dots + A_L^{\frac{k'+3}{2}} \end{pmatrix}.$$

By (5'), (6'), (7') and (8'), then $A_M^{k'+1} = A_M^{k'+2}$. Hence $k(M(D)) \leq k'+1 = k(T(D)) + 1$.

Case b Let k' be even. Similarly, we can prove that $k(M(D)) \leq k(T(D)) + 1$.

Hence we have proved that $k(T(D)) \leq k(M(D)) \leq k'(T(D)) + 1$.

Similarly we can prove that $k(T(D)) \leq k(M'(D)) \leq k(T(D)) + 1$. Our proof thus follows.

Remark 4 Let $D_4 = (V(D_4), A(D_4)), V(D_4) = \{1, 2, 3, 4\}, A(D_4) = \{(1, 2), (2, 1), (2, 3), (3, 4), (4, 1)\}$. Then $k(M(D_4)) = k(M'(D_4)) = k(T(D_4)) = 6$. Note that the digraph D_1 in Remark 1 shows that there are digraphs D such that $k(M(D)) = k(T(D)) + 1$ or $k(M'(D)) = k(T(D)) + 1$. Thus the bounds in Theorem 3.5 are obtained.

Corollary 3.6 Let D be a digraph. Then $k(M(D)) - 1 \leq k(M'(D)) \leq k(M(D)) + 1$.

Proof By Theorem 3.5, it follows that $k(T(D)) \leq k(M(D)) \leq k(T(D)) + 1$, and $k(T(D)) \leq k(M'(D)) \leq k(T(D)) + 1$. Hence $k(M(D)) - 1 \leq k(M'(D)) \leq k(M(D)) + 1$.

Remark 5 Let $D_5 = (V(D_5), A(D_5)), V(D_5) = \{1, 2, 3\}$, and $A(D_5) = \{(1, 2), (2, 2), (2, 3)\}$. Then $k(M(D_5)) = 3, k(M'(D_5)) = 2$. Hence $k(M'(D_5)) = k(M(D_5)) - 1$. This shows that there are digraphs D such that $k(M'(D)) = k(M(D)) - 1$. The digraph D_4 in Remark 4 shows that there are digraphs D such that $k(M'(D)) = k(M(D))$. The digraph D_2 in Remark 2 shows that there are digraphs D such that $k(M'(D)) = k(M(D)) + 1$.

Theorem 3.7 Let D be a digraph. If there are not both sources and sinks in D , then $k(M(D)) \leq k(M'(D)) \leq k(M(D)) + 1$.

Proof By Corollary 3.6, we only need to prove that $k(M(D)) \leq k(M'(D))$. Let $k(M'(D)) = k$; then $A_{M'}^k = A_{M'}^{k+1}$. Hence we only need to prove that $A_M^k = A_M^{k+1}$. We distinguish the following two cases:

Case 1 We suppose that there are no sources in D .

Subcase 1.1 Let k be odd. Since $A_{M'}^k = A_{M'}^{k+1}$, by Lemma 2.3 we have

$$\begin{aligned} A_{M'}^k &= \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^{k-1} + A_L^{k-2} + \dots + A_L^{\frac{k-1}{2}} \end{pmatrix} \\ &= \begin{pmatrix} A^{k+1} + A^k + \dots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix} = A_{M'}^{k+1}. \end{aligned}$$

Hence

- (1) $A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} = A^{k+1} + A^k + \dots + A^{\frac{k+1}{2}}$;
- (2) $(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 = (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0$;
- (3) $B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) = B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})$;
- (4) $A_L^{k-1} + A_L^{k-2} + \dots + A_L^{\frac{k-1}{2}} = A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}}$.

We multiply both sides of the equality (4) by A_L ; then

$$A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+3}{2}} = A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+3}{2}}.$$

Hence

$$(1') \quad A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} = A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+1}{2}}.$$

By (4), then

$$(A_L^{k-1} + A_L^{k-2} + \dots + A_L^{\frac{k+1}{2}})B_1^T = (A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}})B_1^T.$$

by Lemma 2.1, then

$$(5) \quad B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}}) = B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}).$$

Let $C_1 = A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}}$, and $C_2 = A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}$. By (5), then $B_1^T C_1 = B_1^T C_2$. Since there are not sources in D , By Lemma 2.8, then $C_1 = C_2$.

Hence

$$(2') \quad A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}} = A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}.$$

By (2'), (2), (3), and (1'), then

$$\begin{aligned} A_M^k &= \begin{pmatrix} A^{k-1} + A^{k-2} + \dots + A^{\frac{k+1}{2}} & (A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}})B_0 \\ B_1^T(A^{k-1} + A^{k-2} + \dots + A^{\frac{k-1}{2}}) & A_L^k + A_L^{k-1} + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix} \\ &= \begin{pmatrix} A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}} & (A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}})B_0 \\ B_1^T(A^k + A^{k-1} + \dots + A^{\frac{k+1}{2}}) & A_L^{k+1} + A_L^k + \dots + A_L^{\frac{k+1}{2}} \end{pmatrix} = A_M^{k+1}. \end{aligned}$$

Hence $k(M(D)) \leq k = k(M'(D))$.

Subcase 1.2 Let k be even. Similarly, we can prove $k(M(D)) \leq k = k(M'(D))$.

Case 2 We suppose that there are no sinks in D .

Let D^* denote the converse digraph of D . Then there are no sources in D^* . By Case 1, we have $k(M(D^*)) \leq k(M'(D^*))$. It is easy to prove that $M(D^*) = M^*(D)$,

$M'(D^*) = M'^*(D)$, $k(M(D)) = k(M^*(D))$ and $k(M'(D)) = k(M'^*(D))$. Hence $k(M(D)) \leq k(M'(D))$.

By Cases 1 and 2, Theorem 3.7 holds.

By Theorem 3.7, the following corollary is obvious.

Corollary 3.8 Let D be a strongly connected digraph; then

$$k(M(D)) \leq k(M'(D)) \leq k(M(D)) + 1.$$

4 Some problems

In this section, we will pose some problems on the classification of digraphs by their indices of convergence.

Let D be a digraph. Theorem 3.5 shows that $k(T(D)) \leq k(M(D)) \leq k(T(D)) + 1$ and $k(T(D)) \leq k(M'(D)) \leq k(T(D)) + 1$. Corollary 3.6 shows that $k(M(D)) - 1 \leq k(M'(D)) \leq k(M(D)) + 1$ and Theorem 3.7 shows that if there are not both sources and sinks in D then $k(M(D)) \leq k(M'(D)) \leq k(M(D)) + 1$. Hence the following problems of characterization of digraphs are worth considering.

Problem 1 Determine the digraphs D such that $k(M(D)) = k(T(D))$ or $k(M(D)) = k(T(D)) + 1$, respectively.

Problem 2 Determine the digraphs D such that $k(M'(D)) = k(T(D))$ or $k(M'(D)) = k(T(D)) + 1$, respectively.

Problem 3 Determine the digraphs D such that $k(M'(D)) = k(M(D)) - 1$, or $k(M'(D)) = k(M(D))$, or $k(M'(D)) = k(M(D)) + 1$, respectively.

Problem 4 Suppose that there are not both sources and sinks in D . Determine the digraphs D such that $k(M'(D)) = k(M(D))$ or $k(M'(D)) = k(M(D)) + 1$, respectively.

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