

by

Bolian Liu

Department of Mathematics

South China Normal University

Guangzhou

People's Republic of China

Dingjun Lou

Department of Computer Science

Zhongshan University

Guangzhou 510275

People's Republic of China

and

Kewen Zhao

Department of Mathematics

South China Normal University

Guangzhou

People's Republic of China

Abstract

Let G be a 2-connected graph. It is proved that if for each pair of nonadjacent vertices u and v of G , $|N(u) \cup N(v)| \geq \nu - \delta + 1$, then G is vertex-pancyclic, which implies a conjecture of R. J. Faudree, R. J. Gould, M. S. Jacobson and L. Lesniak.

All graphs considered are finite, undirected and simple.

A graph G is said to be vertex-pancyclic if for each vertex v in G , v is contained in a cycle of length m in G for each m such that $3 \leq m \leq \nu(G)$.

We define $NCC(G) = \min \{ |NC(x) \cup NC(y)| \mid x, y \in V(G), x \neq y \text{ and } xy \notin E(G) \}$. And we write NC for $NCC(G)$ when no confusion arises. Let C be a cycle of G . Let u be a vertex in $V(C)$. We give C an orientation. Then u^+ denotes the successor of u on C in the orientation and u^- denotes the predecessor of u on C in the orientation. Let $S \subseteq V(C)$. Then $S^+ = \{ x^+ \mid x \in S \}$ and $S^- = \{ x^- \mid x \in S \}$. Let v be a vertex in $V(G) \setminus V(C)$. $N_C(v)$ denotes $NC(v) \cap V(C)$. Suppose $N_C(v) \neq \emptyset$. An A-structure on $N_C(v)$ is a pair of vertices x and y such that $x, y \in N_C(v)$ and $x^+ = y$. A suc-J-structure on S is an edge x^+y^+ such that $x, y \in S$, $x^+ \neq y$ and $y^+ \neq x$. A pre-J-structure on S is an edge x^-y^- such that $x, y \in S$, $x^- \neq y$ and $y^- \neq x$. Because of the obvious similarity between suc-J-structures and pre-J-structures, for ease of notation and presentation, we frequently give proofs only using suc-J-structures (or pre-J-structures). We denote by $C^+[u, v]$ the path on C from u to v in the orientation and by $C^-[u, v]$ the path on C from u to v in the reverse orientation. The end vertices u and v are included.

Let G and H be two graphs such that $E(G) \cap E(H) = \emptyset$.

We use $G+H$ to denote the graph with vertex set $V(G) \cup V(H)$ and edge set $E(G) \cup E(H)$.

For terminology and notation not defined in this paper, the reader is referred to [2].

In [3], Faudree, Gould, Jacobson and Lesniak conjectured that if G has order ν , connectivity t and satisfies $NC \geq \nu - t$ with $\delta \geq t + 1$, then G is vertex-pancyclic. Song [5] reformulated their conjecture on the Chinese Symposium on Cycle Problems in Graph Theory in the form that if G is 2-connected and $NC \geq \nu - \delta + 1$, then G is vertex-pancyclic. Obviously, Song's conjecture implies the conjecture by Faudree et al. In this paper, we prove Song's conjecture. The main idea of the proof is similar to that in [1].

2. The main result

Lemma 1 Let G be a 2-connected graph. If $NC \geq \nu - \delta + 1$, then every vertex of G lies on a triangle. Furthermore, if $\nu \geq 4$ then every vertex of G lies on a 4-cycle.

Proof Let u be a vertex of G . If there are two distinct vertices x, y in $NC(u)$ such that $xy \in E(G)$, then $uxyu$ is a triangle containing u . Otherwise $|NC(x) \cup NC(y)| \leq \nu - |NC(u)| \leq \nu - \delta$ for any two vertices $x, y \in NC(u)$, contradicting $NC \geq \nu - \delta + 1$.

Now assume $\nu \geq 4$. First, we claim $\delta \geq 3$. Suppose $d(x)$

≤ 2 for a vertex x in G . As $\nu(G) \geq 4$, there is a vertex y such that $xy \notin E(G)$. But $x, y \notin N(x) \cup N(y)$, which contradicts the hypothesis that $|N(x) \cup N(y)| \geq \nu - \delta + 1 \geq \nu - 2 + 1 = \nu - 1$.

Assume that $T = uvw$ is a triangle containing u .

Suppose there is no 4-cycle containing u . Since $\delta \geq 3$, there is a vertex t in $N(u) \setminus \{v, w\}$. Then $tv, tw \notin E(G)$, otherwise there is a 4-cycle containing u . Now $(N(t) \setminus \{u\}) \cap (N(w) \setminus \{u\}) = \emptyset$ and $(N(v) \setminus \{u\}) \cap (N(w) \setminus \{u\}) = \emptyset$, otherwise there is a 4-cycle containing u . So $(N(v) \cup N(t)) \cap (N(w) \setminus \{u, v\}) = \emptyset$. But $v, t \notin N(v) \cup N(t)$. Hence $|N(v) \cup N(t)| \leq \nu - |N(w) \setminus \{u, v\}| - |\{t, v\}| = \nu - d(w) < \nu - \delta + 1$, contradicting the hypothesis. \square

Lemma 2 Let G be a 2-connected graph with cycle C and let $u \in V(G) \setminus V(C)$. Suppose $NC \geq \nu - \delta + 1$. If $|N_C(u)| \geq 2$ and there exist two distinct vertices $u_1^+, u_2^+ \in N_C^+(u)$ (or $u_1^-, u_2^- \in N_C^-(u)$) such that there is no edge from $\{u_1^+, u_2^+\}$ ($\{u_1^-, u_2^-\}$) to $N(u) \setminus V(C)$, then there is a cycle C' of length $|V(C)| + 1$ such that $V(C') = V(C) \cup \{u\}$.

Proof We give C an orientation. If there is an A-structure on $N_C(u)$ such that $x, y \in N_C(u)$ and $x^+ = y$. Then $C' = C^+[y, x] + xuy$ is the desired cycle. So suppose there is no A-structure on $N_C(u)$. And then $u_1^+ \neq u_2^+$ and $u_2^+ \neq u_1^+$.

Suppose there is a suc-J-structure x^+y^+ on $N_C(u)$.

Then $C' = C^+[y^+, x] + xuy + C^-[y, x^+] + x^+y^+$ is the desired

cycle. Otherwise, there is no suc-J-structure on $N_C(u)$. By the hypothesis, $(N_C(u) \setminus V(C)) \cap (N_C(u_1^+) \cup N_C(u_2^+)) = \emptyset$. So $|N_C(u_1^+) \cup N_C(u_2^+)| \leq \nu - |N_C^+(u)| - |N_C(u) \setminus V(C)| = \nu - d(u) < \nu - \delta + 1$, contradicting $NC \geq \nu - \delta + 1$. \square

Theorem 3 Let G be a 2-connected graph. If $NC \geq \nu - \delta + 1$, then G is vertex-pancyclic.

Proof By Lemma 1, when $\nu = 3, 4$, Theorem 3 holds. Thus we may assume $\nu \geq 5$.

Suppose G is not vertex-pancyclic. Let v be a vertex of G which does not lie on any cycle of length r for some r ($3 \leq r \leq \nu$). By Lemma 1, we will assume that m is the minimum number such that $3 \leq m \leq \nu - 2$ and there is a cycle C of length m in G containing v but there is no cycle of length $m+2$ in G containing v . We give C an orientation.

Claim 1 For each $u \in V(G) \setminus V(C)$, there is no edge from $N_C^+(u) \cup N_C^-(u)$ to $N_C(u) \setminus V(C)$, for otherwise a cycle of length $m+2$ containing v results.

Since G is 2-connected and $m \leq \nu - 2$, there are two distinct vertices $x, y \in V(G) \setminus V(C)$ such that $N_C(x) \neq \emptyset$ and $N_C(y) \neq \emptyset$. Now we consider the following two cases.

Case 1 There are two distinct vertices $x, y \in V(G) \setminus V(C)$ such that $|N_C(x)| \geq 2$ and $|N_C(y)| \geq 2$.

Subcase (1.1) There are two distinct vertices $x_1, x_2 \in N_C(x)$ and two distinct vertices $y_1, y_2 \in N_C(y)$ such that $|\langle x_1, x_2 \rangle \cap \langle y_1, y_2 \rangle| \leq 1$.

Subcase (1.1.1) There is an A-structure on $N_C(x)$ (or on $N_C(y)$).

If so, then there is a cycle C' of length $|VCC|+1$ such that $V(C') = VCC \cup \{x\}$. We give C' an orientation. By the assumption of Case (1.1), either $|N_C^+(y) \cap (N_C^+(y) \cup N_C^-(y))| \geq 2$ or $|N_C^-(y) \cap (N_C^+(y) \cup N_C^-(y))| \geq 2$. By Claim 1 and the argument of Lemma 2 on C' , we have a cycle C'' of length $m+2$ such that $V(C'') = V(C') \cup \{y\}$, a contradiction.

Subcase (1.1.2) There is neither an A-structure on $N_C(x)$, nor an A-structure on $N_C(y)$.

Without loss of generality, assume $|N_C(x)| \leq |N_C(y)|$. By Claim 1 and Lemma 2 (and the proof of Lemma 2), there is a suc-J-structure $x_1^+x_2^+$ on $N_C(x)$ and there is a cycle $C' = C^+[x_2^+, x_1^+] + x_1x_2 + C^-[x_2^+, x_1^+] + x_1^+x_2^+$ such that $V(C') = VCC \cup \{x\}$. We give C' an orientation such that C' and C have the same orientation on $C^+[x_2^+, x_1^+]$.

If $|N_C(y) \setminus \{x_1^+, x_2^+\}| \geq 1$, then either $|N_C^+(y) \cap (N_C^+(y) \cup N_C^-(y))| \geq 2$ or $|N_C^-(y) \cap (N_C^+(y) \cup N_C^-(y))| \geq 2$. By Claim 1 and the argument of Lemma 2 on C' , we have a cycle C'' of length $m+2$ such that $V(C'') = V(C') \cup \{y\}$, a contradiction.

So $N_C(y) \subseteq \{x_1^+, x_2^+\}$. Since there is no A-structure on $N_C(y)$, $N_C(y) = \{y_1, y_2\}$, where $y_1 \in \{x_1^+, x_2^+\}$ and $y_2 \in \{x_2^+, x_2^+\}$. But $|N_C(x)| \leq |N_C(y)|$, so $N_C(x) = \{x_1, x_2\}$. There are now two subcases.

Subcase (1.1.2.1) $x_1 = y_1$ and $x_2^+ = y_2$ (or $x_2 = y_2$ and $x_1^+ = y_1$).

By Claim 1 and Lemma 2, there is a suc-J-structure

$y_1^+ y_2^+$ on $N_C(y)$ and $C' = C^+[y_2^+, y_1^+] + y_1 y y_2 + C^-[y_2^+, y_1^+] + y_1^+ y_2^+$ is a cycle of length $m+1$ containing v . We give C' an orientation such that C' and C have the same orientation on $C^+[y_2^+, y_1^+]$. Then $|N_C^-(x) \cap (N_C^+(x) \cup N_C^-(x))| \geq 2$. By Claim 1 and Lemma 2, there is a cycle C'' of length $m+2$ containing v , a contradiction.

Subcase (1.1.2.2) $x_1^+ = y_1$ and $x_2^+ = y_2$.

Clearly, it follows that $C'' = C^+[y_2^+, x_1^+] + x_1 x x_2 + C^-[x_2^+, y_1^+] + y_1 y y_2$ is a cycle of length $m+2$ containing v , a contradiction.

Subcase (1.2) $N_C(x) = N_C(y) = \langle x_1, x_2 \rangle$.

Subcase (1.2.1) $xy \in E(G)$.

If $x_1 = x_2^+$ or $x_2 = x_1^+$, then it contradicts Claim 1.

So there is no A-structure on $N_C(x)$. By Claim 1 and Lemma 2, there is a suc-J-structure $x_1^+ x_2^+$ on $N_C(x)$. Then $C'' = C^+[x_2^+, x_1^+] + x_1 x y x_2 + C^-[x_2^+, x_1^+] + x_1^+ x_2^+$ is a cycle of length $m+2$ containing v , a contradiction.

Subcase (1.2.2) $xy \notin E(G)$.

Let $w \in \langle x_1^+, x_2^+ \rangle \setminus \langle x_1, x_2 \rangle$. Then by Claim 1, $(N(x) \cup N(y)) \cap (N(w) \setminus \langle x_1, x_2 \rangle) = \emptyset$. Also $x, y \notin N(x) \cup N(y)$. So $|N(x) \cup N(y)| \leq \nu - d(w) < \nu - \delta + 1$, a contradiction.

Case 2 There is at most one vertex $x \in V(G) \setminus V(C)$ such that $|N_C(x)| \geq 2$.

Since $m \leq \nu - 2$ and G is 2-connected, there is a vertex $y \in V(G) \setminus V(C)$ such that $|N_C(y)| = 1$. Let $N_C(y) = \langle y_C \rangle$.

Claim 2 $y_C^- y_C^+ \in E(G)$.

Otherwise, by Claim 1, $NC \leq |NC(y_C^-) \cup NC(y_C^+)| \leq \nu -$

$$|NC(y) \setminus \langle y_C \rangle| - |\langle y_C^-, y_C^+ \rangle| \leq \nu - \delta, \text{ a contradiction.}$$

Claim 3 $y_C^- w^+ \in E(G)$ if $y_C^- w \in E(G)$ for each $w \in V(C)$.

Suppose $y_C^- w \in E(G)$. If there is an edge from w^+ to $NC(y) \setminus V(C)$, then there is a cycle C'' of length $m+2$ containing v , a contradiction. Now suppose $y_C^- w^+ \notin E(G)$ and $NC(w^+) \cap (NC(y) \setminus V(C)) = \emptyset$. Then $NC \leq |NC(y_C^-) \cup NC(w^+)| \leq \nu - |NC(y) \setminus \langle y_C \rangle| - |\langle y_C^-, w^+ \rangle| \leq \nu - \delta$, a contradiction, then Claim 3 follows.

Now by Claims 2 and 3, y_C^- is adjacent to all vertices on C .

Claim 4 For any two distinct vertices $u, w \in V(C) \setminus \langle y_C \rangle$, $uw \in E(G)$.

By Claim 3, if either u or w is y_C^- then the claim follows. So $u \neq y_C^-$ and $w \neq y_C^-$. Since $y_C^- u^-, y_C^- w^- \in E(G)$ and there is no cycle of length $m+2$ containing v , it must be the case that neither u nor w is adjacent to any vertex in $NC(y) \setminus \langle y_C \rangle$. Hence $NC \leq |NC(u) \cup NC(w)| \leq \nu - |NC(y) \setminus \langle y_C \rangle| - |\langle u, w \rangle| \leq \nu - \delta$ unless $uw \in E(G)$.

Consequently the graph induced by $V(C) \setminus \langle y_C \rangle$ is complete etc.

Let D be the component of $G - V(C)$ containing y . Let $P = y_1 y_2 \dots y_k$ be a shortest path in D with $y_1 = y$ and y_k adjacent to a vertex in $V(C) \setminus \langle y_C \rangle$. Since $|N_C(y)| = 1$, $k \geq 2$.

If $k = 2$, a cycle C'' of length $m+2$ such that $V(C'') = V(C) \cup \{y_1, y_2\}$ results, a contradiction.

If $k = 3$ and if either $vy_1 \in E(G)$ or $vy_3 \in E(G)$, then it follows that v lies on a cycle of length $m+2$ which contains $\{y_1, y_2, y_3\}$ and all vertices but a vertex of C , a contradiction.

Otherwise $vy_1, vy_3 \notin E(G)$.

If $|N_C(y_1) \cup N_C(y_3)| \geq 3$, then $|N_C(y_3) \setminus \{y_C\}| \geq 2$, and v lies on a cycle of length $m+2$ which contains $\{y_1, y_2, y_3\}$ and misses a vertex in $N_C(y_3)$.

Hence we may suppose $|N_C(y_1) \cup N_C(y_3)| \leq 2$. If $|V(C)| \geq 4$, $v \in \{y_C^+, y_C^-\}$ and $N_C(y_3) \setminus \{y_C\} \subseteq \{y_C^+, y_C^-\}$, then v lies on a cycle of length $m+2$ which contains $\{y_1, y_2, y_3\}$ and misses exactly one vertex in $V(C)$, a contradiction. So suppose this is not the case. Then $N(v) \cap (N_C(y_1) \setminus V(C)) = \emptyset$, for otherwise the case that $k = 2$ results. Additionally, $N(v) \cap (N_C(y_3) \setminus V(C)) = \emptyset$, otherwise, by a similar argument as in the case that $k = 2$, v lies on a cycle of length $m+2$. Hence $NC \leq |N_C(y_1) \cup N_C(y_3)| \leq \nu - |N(v) \setminus (N_C(y_1) \cup N_C(y_3))| - |\{y_1, y_3\}| \leq \nu - d(v) \leq \nu - \delta$, a contradiction.

Suppose $k \geq 4$. Let $u \in V(CD) \setminus (N(y) \cup \{y, x\})$. Since $|N_C(u)| \leq 1$, we can find a vertex $w \in V(CD) \setminus \{y_C\}$ such that $uw \notin E(G)$ and $N(w) \cap (N_C(y) \setminus \{y_C\}) = \emptyset$. If $N(u) \cap (N_C(y) \setminus \{y_C\}) = \emptyset$, then $NC \leq |N_C(u) \cup N_C(w)| \leq \nu - |N_C(y) \setminus \{y_C\}| - |u, w| < \nu - \delta$, a contradiction. Hence $N(u) \cap (N_C(y) \setminus \{y_C\}) \neq \emptyset$ for all $u \in V(CD) \setminus (N_C(y) \cup \{y, x\})$. So $k = 4$ and $x = y_4$. And x is adjacent to all vertices in $V(CD) \setminus \{y_C\}$, otherwise we take a vertex w in $V(CD)$ such that $xw \notin E(G)$, then

$NC \leq |NCx \cup NCw| \leq \nu - |NCy \setminus \{y_c\}| - |\{x, w\}| < \nu - \delta$, a contradiction.

We now have a path $P = y_1 y_2 y_3 y_4$ in D , where $y = y_1$ and $x = y_4$. By the assumption on P , $N_C(y_3) \setminus \{y_c\} = \emptyset$. Let $u \in V(C) \setminus \{y_c\}$.

Suppose $N_C(y_3) \cap (NCu \setminus \{y_c, y_4\}) = \emptyset$. Obviously, $N_C(y_1) \cap (NCu \setminus \{y_c\}) = \emptyset$. Thus $NC \leq |N_C(y_1) \cup N_C(y_3)| \leq \nu - |NCu \setminus \{y_c, y_4\}| - |\{y_1, y_3\}| \leq \nu - d(u) < \nu - \delta + 1$, a contradiction.

Finally, suppose $w \in N_C(y_3) \cap (NCu \setminus \{y_c, y_4\})$. Then either $w \in N_C(y) \setminus \{y_c\}$ or $NC(w) \cap (NCy) \setminus \{y_c\} \neq \emptyset$. So we have a shorter path than P in D , contradicting the assumption on P . The proof of this theorem is complete. \square

Remark 1 $K_{m,m}$ shows that the bound on NC in Theorem 3 is the best possible.

Remark 2 Recently the authors learned that Lin and Song [4] have obtained an analogous result for edge pancyclicity.

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