

# On clique polynomials

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## Abstract

Let  $G$  be a simple graph. We assign a polynomial  $C(G; x)$  to  $G$ , called the clique polynomial, where the coefficient of  $x^i$ ,  $i > 0$ , is the number of cliques of  $G$  with  $i$  vertices, and the constant term is 1. Fisher and Solow (1990), proved that this polynomial always has a real root. We prove this result by a simple and elementary method, which also implies the following results. If  $\zeta_G$  is the greatest real root of  $C(G; x)$  then for an induced subgraph  $H$  of  $G$ ,  $\zeta_H \leq \zeta_G$ , and for a spanning subgraph  $H$  of  $G$ ,  $\zeta_H \geq \zeta_G$ . As a consequence of the first inequality we have  $\alpha(G) \leq -1/\zeta_G$ , where  $\alpha(G)$  denotes the independence number of  $G$ .

## 1 Introduction

Throughout this paper we consider simple graphs, i.e. finite undirected graphs with no loops and multiple edges, and we use the terminology and notation of [1].

The *dependence polynomial* was first introduced by Fisher [2], who studied the following problem: How many  $n$  letter words can be made from an  $m$  letter alphabet if certain pairs of letters commute? Fisher and Solow [3] defined the dependence polynomial as follows:

$$f_G(x) = 1 - c_1x + c_2x^2 - c_3x^3 + \cdots + (-1)^\omega c_\omega x^\omega;$$

where  $\omega$  is the size of the largest clique in  $G$  and  $c_i$  denotes the number of complete subgraphs of size  $i$  in  $G$ . For a set  $S$  of words with an operation on them we assign a graph  $G_S$  such that  $V(G_S) = S$  and two vertices are joined iff they commute. Fisher [2] proved that the generating function for the above problem is precisely  $\frac{1}{f_{G_S}(x)}$ .

If we change the signs of all negative coefficients in  $f_G(x)$  to positive signs, we obtain a polynomial which is called the *clique polynomial* of  $G$ . Using the notation

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of [4] we denote it by  $C(G; x)$ . So we have:

$$C(G; x) = 1 + c_1x + c_2x^2 + c_3x^3 + \cdots + c_\omega x^\omega.$$

In [3] Fisher and Solow showed that the dependence polynomial of a graph always has a real root. In fact they prove that the smallest root (in absolute value) of  $f_G(x)$  is real. This result immediately implies the existence of a real root for the clique polynomial.

In this paper we give a simple proof of the later result. In addition, we show that there is some relation between the largest negative root of  $C(G; x)$  and that of  $C(H; x)$ , for special subgraphs  $H$  of  $G$ .

## 2 Results

We first present the following observation and then use it as the main tool to prove our main theorems.

**Lemma 1.** *Let  $G$  be a graph and let  $v \in V(G)$ . Then*

- (a)  $C(G; x) = C(G \setminus v; x) + xC(G[N(v)]; x)$ ; where  $N(v)$  is the neighborhood of  $v$ .
- (b)  $C(G; x) = C(G - uv; x) + x^2C(G[N(u) \cap N(v)]; x)$ ; where  $uv \in E(G)$ .

**Proof.** Let  $A_i$  be an  $i$ -clique of  $G$ . (a) Either  $v \notin A_i$ , then  $A_i$  is an  $i$ -clique in  $G \setminus v$ ; or  $v \in A_i$ , then  $A_i$  is obtained from an  $(i-1)$ -clique of  $G[N(v)]$ . Summing up the number of these two kinds of  $i$ -cliques we obtain relation (a).

(b) Either  $A_i$  does not contain the edge  $uv$ , then  $A_i$  is an  $i$ -clique in  $G - uv$ ; or it does contain  $uv$ , then  $A_i$  is obtained from an  $(i-2)$ -clique of  $G[N(u) \cap N(v)]$ . Summing up the number of these two kinds of  $i$ -cliques we obtain relation (b).  $\square$

To pursue our study we need the following notation:

**Notation.** *Let  $G$  be a graph and let  $\mathcal{Z}(G)$  be the set of negative real roots of  $C(G; x)$ . If  $\mathcal{Z}(G)$  is non-empty then define  $\zeta_G$  to be  $\max \mathcal{Z}(G)$  and otherwise to be  $-\infty$ .*

The following theorem plays an essential role where we reprove the result of Fisher and Solow. Also it presents a nice property of  $\zeta_G$  in conjunction with induced subgraphs.

**Theorem 1.** *If  $G$  is a graph and  $H$  is one of its induced subgraph, then  $\zeta_H \leq \zeta_G$ .*

**Proof.** Let  $n = |V(G)|$ . We prove the theorem by induction on  $n$ . For  $n = 1$  and 2 the assertion is obvious. If  $H$  is an arbitrary proper induced subgraph of  $G$ , then we can find a vertex  $v$  of  $G$  such that  $H$  is an induced subgraph of  $G \setminus v$ . Hence it is sufficient to prove the theorem for induced subgraphs of the form  $G \setminus v$ , for some  $v \in V(G)$ . Now, let  $v \in V(G)$ . If  $\mathcal{Z}(G \setminus v) = \emptyset$  then there is nothing to prove. So

we can assume that  $\mathcal{Z}(G \setminus v)$  is not empty. On other hand, by part (a) of Lemma 1 we have:

$$C(G; x) = C(G \setminus v; x) + xC(G[N(v)]; x).$$

Substituting  $\zeta_{G \setminus v}$  in both sides of the above equation and applying induction we have  $C(G[N(v)], \zeta_{G \setminus v}) \geq 0$ , thus  $C(G, \zeta_{G \setminus v}) \leq 0$ . On the other hand  $C(G, 0) = 1$ . So the theorem is proved.  $\square$

**Theorem 2.** *For every graph  $G$ ,  $-1 \leq \zeta_G < 0$ .*

**Proof.** Let  $u$  be a vertex of  $G$ , and  $H$  be the subgraph induced on  $u$ . Clearly  $\zeta_H = -1$ . Thus by the above theorem we must have  $\zeta_G \geq -1$ , as desired.  $\square$

Turan's theorem for triangle-free graphs is a consequence of Theorem 2 :

**Corollary 1.** *If  $G$  is a triangle-free graph then  $|E(G)| \leq |V(G)|^2/4$ .*

**Proof.** Since  $G$  has no triangle we have:

$$C(G; x) = 1 + |V(G)|x + |E(G)|x^2.$$

By Theorem 2,  $C(G; x)$  has a real root which implies that the discriminant of this polynomial i.e.  $|V(G)|^2 - 4|E(G)|$  is non-negative; as claimed.  $\square$

The two following propositions are obtained by considering some special induced subgraphs.

**Proposition 1.** *Let  $G$  be a graph and  $\alpha(G)$  be the independence number of  $G$ . Then*

$$\alpha(G) \leq -1/\zeta_G.$$

**Proof.** Consider the subgraph  $H$  induced by an independent set of size  $\alpha(G)$ . We have  $\zeta_H = -1/(\alpha(G))$  and by Theorem 1,  $\zeta_H \leq \zeta_G$ . This proves the proposition.  $\square$

**Proposition 2.** *Let  $G$  be a graph which is not complete and let  $g(G)$  be the girth of  $G$ . Then*

$$g(G) \leq \frac{-1}{\zeta_G^2 + \zeta_G}$$

**Proof.** Consider a cycle of  $G$  with the size  $g(G)$ . This is an induced subgraph of  $G$ . Calculating the  $\zeta$  of this cycle and applying Theorem 1 we obtain the desired inequality.  $\square$

**Remark 1.** By the same method one can prove a similar assertion with  $g(G)$  replaced by the length of the the smallest odd cycle.

The following corollary is an immediate consequence of Proposition 1:

**Corollary 2.** *For every graph  $G$ ,  $\chi(G) \geq -|V(G)|\zeta_G$ .*

**Theorem 3.** If  $G$  is a graph and  $H$  is a spanning subgraph of  $G$ , then  $\zeta_G \leq \zeta_H$ .

**Proof.** It is enough to prove the theorem in the case of  $H = G - e$  where  $e$  is an edge of  $G$ . Suppose  $e = uv$  for  $u, v \in V(G)$ . By the part (b) of Lemma 1 we have:

$$C(G; x) = C(G - uv; x) + x^2 C(G[N(u) \cap N(v)]; x),$$

where  $uv \in E(G)$ . Substitute  $\zeta_G$  in both sides of the above equation. We obtain:

$$C(G - uv; \zeta_G) = -\zeta_G^2 C(G[N(u) \cap N(v)]; \zeta_G). \quad (1)$$

On the other hand  $G[N(u) \cap N(v)]$  is an induced subgraph of  $G$  and therefore by Theorem 1 the right hand side of equation (1) is negative, which implies that  $C(G - uv; \zeta_G)$  is negative also. This together with the fact that  $C(G - uv; 0) = 1$ , implies the assertion.  $\square$

We can apply Theorem 3 to prove some necessary conditions for existence of Hamiltonian cycles and perfect matchings which are useful in some special cases.

**Corollary 3.** Let  $G$  be a graph with  $n$  vertices. We have:

(a) If  $n \geq 4$  and  $\zeta_G > \frac{-1+\sqrt{1-4/n}}{2}$ , then  $G$  is not Hamiltonian.

(b) If  $n \geq 2$  and  $\zeta_G > -1 + \sqrt{1 - 2/n}$ , then  $G$  does not have perfect matching.

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