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#### Abstract

Let $G$ be a simple graph on $n$ vertices having edge-connectivity $\kappa^{\prime}(G)>0$ and minimum degree $\delta(G)$. We say $G$ is $k$-critical if $\kappa^{\prime}(G)=k$ and $\kappa^{\prime}(G-e)<k$ for every edge $e$ of $G$. In this paper we prove that a k-critical graph has $\kappa^{\prime}(\mathrm{G})=\delta(\mathrm{G})$. We describe a number of classes of k-critical graphs and consider the problem of determining the edge-maximal ones.


## 1. INTRODUCTION

For our purposes graphs are undirected, finite, loopless and have no multiple edges. For the most part our notation and terminology follows that of Bondy and Murty [2]. Thus $G$ is a graph with vertex set $V(G)$, edge set $E(G), \nu(G)$ Vertices and $\varepsilon(G)$ edges. However, we denote the complement of $G$ by $\bar{G}, K_{n}$ denotes the complete graph on $n$ vertices, $K_{n, m}$ the complete bipartite graph with bipartioning sets of order $n$ and $m$, and $C_{\ell}$ a cycle of length $\ell$. The join of disjoint graphs $G$ and $H$, denoted $G \vee H$, is the graph obtained by joining each vertex of $G$ to each vertex of $H$.

A good deal of graph theory is concerned with the characterization of graphs having certain specified properties. Graph parameters of particular practical interest include: minimum degree $\delta(G)$, connectivity $\kappa(G)$, edge-connectivity $\kappa^{\prime}(G)$, diameter $d(G)$, chromatic number $\chi(G)$, and various covering numbers (vertex, edge, clique, etc.). In studying such parameters it is often useful to restrict attention to the so called "critical graphs".

Let $P$ be a graph parameter. A graph $G$ is said to be P-edge (vertex)-critical if $P(G-e) \neq P(G)(P(G-v) \neq P(G))$ for every edge $e$ (vertex $v$ ) of $G$. For a given $P$, the problem that arises is that of characterizing the class of P-edge-critical and class of P-vertex-critical graphs. In particular those that are edge-minimal or edge-maximal. This problem has been investigated for the edge case when $P$ is: connectivity (Halin [9]); diameter (Caccetta and Haggkvist [3], Fan [8]); chromatic index (Yap [16]); and the vertex covering number (Lovasz and Plummer [14]), and for the vertex case when $P$ is: connectivity (Chartrand [5], Entringer [7], Hamidoune [11], Krol and Veldman [13]); edge-connectivity (Cozzens and Wu [6]). The anologous problem for "edge addition" has been considered for diameter (Caccetta and Smyth [4], Ore [15]).

The object of this paper is to study graphs that are edge-critical with respect to the parameter $\kappa^{\prime}$. For simplicity we say a graph $G$ is $k$-critical if $\kappa^{\prime}(G)=k$ and $\kappa^{\prime}(G-e)<k$ for every edge e of $G$. Observe that: $K_{n}$ is ( $n-1$ )-critical; $K_{n, m}$ is t-critical, where $t=\min \{m, n\}$; every tree is 1 -critical; $C_{n}$ is 2-critical; and $K_{1} \vee C_{n}$ is 3 -critical. We prove that a k-critical graph $G$ has $k=\delta(G)$. This is anologous to the corresponding result of Halin [10] for edge-critical graphs with respect to $\kappa$. In addition, we shall consider the problem of determining the maximum number of edges in a k-critical graph.

## 2. RESULTS

Let $\mathscr{C}(\mathrm{n}, \mathrm{k})$ denote the class of $k$-critical graphs on n vertices. We begin our discussion with some constructions.

It is very well known that for any graph $G \kappa^{\prime}(G) \leq \delta(G)$. Further, given any positive integers $a$ and $b$ with $a \leq b$ there exists $a$ graph $G$ on $n \geq b+1$ vertices such that $\kappa^{\prime}(G)=a$ and $\delta(G)=b$. $A$ class of graphs corresponding the case $a=b$ is sometimes referred to as the Harary graphs and are described in standard texts (p.48, [2]). Let $H(n, r)$ denote the class of graphs on $n$ vertices having minimum
degree and edge-connectivity $r$ and having $\left\lceil\frac{1}{2} n r\right\rceil$ edges. Observe that for $k \geq 2, H(n, k) \subseteq \mathscr{C}(n, k)$. In fact, this class is edge-minimal. The following is an immediate consequence of the definition of criticality.

Lemma 1. Let $G$ be a graph with $\kappa^{\prime}(G)=\delta(G)=k$ and every edge of $G$ is incident to at least one vertex of minimum degree. Then $G$ is k-critical.

Thus we have one class of critical graphs. Let $A(n, k)$ denote the subclass of $\mathscr{C}(n, k)$ consisting of those graphs in which every edge is incident to a vertex of minimum degree. Clearly $K_{k, n-k} \in A(n, k)$ for $n \geq 2 k$. Later we shall show that for $n \geq 3 k, K_{k, n-k}$ is an edge maximal graph of $\mathscr{A}(\mathrm{n}, \mathrm{k})$. We now construct a class of graphs in $A(n, k)$.

Let $H \in H(n-x, k-x)$ for $1 \leq x \leq n-k$, and define $G=H \vee \bar{K}_{x}$. If $\mathrm{n}-\mathrm{x}$ and $\mathrm{k}-\mathrm{x}$ are both odd and $\mathrm{x} \neq \mathrm{n}-\mathrm{k}$, then G contains an edge $e=u v$ with $u \in H$ and $v \in \bar{K}_{x}$ such that $G-e$ is $k$-edge connected; in fact $G-e$ is $k$-critical. Thus if we let $G^{\prime}=G-e$ if both $n-x$ and $k-x$ are odd and $G^{\prime}=G$ otherwise, then $G^{\prime} \in \mathscr{A}(n, k)$ and has ( $\left.n-x\right) x$ $+\left\lfloor\frac{1}{2}(n-x)(k-x)\right\rfloor$ edges. Figure 1 below illustrates this construction. Note that in our illustration the " $=$ " means all edges

H
$\overline{\mathrm{K}}_{2}$


G


G'

Figure 1
between the vertices of $H$ and the vertices of $\bar{K}_{2}$. We use this notation in all our diagrams. In Theorem 2 we show that $G^{\prime}$ is edge-maximal for $n<3 k$.

The graphs drawn in Figure 2 below show that $\mathcal{A}(\mathrm{n}, \mathrm{k}) \neq \mathscr{E}(\mathrm{n}, \mathrm{k})$. In fact, it is easy to construct graphs in the class $\mathscr{C}(n, k) \backslash \mathcal{A}(n, k)$. One construction is the following. Let $n=2 k t+r, 0 \leq r \leq 2 k-1$. The graph $G$ obtained by adding edges to the graph ( $2 t-1$ ) $\bar{K}_{k} \cup \bar{K}_{k+r}$ as shown in Figure 3 is in the class $\mathscr{C}(n, k) \backslash \mathcal{A}(n, k)$ for $t \geq 2$. Note that a line joining two graphs means a "perfect matching" between the two graphs. We adopt this convention in all our diagrams.

$G_{1} \in \mathscr{C}(9,2) \backslash A(9,2)$

$G_{2} \in \mathscr{C}(8,3) \backslash A(8,3)$

Figure 2.


Figure 3.

Let $G$ be a graph with a cut vertex $v$. That is $G=G_{1} \cup G_{2}$ with $V\left(G_{1}\right) \cap V\left(G_{2}\right)=\{v\}$. Then $\kappa^{\prime}(G)=\min \left\{\kappa^{\prime}\left(G_{1}\right), \kappa^{\prime}\left(G_{2}\right)\right\}$. We thus have the following simple but useful property.

Lemma 2. Let $G_{i} \in \mathscr{C}\left(n_{i}, k\right), 1 \leq i \leq t$. Then the graph $G$ whose blocks are $G_{1}, G_{2}, \ldots, G_{t}$ is in the class $\mathscr{C}\left(n_{1}+n_{2}+\ldots+n_{t}-t+1, k\right)$.

This lemma provides a procedure for building larger critical graphs from smaller ones.

Let $\rho(u, v)$ denote the maximum number of edge-disjoint paths between vertices $u$ and $v$ in $G$. Menger's theorem states that

$$
\begin{aligned}
\kappa^{\prime}(G)= & \min \quad\{\rho(u, v)\} . \\
& u, v \in V(G)
\end{aligned}
$$

We make use of this result in our next lemma.

Lemma 3. Let $G$ be a $k$-edge-connected graph. Then $G \in \mathscr{C}(n, k)$ if and only if $\rho(u, v)=k$ for every pair of adjacent vertices $u, v$ in $G$.

Proof: Suppose $G \in G(n, k)$ and let $e=u v$ be an edge of $G$. Consider the graph $G^{\prime}=G-e$. We have $\kappa^{\prime}\left(G^{\prime}\right)=k-1$. Let $E^{\prime}$ be an edge-cut set of $G^{\prime}$ having $k-1$ elements. Since $k(G)=k$, the graph $G^{\prime \prime}=G^{\prime}-$ $E^{\prime}$ consists of exactly two components. Further, $G^{\prime \prime}+e$ is connected. Hence the vertices $u$ and $v$ are in different components of $G^{\prime \prime}$. Consequently the set $E=E^{\prime} \cup\{e\}$ is an edge-cut set of $G$ having $k=$ $\kappa(G)$ elements. Thus $\rho(u, v) \leq k$. Menger's theorem now implies that $\rho(u, v)=k$ as required.

Conversely, if $\rho(u, v)=k$ for every pair of adjacent vertices $u$, $v$ in $G$, then $\kappa^{\prime}(G-u v) \leq k-1$ and hence $\kappa^{\prime}(G) \leq k$. Now since $\kappa(G) \geq$ $k$ we have $G \in \mathscr{C}(n, k)$. This completes the proof of the lemma.

Thus we can test whether or not a graph $G$ is k-critical using standard network flow algorithms.

We now prove the main result of this paper.

Theorem 1. If $G$ is a $k$-critical graph, then $\delta(G)=k$.
Proof: Let $G \in \mathscr{C}(n, k)$. Then $n \geq k+1$. If $n=k+1$, then $G=K_{k+1}$ and hence $\delta(G)=k$. So we suppose that $n \geq k+2$. We prove the theorem by contradiction. Assume $\delta(\mathrm{G})>\mathrm{k}$.

Let $\mathcal{E}(\mathrm{k})$ denote the set of edge-cut sets of $G$ having $k$ elements. If $E^{\prime} \in \mathscr{E}(k)$, then $G-E^{\prime}$ consists of two components. Let $E^{*}$ denote an element of $\mathcal{E}(\mathrm{k})$ such that $G-E^{*}$ has the smallest possible component. Let $G_{1}$ and $G_{2}$ denote the components of $G-E^{*}$ and suppose, without loss of generality, that $n_{1}=\left|V\left(G_{1}\right)\right| \leq n_{2}=\left|V\left(G_{2}\right)\right|$.

Let $A_{i}$ denote the set of vertices of $G_{i}, i=1,2$, that are incident to an edge of $E^{*}$. We prove the theorem by showing that $n_{1}=$ 1. Suppose that $n_{1} \geq 2$. We will show that $n_{1} \geq k+2$. This is certainly the case if $G_{1}-A_{1} \neq \phi$ as we have assumed that $\delta \geq k+1$. So suppose that every vertex of $G_{1}$ is in $A_{1}$. We have

$$
\begin{aligned}
\sum_{u \in V\left(G_{1}\right)} d_{G_{1}}(u) & =\sum_{u \in V\left(G_{1}\right)} d_{G}(u)-k \\
& \geq n_{1}(k+1)-k \\
& =k\left(n_{1}-1\right)+n_{1},\left(n_{1} \leq k\right) \\
& \geq n_{1}\left(n_{1}-1\right)+n_{1}>n_{1}\left(n_{1}-1\right)
\end{aligned}
$$

a contradiction. Thus $n_{1} \geq k+2$. Since $\left|A_{1}\right| \leq k$, we must have at least two vertices of $G_{1}$ not in $A_{1}$. Hence there exists an edge $e=x y$ in $G_{1}$ with $x, y \notin A_{1}$. Since $G$ is $k$-critical, $\kappa^{\prime}(G-e)=k-1$. Now since $n_{2} \geq n_{1}, G_{2}$ contains vertices which are not in $A_{2}$. Let $z$ be one such vertex. Clearly $z$ is joined to the vertices of $A_{1}$ by $k$-edge disjoint paths.

Since $\kappa^{\prime}(G)=k$ the vertices $x$ and $y$ must each be joined to the vertices of $A_{1}$ by at least $k$-edge disjoint paths. In fact, the choice of $E^{*}$ ensures that there are at least $k+1$ such paths.
This contradicts Lemma 3. Hence $n_{1}=1$. This completes the proof of our theorem.

We mentioned earlier in this section that the Harary graphs were edge-minimal members of $\mathscr{C}(n, k)$. The problem of determining the edge-maximal members of $\mathscr{C}(n, k)$ seems to be difficult. Our next result determines the maximum number of edges for a graph $G \in \mathcal{A}(\mathrm{n}, \mathrm{k})$.

Theorem 2. Let $G$ be an edge-maximal graph of $A(n, k)$. Then

$$
\varepsilon(G)= \begin{cases}k(n-k), & \text { if } n \geq 3 k \\ \left\lfloor\frac{1}{8}(n+k)^{2}\right\rfloor, & \text { otherwise }\end{cases}
$$

Proof: By Theorem $1 \delta(G)=k$. We denote the set of vertices of $G$ having degree $k$ by $X$ and the remaining vertices by $\bar{X}$. Let $n_{1}=|X|$. Since $G \in \mathbb{A}(\mathrm{n}, \mathrm{k})$, we must have $\mathrm{n}_{1} \geq \mathrm{k}+1$. Simple counting gives:

$$
\varepsilon(G) \leq \begin{cases}n_{1} k & \text { if } n_{1} \leq n-k  \tag{1}\\ n_{1}\left(n-n_{1}\right)+\left\lfloor\frac{1}{2} n_{1}\left(k-n+n_{1}\right)\right], & \text { otherwise }\end{cases}
$$

Let $g\left(n_{1}\right)$ denote the right hand side of (1).

## Clearly

$$
\begin{aligned}
\max & \left\{g\left(n_{1}\right)\right\} & =g(n-k) \\
n_{1} \leq n-k & & =k(n-k) .
\end{aligned}
$$

This maximum is attained by the graph $K_{k, n-k}$. For $n_{1} \geq n-k$, we have for fixed $n$ and $k$

$$
\begin{equation*}
g\left(n_{1}+1\right)-g\left(n_{1}\right)=\left\lfloor\frac{1}{2}(n+k-1)\right\rfloor-n_{1}+\delta\left(n_{1}\right) \cdot \delta(n-k-1) \tag{2}
\end{equation*}
$$

where $\delta(x)=0$ or 1 according to whether $x$ is even or odd. There is some algebra involved in establishing (2), but it is fairly elementary.

From (2) we deduce that $g\left(n_{1}\right)$ monotonically increases in $n_{1}$ for $n_{1} \leq\left[\frac{1}{2}(n+k-1)\right]$ and monotonically decreases in $n_{1} \geq\left[\frac{1}{2}(n+k+\right.$ 1) . Now since $n_{1} \geq n-k, g\left(n_{1}\right)$ is decreasing in $n_{1}$ for $n \geq 3 k$. Hence

$$
\begin{array}{lll}
\max & \left\{g\left(n_{1}\right)\right\} & \leq g(n-k) \\
n \geq 3 k & & =k(n-k) .
\end{array}
$$

For $n<3 k, g\left(n_{1}\right)$ attains its maximum value at $n_{1}=\left[\frac{1}{2}(n+k+1)\right]$. It is a straight forward algebraic exercise to verify that $g\left(\left\lvert\, \frac{1}{2}(n+k\right.\right.$ $+1)\rfloor$ ) $=\left\lfloor\frac{1}{8}(n+k)^{2}\right\rfloor$. An example of a graph in $\mathcal{A}(n, k)$ having this number of edges is the graph G' (described following Lemma 1) with $x=n-\left\lfloor\frac{1}{2}(n+k+1)\right\rfloor$.

Now

$$
(n+k)^{2}-8 k(n-k)=(n-3 k)^{2} \geq 0
$$

Hence

$$
\left\lfloor\frac{1}{8}(n+k)^{2}\right\rfloor \geq k(n-k)
$$

always. This completes the proof of the theorem.

It would be interesting to determine whether or not the edge-maximal graphs of $\mathscr{C}(n, k)$ coincide with the edge-maximal graphs of $\mathscr{A}(\mathrm{n}, \mathrm{k})$. Krol and Veldman [13] have conjectured that for $\kappa$-vertex-critical graphs the analogous question is true for $\kappa \geq 3$.

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